

Atmel ATPL210A

PRIME compliant Power Line Communications SoC

DATASHEET

Features

- Core
 - ADD8051C3A enhanced 8051 core
 - Speedups up to x5 vs. standard 8051 microcontroller
- Modem
 - Power Line Carrier Modem for 50 and 60 Hz mains
 - 97-carrier OFDM PRIME compliant
 - Baud rate Selectable: 21400 to 128600 bps
 - Differential BPSK, QPSK, 8-PSK modulations
- Memories
 - 32Kbytes on-chip SRAM
 - Up to 256Kbytes external SRAM
- In-circuit serial flash programming
- Auto boot-loading program from serial flash
- Automatic Gain Control and signal amplitude tracking
- Embedded on-chip DMAs
- Automatic code encryption during boot loading
- Media Access Control
 - · Viterbi decoding and CRC PRIME compliant
 - 128-bit AES encryption
 - Channel sensing and collision pre-detection
- Peripherals
 - Two 2-wire UARTs
 - Two SPI. SPI to serial flash and External RTC. Buffered SPI to external metering IC
 - Programmable Watchdog
 - Up to 14 I/O lines
- Package
 - 120-lead LQFP, 14 x 14 mm, pitch 0.4 mm
 - Pb-free and RoHS compliant
- Typical Applications
 - Automated Meter Reading (AMR) & Advanced Meter Management (AMM)
 - Street lighting
 - Home Automation

Description

The ATPL210A is a Power Line Communications System on Chip, which implements a full PRIME compliant PLC modem. It includes an enhanced 8051 microcontroller (IP core ADD8051C3A), a Medium Access Controller (MAC) (IP core ADD1221) and a Modem circuit (IP core ADD1321) for power line medium using OFDM modulation compatible with PRIME specifications.

ATPL210A is oriented to high performance & robust AMR systems. The ATPL210A is designed to be used by meter manufacturers to provide a low cost and compact solutions for AMR (Automated Meter Reading) & AMM (Advanced Meter Management) systems using narrow band power line communications.

This device has been developed to reduce CPU computational load in PLC systems running PRIME protocols. ATPL210A includes all necessary resources to be used as main controller in metering applications, and allows an external device to communicate according to PLC PRIME specifications.



Table of Contents

1.	Block Diagram						
2.	Pac 2.1 2.2	kage and Pinout 120-Lead LQFP Package Outline 120-Lead LQFP Pinout	8				
3.	Pin	Description	11				
4	PRI	ME overview	18				
т.	4.1	Kev Features					
	4.1	PRIME physical layer overview					
	7.2	4.2.1 Main features					
		4.2.2 Orthogonal Frequency Division Multiplexing (OFDM)					
	4.3	PRIME MAC layer overview					
		4.3.1 Base Node					
		4.3.2 Service Node					
	4.4	Convergence Layer	21				
5.	Proc	cessor and Architecture	23				
	5.1	ADD8051C3A Microcontroller Description	23				
	5.2	Core Pinout Description					
	5.3	Memory Organization					
		5.3.2 Program Memory					
		5.3.3 Extended Addressing					
		5.3.4 Data Memory					
	5.4	5.3.5 SFR Registers Instruction Set					
	5.4	5.4.1 Program Status Word					
		5.4.2 Addressing Modes					
		5.4.3 Arithmetic Instructions					
		5.4.4 Logical Instructions					
		5.4.5 Data Transfer Instructions					
		5.4.6 Boolean Instructions					
	5.5	5.4.7 Jump Instructions CPU Timing					
	5.5	5.5.2 Reset Modes					
		5.5.3 Power Saving Modes					
		5.5.4 Idle Mode					
		5.5.5 Power-Down Mode					
	5.6	Interrupts					
		5.6.1 Interrupt Enabling					
		5.6.2 Interrupt Priorities					
	5.7	5.6.3 Interrupt Handling I/O Ports					
	5.7	5.7.2 I/O Configurations					
		5.7.3 Read-Modify-Write Feature					
		5.7.4 Accessing External Memories					
	5.8	Debug Mode					
6.	Time	ers 55					
	6.1	Timer 0 and Timer 1	55				
	0.1	6.1.1 Timer Mode 0					

	6.1.1	Timer Mode 0	56
	6.1.2	Timer Mode 1	57
	6.1.3	Timer Mode 2	57
	6.1.4	Timer Mode 3	57
6.2	Timer 2.		58
	6.2.2	Capture mode	59



	6.3	 6.2.3 Auto-Reload mode 6.2.4 Clock-Out mode 6.2.5 Baud rate Generator mode Watchdog (timer 3) 	60 61
7.			
1.		ndard Serial Interfaces	
	7.1	Serial Port modes 7.1.1 Mode 0 (shift register mode)	
		7.1.2 Mode 1 (8-bit UART)	
		7.1.3 Modes 2 and 3 (9-bit UART)	65
	7.2	Serial Port Timers (Timers 11 &12)	
8.	Seria	al Peripheral Interfaces SPI0 and SPI1	69
	8.1	SPI description	69
		8.1.2 SPI clock phase, polarity and operation	71
	8.2	8.1.3 SPI0 write collision SPI Modes	
	0.2 8.3	SPI Modes	
0		·	
9.	Penp	pheral Registers	
10	. Boot	t Loader	78
	10.2	Pin Description	79
	10.3	5 5	
		10.3.1 In-System programming	
	10.4	10.3.2 SPI Flash programming System startup	
	10.1	10.4.2 Encrypted firmware requirements	
		10.4.3 Supported Devices	83
	105	10.4.3.1 Serial Number Device	
	10.5	Boot loader registers	
	_		
11		a Memory Extension Block	
	11.1		
	11.2	DMEB Configuration Registers 11.2.2 TABLE_ELEMENT_SIZE registers	
		11.2.3 TABLE_INDEX registers	
		11.2.4 TABLE_ELEMENT_INIT registers	
	44.0	11.2.5 TABLE_OFFSET registers	
	11.3	Code Example	92
12	. Direc	ct Memory Access (DMA)	96
	12.1	DMA Management	
	12.2	DMA Hardware	
	12.3 12.4		
	12.4	•	
	12.6	DMA Flags	
	12.7		
	12.8	Physical DMA channels	
	12.9	PHY_RX DMA channel	
		0 PHY_TX DMA channel	
	12.11	1 DMA Configuration Registers 12.11.1 PHY_SFR Register	
		12.11.1 DMA_SFR Register	
		12.11.3 PHY_TX Registers	104
		12.11.4 PHY_RX Registers	
		12.11.5 ADDR_PHY_INI_RX Registers 12.11.6 ADDR_PHY_INI_TX Registers	



			ADDR_DMA1_SRC Registers	
			ADDR_DMA1_DTN Registers	
			SIZE_DMA1_SRC Registers	
		12.11.10)SIZE_DMA1_DTN Registers SIZE_RQ_DMA1 Registers	112
			2DMA1_CHN_CTL Register	
			ADDR_DMA2_SRC Registers	
			ADDR_DMA2_DTN Registers	
			SIZE_DMA2_SRC Registers	
			SIZE_DMA2_DTN Registers	
			SIZE_RQ_DMA2 Registers	
13.	ATPL	_210 M	AC Coprocessor 1	20
	13.2		edundancy Check (CRC)	
	13.3		ed Encryption Standard	
	13.4		RIME Software Stack	
	13.5	MAC Co	processor Registers	125
		13.5.1	SNA Registers	125
		13.5.2	VITERBI_BER_HARD Register	
		13.5.3	VITERBI_BER_SOFT Register	
		13.5.4	ERR_CRC32_MAC Registers	
		13.5.5 13.5.6	ERR_CRC8_MAC Registers ERR_CRC8_AES Registers	
		13.5.7	ERR_CRC8_MAC_HD Registers	
		13.5.8	ERR_CRC8_PHY Registers	
		13.5.9	FALSE_DET_CONFIG Register	
		13.5.10	FALSE_DET Registers	
			MAX_LEN_DBPSK Register	
		13.5.12	MAX_LEN_DBPSK_VTB Register	136
			MAX_LEN_DQPSK Register	
			MAX_LEN_DQPSK_VTB Registers	
			MAX_LEN_D8PSK Registers MAX_LEN_D8PSK_VTB Register	
			AES_PAD_LEN Register	
			AES_DATA_IN Registers	
			AES_DATA_OUT Registers	
		13.5.20	KEY_PERIPH Registers	144
			KEY_PHY Registers	
		13.5.22	AES_SFR Register	146
14.	ATPL	_210 PF	RIME PHY Layer 1	47
			0 PHY Layer	
		14.1.2	Transmission and Reception branches	
		14.1.3	Carrier Detection.	
		14.1.4	Analog Front End control	149
			14.1.4.1 Power Supply Sensing: VSENSE and PSENSE	
			14.1.4.2 Gain Control	
			14.1.4.3 Line Impedance Control	
	14.2	DHV nau	14.1.4.4 TxRx Control	
	14.3		otocol Data Unit (PPDU) Format	
	14.4		rvice Specification	
	14.5		/er registers	
	14.5	14.5.1	PHY_SFR Register	
		14.5.2	SYS_CONFIG Register	
		14.5.3	PHY_CONFIG Register	
		14.5.4	ATTENUATION Register	
		14.5.5	ATT_CHIRP Register	
		14.5.6	ATT_SIGNAL Register	
		14.5.7	TX_TIME Registers.	
		14.5.8	TIMER_FRAME Registers	101





	14.5.10	TIMER_BEACON_REF Registers
	14.5.12	RSSI_AVG Register
		CINR_MIN Register
		CINR_AVG Register
		CINR_MAX Register
		EVM_HEADER Registers
		EVM_PAYLOAD Registers
		EVM_PAYLOAD_ACUM Registers
		RMS_CALC Register
	14.5.22	VSENSE_CONFIG Register
		NUM_FAILS Register
	14.5.24	TTRANS Register
		AGC0_KRSSI Register
		ZERO_CROSS_TIME Registers
	14.5.28	ZERO_CROSS_CONFIG Register
	14.5.29	PSENSECYCLES Registers
		MEAN Registers
		PMAX Registers
	14.5.32	P_TH Registers
		MAXPOT Registers
		NUMCYCLEŠ Register
	14.5.36	A_NUMMILIS Register
		EMIT_CONFIG Register
		AFE_CTL Register
	14.5.40	PHY_ERRORS Registers
	14.5.41	FFT_MODE Registers 194
	14.5.42	AGC_CONFIG Register
	14.5.43	SAT_TH Registers
		AGC1_TH Registers
		AGC_PADS Register
Elect		aracteristics
15.1	Absolute	Maximum Ratings
15.2	Recomm	nended Operating Conditions
15.3	DC Char	acteristics
		V-I curves
15.4		onsumption
15.5		Data
15.6 15.7		r
10.7	r ower o	
Mech	nanical (Characteristics
Reco		led mounting conditions 212
17.1		ns of Standard Reflow212
17.2	Manual S	Soldering
Orde	ring Info	prmation
Revis	sion His	tory 215

15.

16.

17.

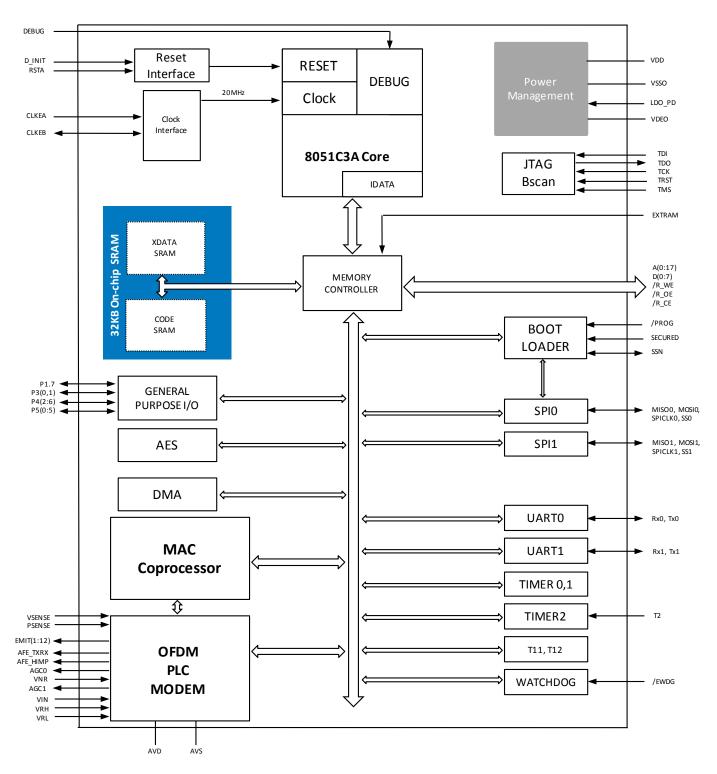
18.

19.



1. Block Diagram



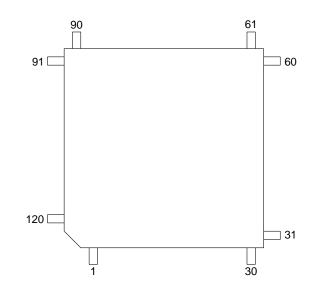




2. Package and Pinout

2.1 120-Lead LQFP Package Outline

Figure 2-1. Orientation of the 120-Lead Package







2.2 120-Lead LQFP Pinout

Table 2-1. ATPL210A 120-Lead LQFP pinout

PinNo	Pin Name	I/O	l(mA)	Res	HY	
1	A17	0	±2	-	-	
2	A9	0	±2	-	-	
3	GND	Р	-	-	-	
4	VCC	Р	-	-	-	
5	/R_WE	0	±2	PU	-	
6	D4	I/O	±2	-	-	
7	D3	I/O	±2	-	-	
8	D5	I/O	±2	-	-	
9	D2	I/O	±2	-	-	
10	D6	I/O	±2	-	-	
11	D1	I/O	±2	-	-	
12	D7	I/O	±2	-	-	
13	D0	I/O	±2	_	-	
14	VCC	Р	-	-	-	
15	GND	Р	-	-	-	
16	VDD	Р	-			
17	/R_OE	_OE 0 ±		-	-	
18	/R_CE	0	±2	-	-	
19	A8	0	±2	2 -		
20	A0	0	±2			
21	A7	0	±2			
22	A1	0	±2			
23	A6	0	±2			
24	A2	0	±2			
25	A5	0	±2	-	-	
26	GND	Р	-	-	-	
27	VCC	Р	-	-	-	
28	A3	0	±2	-	-]	
29	A4	0	±2	-	-	
30	P4.5/MISO1	I/O	±2	PU	-	
31	P4.4/MOSI1	I/O	±2	PU	-	
32	P4.3/SPICLK1	I/O	±2	PU	-	
33	P4.2/SS1	I/O	±2	PU	-	
34	P3.0/RxD0	I/O	±2	PU	-	
35	P3.1/TxD0	I/O	±2	PU	-	
36	VCC	Р	-	-	-	
37	GND	Р	-	-	-	
38	EMIT.1	0	±Χ	-	-	

PinNo	Pin Name	I/O	l(mA)	Res	HY
39	EMIT.2	0	±X -		-
40	EMIT.3	O ±X -		-	
41	EMIT.4	0	±Χ	-	-
42	VCC	Р	-	-	-
43	GND	Р	-	-	-
44	EMIT.5	0	±Χ	-	-
45	EMIT.6	0	±Χ	-	-
46	EMIT.7	0	±Χ	-	-
47	EMIT.8	0	±Χ	-	-
48	VCC	Р	-	-	-
49	GND	Р	-	-	-
50	EMIT.9	0	±Χ	-	-
51	EMIT.10	0	±Χ	-	-
52	EMIT.11	0	±Χ	-	-
53	EMIT.12	0	±Χ	-	-
54	VCC	Р	P		-
55	55 GND P		-	-	-
56	AFE_HIMP	0	±Χ	-	-
57	AFE_TXRX	0	±Χ	-	-
58	VSENSE	I	-	-	Y
59	PSENSE	I	-	-	Y
60	VNR	I	-	- Y	
61	TDI	I	I - PU		-
62	TMS	I	-	PU	-
63	TDO	0	±2	-	-
64	GND ⁽¹⁾	Р	-	-	-
65	GND	Р	-	-	-
66	VCC	Р	-	-	-
67	TRST	I	-	PU	-
68	тск	I	-	-	
69	RSTA	I	-	PD	Y
70	D_INIT	I	-	PD	Y
71	GND	Р	-	-	-
72	VCC	Р	-	-	-
73	GND	Р	-	-	-
74	VDD	Р	-	-	-
75	LDO_PD	I	-	-	-
76	VSS0	Р	-	-	-



PinNo	Pin Name	I/O	I(mA)	Res	HY
77	VDE0	0 P		-	-
78	VDE0 P		-	-	
79	GND	Р	-	-	-
80	GND	Р	-	-	-
81	VCC	Р	-	-	-
82	CLKEA	Ι	-	-	-
83	GND	Р	-	-	-
84	CLKEB	I/O	-	-	-
85	85 VCC P -		-	-	
86	86 /EWDG I -		PD	Y	
87	87 DEBUG I		-	PD	Y
88	88 EXTRAM I -		-	Y	
89	/PROG I - PU		PU	Y	
90	SECURED	I	I - PD Y		Y
91	P5.4/RxD1	I/O	±2	PU	-
92	P5.5/TxD1	I/O	±2	PU	-
93	NC	-	-	-	-
94	P4.6/T2/AGC1	I/O	±Χ	PU	-
95	AGC0	0	±Χ	-	-
96	GND	Р	-	-	-
97	VCC	Р	-		
98	AVS2	Р	-	-	-

PinNo	Pin Name	I/O	I(mA)	Res	HY
99	AVD2	Р	-	-	-
100	AVS1	Р	-	-	-
101	AVD1	Р	-	-	-
102	VRH	I	-	-	-
103	VIN	I	-	-	-
104	VRL	I	-	-	-
105	GND	Р	-	-	-
106	VCC	Р	-	-	-
107	P5.0/SS0	I/O	±2	PU	-
108	P5.3/MISO0	I/O	±2	PU	-
109	P5.2/MOSI0	I/O	±2	PU	-
110	P5.1/SPICLK0	I/O	±2	PU	-
111	P1.7/SSN	I/O	±2	PU	-
112	A13	0	±2	-	-
113	A14	0	±2	-	-
114	GND	Р	-	-	-
115	VCC	Р	-	-	-
116	A12	0	±2	-	-
117	A15	0	±2	-	-
118	A11	0	±2	-	-
119	A16	0	±2	-	-
120	A10	0	±2	-	-

Notes: 1. Mandatory to be tied down

I/O=pin direction:I=input, O=Output, P=PowerI(mA)=nominal current:+=source, -=sink, X=fixed by external resistor

RES=pin pullup/pulldown resistor: PU=pullup, PD=pulldown,

HY=Input Hysteresis



3. Pin Description

Table 3-1. Pin Description List

Pin Number	Pin Name	Туре	Comments
1, 2, 19, 20, 21, 22, 23, 24, 25, 28, 29, 112,113, 116, 117, 118, 119, 120	A(0:17)	Output	 External Memory data/program Address bus. This bus is enabled when the system is configured to work with an external memory device (EXTRAM pin = '1'). If the system is configured to work with the on-chip SRAM configuration (EXTRAM pin = '0'), then these pins output is '0'.
3, 15, 26, 37, 43, 49, 55, 64, 65, 71, 73, 79,80, 83, 96, 105, 114	GND	Power	Digital Ground
4, 14, 27, 36, 42, 48, 54, 66, 72, 81, 85, 97,106, 115	VCC	Power	3.3v digital supply. Digital power supply must be decoupled by external capacitors
5	/R_WE	Output	 External Memory Write Enable signal. This output is enabled when the system is configured to work with an external memory device (EXTRAM='1'). If the system is configured to work with the on-chip SRAM configuration (EXTRAM='0'), then this output must be left unconnected. /R_WE is active low Internal configuration: 33kΩ typ. pull-up resistor
6, 7, 8, 9, 10, 11, 12, 13	D(0:7)	I/O	 External Memory Data bus. This bus is enabled when the system is configured to work with an external memory device (EXTRAM='1'). If the system is configured to work with the on-chip SRAM (EXTRAM='0'), then these pins are in Hi-Z.
16, 74.	VDD	Power	LDO Power Output. A capacitor in the range 0.1µF-10µF must be connected to each pin
17	/R_OE	Output	 External Memory Output Enable signal. This output is enabled when the system is configured to work with an external memory device (EXTRAM='1'). If the system is configured to work with the on-chip SRAM configuration (EXTRAM='0'), then this pin output is '1'. /R_OE is active low
18	/R_CE	Output	 External Memory Chip Enable signal. This output is enabled when the system is configured to work with an external memory device (EXTRAM='1'). If the system is configured to work with the on-chip SRAM configuration (EXTRAM='0'), then this pin output is '1'. /R_CE is active low



Pin Number	Pin Name	Туре	Comments
			Microcontroller port 4.5 / SPI1 Master In Slave Out.
30	P4.5/MISO1	I/O	 When configured as P4.5, this pin is a pseudo- bidirectional microcontroller I/O port
30	F4.5/MISO1	1/0	 When configured as MISO1, this pin is the SPI1 Master In Slave Out
			 Internal configuration: 33kΩ typ. pull-up resistor
			Microcontroller port 4.4 / SPI1 Master Out Slave In
31	P4.4/MOSI1	I/O	 When configured as P4.2, this pin is a pseudo- bidirectional microcontroller I/O port
51		1/0	 When configured as MOSI1, this pin is the SPI1 Master Out Slave In
			 Internal configuration: 33kΩ typ. pull-up resistor
			Microcontroller port 4.3 / SPI1 Clock
32	P4.3/SPICLK1	I/O	 When configured as P4.3, this pin is a pseudo- bidirectional microcontroller I/O port
			 When configured as SPICLK1, this pin is the SPI1 clock signal
			 Internal configuration: 33kΩ typ. pull-up resistor
	P4.2/SS1		Microcontroller port 4.2 / SPI1 Slave Select
33		I/O	 When configured as P4.2, this pin is a pseudo- bidirectional microcontroller I/O port
55			 When configured as SS1, this pin is the SPI1 Slave Select. Active low
			 Internal configuration: 33kΩ typ. pull-up resistor
			Microcontroller port 3.0 / Standard Serial Port 0 Rx
34	P3.0/RxD0	I/O	 When configured as P3.0, this pin is a pseudo- bidirectional microcontroller I/O port
04	10.0/1000		 When configured as RxD0, this pin is the digital input of the asynchronous standard serial port 0
			Internal configuration: $33k\Omega$ typ. pull-up resistor
			Microcontroller port 3.1 / Standard Serial Port 0 Tx
35	P3.1/TxD0	I/O	 When configured as P3.1, this pin is a pseudo- bidirectional microcontroller I/O port
			 When configured as TxD0, this pin is the digital output of the asynchronous standard serial port 0
			Internal configuration: 33kΩ typ. pull-up resistor
38, 39, 40, 41, 44, 45, 46, 47, 50, 51, 52, 53	EMIT(1:12)	Output	PLC Transmission ports ⁽¹⁾

Pin Number	Pin Name	Туре	Comments
			Analog Front-End High-Impedance
56	AFE_HIMP	Output	• This digital output is used by the chip to select between low-impedance and high-impedance transmission branch (when working with a "two half-H-bridge branches" analog front end configuration). This way, the system adapts its transmission external circuitry to the net impedance, improving transmission behavior. The polarity of this pin can be inverted by hardware. Please refer to the Reference Design for further information.
			Analog Front-End Transmission/Reception
57	AFE_TxRx	Output	 This digital output is used to select between external Transmission and Reception branches. The suitable value depends on the external circuitry configuration. The polarity of this pin can be inverted by hardware. Please refer to the Reference Design for further information.
			Voltage Level Sensing
58	VSENSE	Input	 This input tracks the voltage level in the power supply to avoid power supply malfunction.
			Power Level Sensing
59	PSENSE	Input	 This input tracks the power level in the power supply to avoid power supply malfunction.
			Zero Crossing Detection Signal
60	VNR	Input	 This input detects the zero-crossing of the mains voltage, needed to determine proper switching times. Depending on whether an isolated or a non-isolated power supply is being used, isolation of this pin should be taken into account in the circuitry design. Please refer to the Reference Design for further information.
61	TDI ⁽²⁾	Input	Test Data In
			 Internal configuration: 33kΩ typ. pull-up resistor
62	TMS ⁽²⁾	Input	 Test Mode Select Internal configuration: 33kΩ typ. pull-up resistor
63	TDO ⁽²⁾	Output	Test Data out
67	TRST ⁽²⁾	Input	 Test Reset Internal configuration: 33kΩ typ. pull-up resistor
68	TCK ⁽²⁾	Input	Test Clock
			Asynchronous reset
69	RSTA	Input	 RSTA is a digital input pin used to perform a hardware reset of the ASIC
			RSTA is active high
			 Internal configuration: 33kΩ typ. pull-down resistor



Pin Number	Pin Name	Туре	Comments
70	D_INIT	Input	 Initialization Signal During power-on, D_INIT should be released before asynchronous reset signal RSTA, in order to ensure proper system start up. Not minimum time is required between both releases, Δt>0 D_INIT is active high Internal configuration: 33kΩ typ. pull-down resistor
75	LDO_PD	Input	LDO Power-down This digital input is used to put the internal linear regulator into power down mode '0': Power down mode disabled '1': Power down mode enabled
76	VSS0	Power	LDO ground
77, 78	VDE0	Power	LDO 3.3v power supply
82	CLKEA ⁽³⁾	Input	 External clock reference CLKEA must be connected to one terminal of a crystal (when a crystal is being used) or tied to ground if a compatible oscillator is being used
84	CLKEB ⁽³⁾	I/O	 External clock reference CLKEB must be connected to one terminal of a crystal (when a crystal is being used) or to one terminal of a compatible oscillator (when a compatible oscillator is being used)
86	/EWDG	Input	 Watchdog enable /EWDG digital input enables watchdog timer. This pin is internally connected to the /EW signal of the ADD8051C3A microcontroller '0': Watchdog timer enabled '1': Watchdog timer disabled Internal configuration: 33kΩ typ. pull-down resistor
87	DEBUG	Input	 Debug mode enable DEBUG digital input is internally connected to DBG signal of the ADD8051C3A microcontroller, and it is intended to implement software debugging tools '0': Debug mode disabled '1': Debug mode enabled Internal configuration: 33kΩ typ. pull-down resistor



Pin Number	Pin Name	Туре	Comments
			SRAM device selection
88	EXTRAM	Input	 This digital input allows selecting between working with an external SRAM device or with the internal SRAM embedded on chip. '0': 32Kbytes on-chip SRAM selected. '1': External memory device selected. (The on-chip SRAM is not accessible when working in this
			mode).
			SPI Flash programming pin
89	/PROG	Input	 /PROG digital input is read during power up and sets the system into "execution" mode or "serial flash programming" mode
			'0': Serial flash programming mode
			'1': Execution mode (normal mode)
			 Internal configuration: 33kΩ typ. pull-up resistor
		Input	Encryption enable
90	SECURED ⁽⁴⁾		 SECURED digital input enables encrypted firmware storage and execution when the board configuration supports it
			 '0': Encrypted storage/execution disabled
			 '1': Encrypted storage/execution enabled
			 Internal configuration: 33kΩ typ. pull-down resistor
			Microcontroller port 5.4 / Standard Serial Port 1 Rx
91	P5.4/RxD1	1/0	 When configured as P5.4, this pin is a pseudo- bidirectional microcontroller I/O port
91	F3.4/KXD1	1/0	 When configured as RxD1, this pin is the digital input of the asynchronous standard serial port 1
			 Internal configuration: 33kΩ typ. pull-up resistor
			Microcontroller port 5.5 / Standard Serial Port 1 Tx
02			 When configured as P5.5, this pin is a pseudo- bidirectional microcontroller I/O port
92	P5.5/TxD1	I/O	 When configured as TxD1, this pin is the digital output of the asynchronous standard serial port 1
			 Internal configuration: 33kΩ typ. pull-up resistor
93	NC		No connect



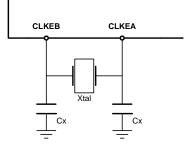
Pin Number	Pin Name	Туре	Comments			
			Microcontroller port 4.6/T2/AGC1			
			 When configured as P4.6, this pin is a pseudo- bidirectional microcontroller I/O port 			
94	P4.6/T2/AGC1	I/O	 When configured as T2, this pin works as the external T2 pin described in Timer2 section 			
			 When configured as AGC1, this pin is managed by AGC hardware logic to drive external circuitry if input signal attenuation is needed 			
			 Internal configuration: 33kΩ typ. pull-up resistor 			
			Automatic Gain Control 0			
95	AGC0	Output	 This digital output is managed by AGC hardware logic to drive external circuitry if input signal attenuation is needed 			
98, 100	AVS1, AVS2	Power	Analog ground			
99, 101	AVD1, AVD2	Power	3.3v analog power			
102	VRH ⁽⁵⁾	Input	Analog input high voltage reference			
103	VIN ⁽⁵⁾	Input	Direct-analog input voltage			
104	VRL ⁽⁵⁾	Input	Analog input low voltage reference			
			Microcontroller port 5.0 / SPI0 Slave Select			
107	P5.0/SS0	I/O	 When configured as P5.0, this pin is a pseudo- bidirectional microcontroller I/O port 			
107	F3.0/330	1/0	 When configured as SS0, this pin is the SPI0 Slave Select. Active low 			
			 Internal configuration: 33kΩ typ. pull-up resistor 			
			Microcontroller port 5.3 / SPI0 Master In Slave Out			
108	P5.3/MISO0	I/O	 When configured as P5.3, this pin is a pseudo- bidirectional microcontroller I/O port 			
100	1 3.3/101000	1/0	 When configured as MISO0, this pin is the SPI0 Master In Slave Out 			
			 Internal configuration: 33kΩ typ. pull-up resistor 			
			Microcontroller port 5.2 / SPI0 Master Out Slave In			
109	P5.2/MOSI0	I/O	 When configured as P5.2, this pin is a pseudo- bidirectional microcontroller I/O port 			
109	F 3.2/100010	1/0	 When configured as MOSI0, this pin is the SPI0 Master Out Slave In 			
			 Internal configuration: 33kΩ typ. pull-up resistor 			



Pin Number	Pin Name	Туре	Comments
			Microcontroller port 5.1 / SPI0 Clock
110	P5.1/SPICLK0	1/0	 When configured as 5.1, this pin is a pseudo- bidirectional microcontroller I/O port
110	F3.1/SFICERU		 When configured as SPICLK0, this pin is the SPI0 clock signal
			 Internal configuration: 33kΩ typ. pull-up resistor
			Microcontroller port 1.7 / Silicon Serial Number
		I/O	 When configured as P1.7, this pin is a pseudo- bidirectional microcontroller I/O port
111	P1.7/SSN		 This pin is the digital input used to read a Serial number if a valid SSN device is being used. This Serial Number is used for encryption purposes. Precaution should be taken if used as generic control port since it searches for a Silicon Serial Number device at start-up and could put out undesirable transient values
			 Internal configuration: 33kΩ typ. pull-up resistor

Notes: 1. Different configurations allowed depending on external topology and net behavior

- 2. This pin is part of the JTAG Boundary Scan interface and is only used for boundary scan purposes
- 3. The crystal should be located as close as possible to CLKEA and CLKEB pins. Recommended value for Cx is 18pF. This value may depend on the specific crystal characteristics



- 4. See supported devices section 10.4.3
- 5. See Reference Design for suitable values



4. **PRIME overview**

The PRIME (PoweRline Intelligent Metering Evolution) initiative was originally conceived as an answer to the need for a future-proof, cost-effective Automatic Meter Management (AMM) solution. During its evolution in the last two years, it has evolved into a solution for an entire Smart Grid environment which will contribute definitively to energy efficiency improvement and ultimately to addressing the pressing issue of climate change.

The overall performance of the system fully complies with requirements for a new, stable metering infrastructure that is set to become a fundamental part of the Smart Grids. The revolutionary openness of the PRIME solution and its focus on interoperability as a way to decrease costs in a competitive framework, has been arising an unprecedented level of interest within the industry and markets. For this very reason a decision was taken by several of the most committed parties to set up the PRIME Alliance.

PRIME Alliance remains unique in that it shuns proprietary technologies and is the result of an open, cost-oriented effort among many different partners seeking to develop a future market for common benefit. Royalties and patents are definitely not part of the core PRIME specification. Ownership of the PRIME specification is uniquely public.

4.1 Key Features

PRIME defines lower layers of a PLC narrowband data transmission system over the electric grid. All the system has been created to be low cost and high performance.

Figure 4-1 below depicts the proposed communication layers and the scope of the specification. The proposed reference model is based on IEEE Std. 802.16 protocol layering.

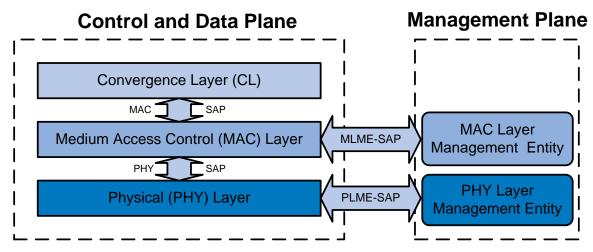


Figure 4-1. PRIME layers

The service-specific **Convergence Layer** (CL) classifies traffic associating it with its proper MAC connection. This layer performs the mapping of any kind of traffic to be properly included in MAC SDUs (Service Data units). It may also include payload header suppression functions. Multiple Convergence sub-layers s are defined in order to accommodate different kinds of traffic into MAC SDUs.

The **MAC layer** provides core MAC functionalities of system access, bandwidth allocation, connection management and topology resolution. It has been defined for a connection oriented Master-Slave environment, and optimized for low voltage power line environments.

The **PHY layer** transmits and receives MAC PDUs (Protocol Data Units) between Neighbor Nodes. It is based on OFDM multiplexing in CENELEC A band and reaches up to 130 kbps raw data rate.



PRIME specifications take advantages of state of the art technologies and adapt them to the needed requirements, simplifying processes, overheads and others, to ensure performance, interoperability between devices and different implementations of elements in the system.

4.2 PRIME physical layer overview

4.2.1 Main features

PRIME PHY layer is designed to transmit and receive over power lines which were originally devised for distribution of power at 50-60Hz AC. The use of this medium for communications at higher frequencies presents some challenging technical problems:

- Distribution networks are usually made of a random variety of conductor types, and terminating into loads of different impedances. Such a network has an amplitude and phase response that varies widely with frequency. Furthermore, the channel characteristics shall also vary with time as the loads on the network change.
- Interference also affects power lines. Electric appliances with different kind of motors, switching power supplies and halogen lamps produce impulse noise that reduces the reliability of communication signals. Due to attenuation, the noise is also location dependent.

PRIME PHY layer uses a combination of approaches that ultimately allow for robust high speed, low cost communications over power lines. A simple yet powerful scheme which is based on adaptively modulated Orthogonal Frequency Division Multiplexing (OFDM), along with forward error correction and data interleaving.

A block diagram representation of a PHY transmitter is shown below in Figure 4-2:

Figure 4-2. PHY transmitter



On the transmitter side, the PHY layer receives its inputs from the Media Access Control layer. If decided by higher layers, the PHY frame after the CRC block is convolutionally encoded and interleaved; however, it will always be scrambled). The output is differentially modulated using a DBPSK, DQPSK or D8PSK scheme. The next step is OFDM, which comprises the IFFT (Inverse Fast Fourier Transform) block and the cyclic prefix generator.

4.2.2 Orthogonal Frequency Division Multiplexing (OFDM)

One of the main novelties of PRIME is that it uses an OFDM approach instead of traditional single carrier solutions that have been used in the past for narrowband power line communications.

OFDM is well known in the literature and in industry. It is currently used in xDSL technologies, terrestrial wireless distribution of television signals (DVB-T, DVB-H and more), and has also been adapted for IEEE's high rate wireless LAN Standards (802.11a and 802.11g). In less than twenty years, in fact, OFDM has developed into a popular scheme for virtually all new telecoms standards: WiMAX, DAB, DRM, 3G cellular telephony, UWB, MoCA, Broadband over Power line...

OFDM is a digital multi-carrier modulation scheme, which uses a large number of closely-spaced orthogonal subcarriers to carry data. To obtain high spectral efficiency, these subcarriers typically overlap in frequency. However the mathematical property of orthogonality allows recovering each of the subcarriers separately at the receiver, so in practice subcarriers do not interfere with each other as would be the case with traditional FDM.



Each subcarrier is modulated with a conventional modulation scheme at a low symbol rate, maintaining data rates similar to conventional single-carrier modulation schemes in the same bandwidth. In practice, OFDM signals are efficiently generated and detected using the well-known Fast Fourier Transform (FFT) algorithm.

The primary advantage of OFDM over single-carrier schemes is its ability to cope with severe channel conditions -for example, attenuation of high frequencies in long power lines, narrowband interference and frequency-selective fading due to multipath- without complex additional mechanisms (e.g. equalization filters). Channel equalization is simplified because OFDM may be viewed as using many slowly-modulated narrowband signals rather than one rapidly-modulated wideband signal.

Additionally, low symbol rate makes the use of a guard interval (or cyclic prefix) between symbols affordable, rendering it possible to handle time-spreading and eliminate intersymbol interference (ISI). For low frequencies like the ones PRIME uses, multipath is not a critical issue so cyclic prefixes will not waste a significant part of OFDM symbols.

4.3 PRIME MAC layer overview

PRIME system is composed of sub networks, each of them defined in the context of a transformer station. A sub network is a tree with two types of nodes, the Base Node and the Services Nodes.

4.3.1 Base Node

The Base Node is at the root of the tree and acts as master node that provides connectivity to the sub network. It manages the sub network resources and connections.

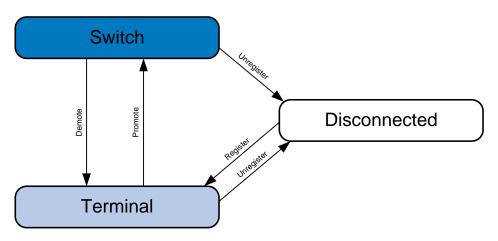
There is only one Base Node in a sub network. This Base Node is initially the sub network itself, and other nodes should follow a process of registering in order to enroll them to this sub network.

4.3.2 Service Node

Any other node of the sub network is a Service Node. Service Nodes are either leaves of the tree or branch points of the tree. These nodes start in a disconnected state and follow certain procedures to establish network connectivity. Each of these nodes is one point of the mesh of the sub network. These nodes have two responsibilities: connecting themselves to the sub network and switching the data of their neighbors in order to propagate connectivity.

Service Nodes change their behavior dynamically from "Terminal" functions to "Switch" functions and vice-versa. Changing of functional states occurs based on certain predefined events in the network.

Figure 4-3. Functional states of a Service Node





As shown in Figure 4-3, the three functional states of a Service Node are:

- **Disconnected**: Service Nodes start in a disconnected state. In this state a node is not capable of communicating or switching the traffic of another node. The primary function of a Service Node in this state is to search for an operational network in its proximity and to try to register itself to it.
- **Terminal**: In this state a Service Node is capable of communicating its traffic by establishing connections, but is not capable of switching the traffic of any other node.
- **Switch**: In this state a Service Node is capable of performing all Terminal functions. Additionally, it is capable of forwarding data to and from other devices in the sub network. It is a branch point in the tree.

The MAC layer provides all necessary features to manage PRIME networks and sub networks: addressing, synchronization (beacon management), dynamic management of the network structure (promotion and demotion of terminals), device registration management, connection setup and management, channel access arbitration, distribution of random sequences for deriving encryption keys, multicast group management...

4.4 Convergence Layer

The Convergence Layer (CL) classifies traffic associating it with its proper MAC connection. This layer performs the mapping of any kind of traffic to be properly included in MAC SDUs, providing access to the core MAC functionalities of system access, bandwidth allocation, connection management and mesh topology resolution. It may also include payload header suppression functions.

The convergence layer is separated into two sub layers:

- The Common Part Convergence Sub layer (CPCS) provides a set of generic services.
- The Service Specific Convergence Sub layer (SSCS) contains services that are specific to one application layer.

There are many SSCS, typically one per application, but only one common part, as shown in Figure 4-4:

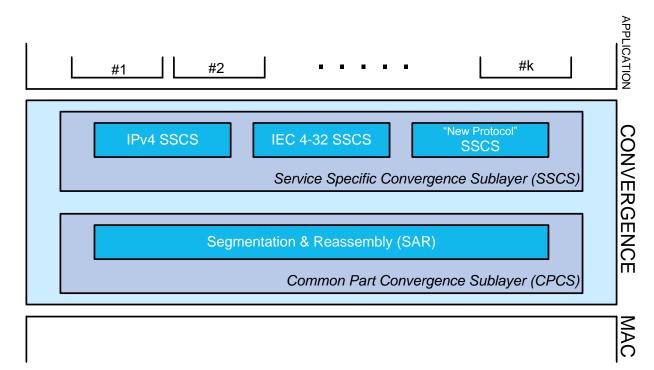


Figure 4-4. Convergence Layer



Several convergence sub layers are defined in order to accommodate different kinds of traffic into MAC SDUs, namely, IP convergence layer as a very useful and universal access to PRIME, and IEC 61334-4-32 as a link towards metering systems.





5. Processor and Architecture

5.1 ADD8051C3A Microcontroller Description

The following document provides a detailed description of the ADD8051C3A architecture and its hardware. ADD8051C3A is fully compatible with any 8051 legacy microcontroller, however there are some hardware differences that deserve to be taken into account.

CPU speed has been improved. The machine cycle has been reduced from 12 to 3 clock cycles, and all the instructions are executed in 1 or 2 machine cycles. Programs are executed a minimum of four times faster than in a standard-8051- architecture device, and up to a five times speedup is achieved for most programs.

The microcontroller is intended to work with an up to 256Kbytes external SRAM device. Program memory and XDATA memory share this external SRAM, where lower addresses are always dedicated to store program memory, since upper ones are configurable to allocate different XDATA memory size configurations depending on bank switching page size and SRAM total size.

The size of the extended program pages can be configured to 32K/16K/8K/0K.

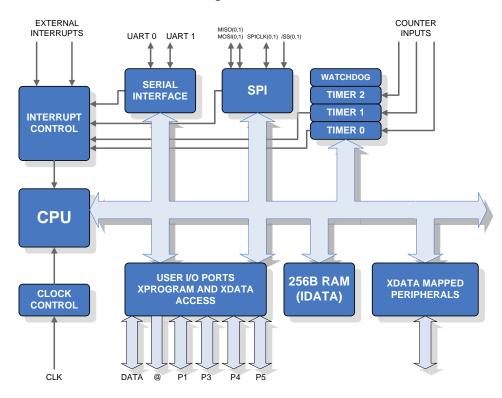
Furthermore, ATPL210 has developed a "compact" memory configuration using a 32Kbytes on-chip SRAM to satisfy applications that use a reduced program code version and in which saving external-circuitry costs and reducing PCB size are prioritary requirements. This configuration can be selected by means of EXTRAM external input:

- EXTRAM='0', 32Kbytes On-chip SRAM selected.
- EXTRAM='1', External memory device selected. On-chip SRAM is not accessible in this mode.

(Unless otherwise stated, External memory device size =256Kbytes is assumed in this user manual).

This document includes a detailed description of registers and peripherals. Some of these items are upgraded versions of the original 8051 microcontroller circuitry, while others are a simplified version.

Figure 5-1. ADD8051C3A microcontroller block diagram





5.2 Core Pinout Description

Figure 5-2. ADD8051C3A core pinout diagram

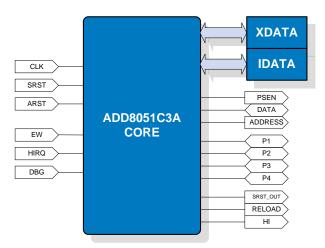


Table 5-1. Core pinout description

Pin	I/O	Description
CLK	1	Clock input.
ARST	1	Asynchronous reset.
SRST	I	SRST Must be held active one clock cycle to ensure proper power-on. External signal D_INIT is connected to this core input.
/EW	I	Watchdog enable Low level active. External signal/EWDG is connected to this core input.
HIRQ	1	Hard idle request High level active. When asserted, the microcontroller finishes the instruction in execution and goes to a special hard idle state. In hard idle state the CPU is stopped while peripheral circuits still run, CPU restarts only after HIRQ deassertion.
DBG	I	Debug mode enable High level active.
PSEN	0	Program Store Enable Read strobe to program memory.
DATA	I/O	Data bus 16-bit bidirectional port to access program and/or data memories.
ADDRESS	0	Address bus Output port to access program memory and XDATA.
SRST_OUT	0	Watchdog timer reset signal.
RELOAD	0	Program reload control signal Software or watchdog activated. Force program reload from a serial storage device to a parallel SRAM.
н	0	Hard idle flag, This flag is set to '1' when the CPU is in hard idle state.



PORTS: P1, P3, P4, P5 (*)	I/O	 I/O Ports. 8-bit pseudo bidirectional I/O ports with pull-up resistors. These ports are by default in pseudo-bidirectional configuration. Configured in this way, port pins that have 1s written to them are pulled high by the internal pull-ups, and in that state can be used as inputs. When a bit in a Port register has a 0 to 1 transition the related pin is driven high using transistor during 1 clock cycle, then the transistor is switched off and the pull-up resistor keeps the logic level. Ports P3, P4 and P5 can be also configured as push-pull mode ports. In push-pull mode the pin is always in output mode and logic 1 is always high driven (see 5.7.2). Some pins of P3, P4 and P5 also serve the functions of various special features of the microcontroller: P3.0 → RxD (serial port 0 input) P3.1 → TxD (serial port 0 output) P3.2 → INT0 (external interrupt 0) P3.3 → INT1 (external interrupt 1) P3.4 → T0 (timer 0 external input) P3.5 → T1 (timer 1 external input) P3.6 → WR(external data memory write strobe) P3.7 → RD (serial port 2 output) P4.1 → TxD (serial port 2 output) P4.2 → SS1 (slave select input) P4.3 → SPICLK1 (SPI1 clock input/output) P4.4 → MOS11 (master out / slave in data) P4.5 → MISO1 (master in / slave out data P4.6 → T2 (timer 2 external input) P5.0 → SS0 (slave select input) P5.1 → SPICLK0 (SPI0 clock input/output) P5.1 → SPICLK0 (SPI0 clock input/output) P5.3 → MISO0 (master out / slave in data) P5.5 → TxD (serial port 1 output)
		(*) each bit of ports 1, 3, 4 and 5 is composed of three lines: data output, data input and output enable (low level active)

5.3 Memory Organization

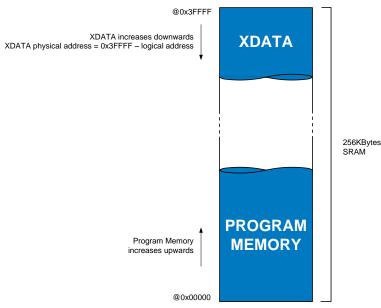
ADD8051C3A has separate logical address spaces for program and data memory, as shown in Figure 5-3. The logical separation of program and data memory allows the data memory to be accessed by 8-bit addresses, which can be quickly stored and manipulated by an 8-bit CPU. Nevertheless, 16-bit data memory addresses can also be generated through the DPTR (Data Pointer Registers).

Program memory must only be read. Writing in Program Memory could lead to unexpected results and must be avoided. There is a block of 64kbytes of base program memory, which can be increased by extended program pages.

XDATA occupies a separate address space from Program Memory. In the ADD8051C3A, the lowest 256 bytes of data memory (IDATA) are on chip. From the microcontroller point of view, up to 64Kbytes of XDATA can be addressed in the Data Memory space (ATPL210A allows to address more than 64Kbytes in data memory, see 11). The CPU generates the internal read and write signals, RD and WR, needed during XDATA Memory accesses. In ATPL210A, Program Memory and XDATA Memory are combined in a single SRAM and accessed by the same bus.



Figure 5-3. 256KB SRAM, XDATA and Program Memory



The program must be uploaded to the SRAM after power up. The source of program must be a non volatile storage unit (i.e. serial flash). A specific module (boot loader peripheral) has been designed to control the booting of the system after power up. The boot loader module can also be used to program the non volatile storage unit (serial flash) using a serial link. The non volatile storage unit must be connected to P5 port pins used for the SPI0 link (see Table 5-1). The serial link to the boot loader shares pins with the microcontroller Serial Port 0. For more information see 10.

5.3.2 Program Memory

After reset, the CPU begins execution from location 0000H and stack pointer (SP) value is set to 07H.

Interrupt Handling

A fixed location in Program Memory is assigned to each interrupt. The interrupt causes the CPU to jump to that location, where it begins executing the service routine.

Example: External Interrupt 0, is assigned to location 0003H. If External Interrupt 0 is going to be used, its service routine must begin at location 0003H. If the interrupt is not going to be used, its service location is available as general purpose Program Memory.

The interrupt service locations are spaced at 8-byte intervals beginning at 0x0003 (i.e.0x0003 for External Interrupt0, 0x000B for Timer0, etc.).

If an interrupt service routine is short enough, it can reside entirely within that 8-byte interval. Longer service routines can use a jump instruction to skip over subsequent interrupt locations, if other interrupts are in use.

Program Memory addresses are always 17 bits wide, even though the actual amount of Program Memory used may be less than 64kbytes.

P0 and P2 ports are no longer used to access external memories, but Special Function Registers (SFR) P0 and P2 are still functional.

- P2 register is used to generate the high order address byte when executing MOVX A,@Ri or MOVX @Ri,A.
- P0 register is the program address control register and it is used to specify the size of extended program pages and to control bank switching



5.3.3 Extended Addressing

ATPL210A combines program code and data in a single SRAM.

ADDRESS to access program is generated as a function of P0(7:6), P0(5:0) and PC(15:0).

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
P0	SX1	SX0	P05	P04	P03	P02	P01	P00

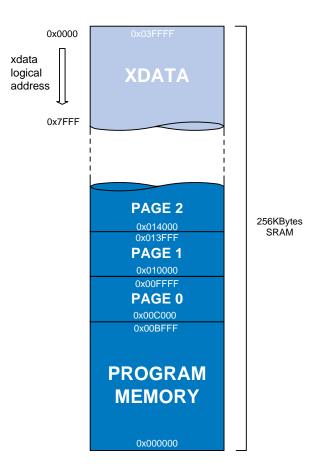
• SX(1:0): Size of extended program pages and common program space

SX(1:0)	Extended	Common
"11"	0KB	64KB
"10"	8KB	56KB
"01"	16KB	48KB
"00"	32KB	32KB

• P0(5:0): Extended program page number

Allowing different configurations as the one shown in Figure 5-4

Figure 5-4. SRAM Extended Addressing. 48KB common memory + 3x16KB extended program pages

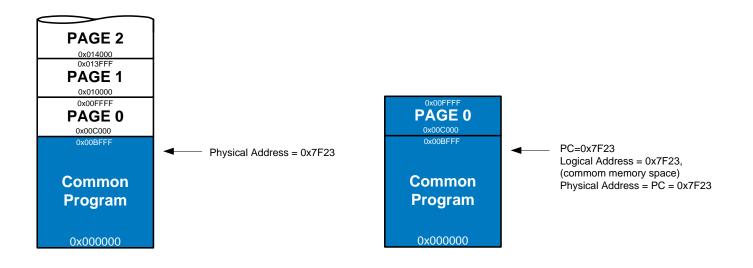




Common and extended memory sizes vary depending on the value in P0(7:6), as shown above. According to these size values, the core knows when the PC is pointing to an address located in common memory and when it is pointing to an address located in extended memory. When PC is pointing to an address in extended memory, then P0(5:0) indicates the extended page number and the physical address is automatically calculated. Some examples are shown below.

Example: SX(1:0)="00" (common memory size = 32KB) PC=0x7F23 (PC pointing to an address in common memory space) P0(5:0)= don't care LOGIC ADDRESS=0x7F23, common memory space PHYSICAL ADDRESS=0x7F23 (common memory space → P0(5:0) is ignored)

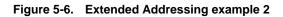
Figure 5-5. Extended Addressing example 1

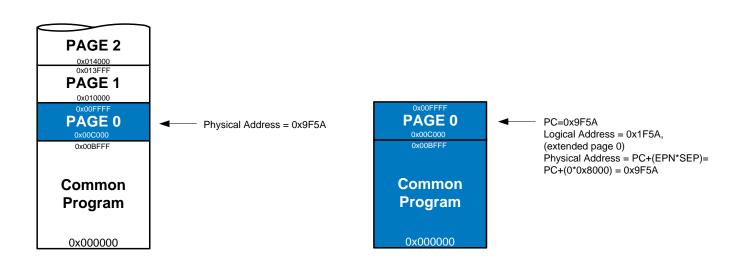






Example: SX(1:0)="00" (common memory size = 32KB) PC=0x9F5A (PC pointing to an address in extended memory) P0(5:0)= 0x00 (extended page 0) LOGIC ADDRESS=0x1F5A, extended page 0 PHYSICAL ADDRESS=0x9F5A

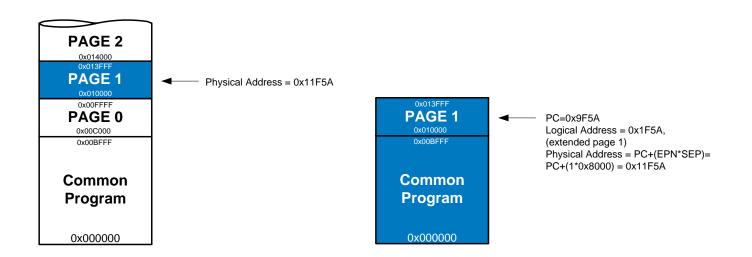






Example: SX(1:0)="00" (common memory size = 32KB) PC=0x9F5A (PC pointing to an address in extended memory) P0(5:0)= 0x01 (extended page 1) LOGIC ADDRESS=0x1F5A, extended page 1 PHYSICAL ADDRESS=0x11F5A





- Thus
 if PC < SCS</th>
 Physical Address = PC

 else
 Physical Address = PC+(EPN*SEP)
- where SCS= Size of Common Space; EPN=Extended Page Number; SEP= Size of Extended Pages(defined by SX(1:0) value).



5.3.4 Data Memory

Figure 5-8 shows the internal and external Logical Data Memory spaces available to the microcontroller user.

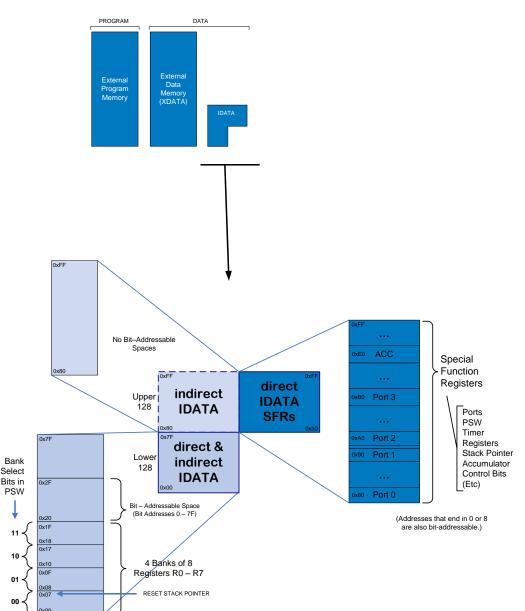


Figure 5-8. Data Memory Space

Logical addressing of the Internal Data Memory is mapped as shown in Figure 5-8. The internal IDATA memory space has a total size of 384 bytes and is divided into three blocks, which are generally referred to as the Lower 128 bytes, the Upper 128 bytes and SFR space (128 bytes size).

Internal Data Memory addresses are always one byte wide, which implies an address space of only 256 bytes. A simple trick is used to accommodate the 384 bytes of IDATA using 8 bit addresses: direct addresses higher than 7FH access SFR memory space, whereas indirect addresses higher than 7FH access the Upper 128 bytes of IDATA. Thus, Figure 5-8 shows the Upper 128 bytes and SFR space occupying the same block of addresses, 80H through FFH, although they are physically separate entities.



In the Lower 128 bytes of IDATA, the lowest 32 bytes are grouped into 4 banks of 8 registers. Program instructions call out these registers as R0 through R7. Two bits in the Program Status Word (PSW) select which register bank is in use. This allows more efficient use of code space, since register instructions are shorter than instructions that use direct addressing.

The next 16 bytes above the register banks form a block of bit-addressable memory space. The 80C51 instruction set includes a wide selection of single-bit instructions, and the 128 bits in this area can be directly addressed by these instructions. The bit addresses in this area are 00H through 7FH.

Summarizing, the Lower 128 bytes are accessed independently of the addressing mode (direct or indirect addressing). The Upper 128 bytes can be accessed only by indirect addressing, whereas SRF space is accessed only by direct addressing.

Table 5-2 gives a look at the Special Function Register (SFR) space. SFRs include the Port registers, timers, peripheral controls, etc. These registers can only be accessed by direct addressing. Sixteen addresses in SFR space are both byte and bit addressable. The bit addressable SFRs are those whose address ends in 0H or 8H. ADD8051C3A includes the same 21 SFRs included in the legacy 8051C plus additional SFRs to implement the extended features.

The CPU generates the RD (P3.7) and WR (P3.6) signals needed during SRAM accesses, as shown below.

5.3.5 SFR Registers

Table 5-2 shows a map of the internal memory area called the Special Function Register (SFR) space. The entries in bold are registers not included in the standard architecture of the 8051, these entries are specific for the ADD8051C3A. Blank entries are not implemented on the chip. User software should not write to these unimplemented locations. Read accesses to these addresses will return random data.

Table 5-2. Special Function Registers (specific ADD8051C3A registers in bold)

F8								Т3	FF
FO	В	CONF							F7
E8									EF
EO	ACC	SPSTAT	SPCTL	SPDAT	SPSTAT1	SPCTL1			E7
D8	SPDAT10	SPDAT11	SPDAT12	SPDAT13	SPDAT14	SPDAT15	SPDAT16	SPDAT17	DF
D0	PSW								D7
C8	T2CON	T2MOD	RCAP2L	RCAP2H	TL2	TH2			CF
C0	P4				P5				C7
B8	IP								BF
BO	P3							IPH	B7
A8	IE								AF
A0	P2	IE2	AUX1	AUX2					A7
98	SCON	SBUF	SCON1	SBUF1	P3M	P4M	P5M		9F
90	P1	IP2HL					T1NP		97
88	TCON	TMOD	TLO	TL1	TH0	TH1	TPR	T1NC	8F
80	PO	SP	DPL	DPH				PCON	87



The functions of some SFRs are described in the text below, while others are described with their related peripherals.

• Accumulator (address: E0H)

ACC is the Accumulator register. Note that mnemonics for accumulator specific instructions usually refer to the Accumulator simply as A.

• B Register (address: F0H)

The B register is used during multiply and divide operations. For other instructions it can be treated as another scratch pad register.

PSW register (address: D0H)

The PSW register contains program status information as detailed below.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
PSW	CY	AC	F0	RS1	RS0	OV	F1	Р

- CY: Carry flag
- AC: Auxiliary carry flag
- F0: General purpose flag
- RS(1:0): Register bank select control bits
- OV: Overflow flag
- F1: User definable flag
- P: Parity flag

• Stack Pointer (address: 81H)

8 bit stack pointer that is initialized to internal memory of 07H. The user program can initialize it at any internal RAM location ranging from 07H to FFH.

Auxiliary 1

The AUX1 register includes the data pointer selection bit, the software reset control and the switch to enable wake-up from power-down using external interrupts.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
AUX1			SRST	GF2	WUPD	0		DPS

- --: Reserved bits
- SRST: Software reset
- GF2: General purpose flag
- WUPD: When set, enables external interrupts driven wake-up from power down
- 0: fixed '0'
- DPS: Data pointer selector, selects between DPTR0 and DPTR1

• P0, P1, P2, P3, P4, P5 registers

P1, P3, P4 and P5 are the SFR registers of Ports 1, 3, 4 and 5 respectively.

Writing a one/zero to a bit in these registers causes the corresponding port output pin to switch high/low. When a port bit is used as an input the corresponding port SFR must be set to '1'.

Ports 3, 4 and 5 are pseudo bidirectional by default, and can be configured to push-pull or pseudo bidirectional using SFRs P3M, P4M and P5M respectively, as follows: when P4M(i) is set, P4(i) is configured as push-pull.

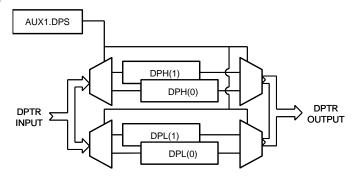


P0 and P2 ports are no longer used to address program and external data. P0 SFR takes control of extended addressing and program banking and P2 SFR is used when MOVX @Ri instructions are executed (in this case, P2 SFR content is used as most significant address byte).

Data Pointer

The Data Pointer Register (DPTR) consists of a high byte (DPH) and a low byte (DPL). Its intended function is to hold a 16-bit address. It may be manipulated as a 16-bit register or as two independent 8-bit registers. The same entry accesses two different data pointers (DPTR0, DPTR1), user software can switch between them using the data pointer selection bit in AUX1.

Figure 5-9. DPTR selection scheme



Serial Data Buffers

The Serial Buffers have actually two separate registers, a transmit buffer and a receive buffer. When data is moved to SBUF, it goes to the transmit buffer and is held for serial transmission. (Moving a byte to SBUF is what initiates the transmission.) When data is moved from SBUF, it comes from the receive buffer.

Timer Registers

Register pairs (TH0, TL0), (TH1, TL1) and (TH2, TL2) are the 16-bit Counting registers for Timer/Counters 0, 1 and 2, respectively.

T3 is the 8-bit counting register of the watchdog timer, see 6.3 for detailed behavior description.

CONF Register

CONF register includes specific configuration features for the flash loader and Standard Serial Interface.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
CONF						RLD	RLT3	BAUD2

- --: Reserved bits
- RLD: When set, it forces program reload from SPI flash
- RLT3: When set, it forces program reload from SPI flash when watchdog timeout occurs
- BAUD2: doubles USART baud rates when set

Control Registers

Special Function Registers IP, IPH, IP2HL, IE, IE2, TMOD, TCON, T2CON, SCONj, SPSTATj and PCON contain control and status bits for the interrupt system, the Timer/Counters, the serial ports and the serial peripheral interfaces



5.4 Instruction Set

The instruction set provides a variety of fast addressing modes for accessing the internal RAM to facilitate byte operations on small data structures. The instruction set provides extensive support for one-bit variables as a separate data type, allowing direct bit manipulation in control and logic systems that require boolean processing.

5.4.1 Program Status Word

The Program Status Word (PSW) contains several status bits that reflect the current state of the CPU. The PSW, shown in 5.3.5, resides in the SFR space. It contains the Carry bit, the Auxiliary Carry (for BCD operations), the two register bank select bits, the Overflow flag, a Parity bit, and two user-definable status flags.

The Carry bit, other than serving the function of a Carry bit in arithmetic operations, also serves as the "Accumulator" for a number in Boolean operations.

The bits RS0 and RS1 are used to select one of the four register banks shown in Figure 5-8. A number of instructions refer to these RAM locations as R0 through R7. The selection of which of the four is being referred to is made on the basis of the RS0 and RS1 at execution time.

The Parity bit reflects the number of 1s in the Accumulator: P = 1 if the Accumulator contains an odd number of 1s, and P = 0 if the Accumulator contains an even number of 1s. Thus the number of 1s in the Accumulator plus P is always even. Two bits in the PSW are uncommitted and may be used as general purpose status flags.

5.4.2 Addressing Modes

• Direct Addressing

In direct addressing the operand is specified by an 8-bit address field in the instruction. Only internal Data RAM and SFRs can be directly addressed.

Indirect Addressing

In indirect addressing the instruction specifies a register which contains the address of the operand. Both internal and external RAM can be indirectly addressed.

The address register for 8-bit addresses can be either R0 / R1 of the selected register bank or the Stack Pointer. The address register for 16-bit addresses can only be the 16-bit "data pointer" register, DPTR.

Register Instructions

The register banks, containing registers R0 through R7, can be accessed by certain instructions which carry a 3-bit register specification within the opcode of the instruction. Instructions that access the registers this way are code efficient, since this mode eliminates an address byte. When the instruction is executed, one of the eight registers in the selected bank is accessed. One of four banks is selected at execution time by the two bank select bits in the PSW.

Register-Specific Instructions

Some instructions are specific to a certain register. For example, some instructions always operate on the Accumulator, or Data Pointer, etc., so no address byte is needed to point to it. The opcode itself does that. Instructions that refer to the Accumulator as A assemble as accumulator specific opcodes.

Immediate Constants

The value of a constant can follow the opcode in Program Memory.

For example:

MOV A, 64H

loads the Accumulator with the number 64H.

Indexed Addressing

Only program Memory can be accessed with indexed addressing, and it can only be read. This addressing mode is intended for reading look-up tables in Program Memory. A 16-bit base register (either DPTR or the Program Counter) points to the base of the table, and the Accumulator is set up with the table entry number.



The address of the table entry in Program Memory is formed by adding the Accumulator data to the base pointer.

Another type of indexed addressing is used in the "case jump" instruction. In this case the destination address of a jump instruction is computed as the sum of the base pointer and the Accumulator data.

5.4.3 Arithmetic Instructions

The menu of arithmetic instructions is listed in Table 5-3, which indicates the addressing modes that can be used with each instruction to access the <byte> operand. For example, the ADD A,<byte> instruction can be written as:

ADD	A, 6EH	(direct addressing)	
ADD	A, @R1	(indirect addressing)	
ADD	A, R6	(register addressing)	
ADD	A, #206(immediate constant)		
	, (,	

The execution times listed in Table 5-3 assume 3 clock cycles per machine cycle (mc), using a 10MHz clock 1 mc is executed in 0.3us. All the arithmetic instructions execute in 1mc except the Multiply and Divide instructions, which take 2mc.

Note that any byte in the internal Data Memory space can be incremented without going through the Accumulator.

One of the INC instructions operates on the 16-bit Data Pointer. The Data Pointer is used to generate 16-bit addresses for external memory, so being able to increment it in one 16-bit operation is a useful feature.

The MUL AB instruction multiplies the Accumulator by the data in the B register and puts the 16-bit product into the concatenated B and Accumulator registers.

The DIV AB instruction divides the Accumulator by the data in the B register and leaves the 8-bit quotient in the Accumulator, and the 8-bit remainder in the B register. Oddly enough, DIV AB finds less use in arithmetic "divide" routines than in radix conversions and programmable shift operations. In shift operations, dividing a number by 2ⁿ shifts its n bits to the right. Using DIV AB to perform the division completes the shift in 2mc and leaves the B register holding the bits that were shifted out.

The DA A instruction is for BCD arithmetic operations. In BCD arithmetic, ADD and ADDC instructions should always be followed by a DA A operation, to ensure that the result is also in BCD. Note that DA A will not convert a binary number to BCD. The DA A operation produces a meaningful result only as the second step in the addition of two BCD bytes.





					\$	EXECUTION	
MNEMONIC	OPERATION	DIR	IND	REG	ІММ	TIMES (mc)	
ADD A, <byte></byte>	A = A + <byte></byte>	Х	х	Х	Х	1	
ADD A, <byte></byte>	A = A + <byte> + C</byte>	Х	х	Х	Х	1	
SUBB A, <byte></byte>	$A = A - \langle byte \rangle - C$	Х	Х	Х	Х	1	
INC A	A = A + 1		Accumul	ator only		1	
INC <byte></byte>	<byte> = <byte> + 1</byte></byte>	Х	Х	Х		1	
INC DPTR	DPTR = DPTR + 1	Data Pointer only			1		
DEC A	A = A - 1		Accumul	ator only		1	
DEC <byte></byte>	<byte> = <byte> - 1</byte></byte>	Х	Х	Х		1	
MUL AB	B*A = B x A	ACC and B only		2			
	A = Int[A/B]			0			
DIV AB	B = Mod[A/B]	ACC and B only 2			2		
DA A	Decimal Adjust		Accumul	ator only		1	

5.4.4 Logical Instructions

Table 5-4 shows the list of logical instructions. The instructions that perform Boolean operations (AND, OR, Exclusive OR, CPL) on bytes perform the operation on a bit-by-bit basis. That is, if the Accumulator contains 00110101B and byte contains 01010011B, then:

ANL A, <byte>

will leave the Accumulator holding 00010001B.

The addressing modes that can be used to access the <byte> operand are listed in Table 5-4.

The ANL A, <byte> instruction may take any of the forms:

ANL	A,6FH	(direct addressing)
ANL	A,@R0	(indirect addressing)
ANL	A,R5	(register addressing)
ANL	A,#42H	(immediate constant)

All of the logical instructions execute in 1mc.



Table 5-4. Logical Instructions

MNEMONIC	OPERATION		ADDRESSING MODES		S	EXECUTION
MINEMICINIC	OFERATION	DIR	IND	REG	ІММ	TIME (mc)
ANL A, <byte></byte>	A = A.AND. <byte></byte>	х	Х	x	Х	1
ANL <byte>,A</byte>	<byte> = <byte> .AND.A</byte></byte>	x				1
ANL <byte>,#data</byte>	<byte> = <byte> .AND.#data</byte></byte>	х				1
ORL A, <byte></byte>	A = A.OR. <byte></byte>	х	х	X	Х	1
ORL <byte>,A</byte>	<byte> = <byte> .OR.A</byte></byte>	х				1
ORL <byte>,#data</byte>	<byte> = <byte> .OR.#data</byte></byte>	х				1
XRL A, <byte></byte>	A = A.XOR. <byte></byte>	х	х	Х	Х	1
XRL <byte>,A</byte>	<byte> = <byte> .XOR.A</byte></byte>	х				1
XRL <byte>,#data</byte>	<byte> = <byte> .XOR.#data</byte></byte>	х				1
CLR A	A = 00H		Accum	ulator only		1
CPL A	A = .NOT.A		Accum	ulator only		1
RL A	Rotate ACC Left 1 bit	Accumulator only		1		
RLC A	Rotate Left through Carry	Accumulator only		1		
RR A	Rotate ACC Right 1 bit	Accumulator only		1		
RRC A	Rotate Right through Carry		Accum	ulator only		1
SWAP A	Swap Nibbles in A		Accum	ulator only		1

Note that Boolean operations can be performed on any byte in the internal Data Memory space without going through the Accumulator. The *XRL <byte>, #data instruction*, for example, offers a quick and easy way to invert port bits, as in *XRL P1, #0FFH*.

If the operation is in response to an interrupt, not using the Accumulator saves time and effort to push it onto the stack in the service routine.

The Rotate instructions (RL, A, RLC A, etc.) shift the Accumulator 1 bit to the left or right. For a left rotation, the MSB rolls into the LSB position. For a right rotation, the LSB rolls into the MSB position.

The SWAP A instruction interchanges the high and low nibbles within the Accumulator. This is a useful operation in BCD manipulations.

5.4.5 Data Transfer Instructions

Internal RAM

Table 5-5 shows the menu of instructions that are available for moving data around within the internal memory spaces, and the addressing modes that can be used with each one. All of these instructions are executed in 1 mc.



MNEMONIC	OPERATION	ADDRESSING MODES				EXE CUTI ON
		DIR	IND	REG	імм	TIME (mc)
MOV A, <src></src>	A = <src></src>	х	x	х	Х	1
MOV <dest>,A</dest>	<dest> = A</dest>	х	x	x		1
MOV <dest>,<src></src></dest>	<dest> = <src></src></dest>	x	x	x	Х	1
MOV DPTR,#data16	DPTR = 16-bit immediate constant				Х	1
PUSH <src></src>	INC SP:MOV"@SP", <src></src>	Х				1
POP <dest></dest>	MOV <dest>,"@SP":DEC SP</dest>	x				1
XCH A, <byte></byte>	ACC and <byte> exchange data</byte>	х	х	х		1
XCHD A,@Ri	ACC and @Ri exchange low nibbles		x			1

Table 5-5. Data Transfer Instructions that Access Internal Data Memory Space

The MOV <dest>, <src> instruction allows data to be transferred between any two internal RAM or SFR locations without going through the Accumulator. Remember, the upper 128 bytes of IDATA can be accessed only by indirect addressing, and SFR space only by direct addressing.

Note that in any 80C51 device, the stack resides in on-chip RAM, and grows upwards. The PUSH instruction first increments the Stack Pointer (SP), then copies the byte into the stack. PUSH and POP use only direct addressing to identify the byte being saved or restored, but the stack itself is accessed by indirect addressing using the SP register. This means the stack can go into the Upper 128 bytes of RAM, if they are implemented, but not into SFR space.

The Data Transfer instructions include a 16-bit MOV that can be used to initialize the Data Pointer (DPTR) for look-up tables in Program Memory, or for 16-bit external Data Memory accesses.

The XCH A, <byte> instruction causes the Accumulator and addressed byte to exchange data. The XCHD A, @Ri instruction is similar, but only the low nibbles are involved in the exchange.

External RAM (from microcontroller point of view)

Table 5-6 shows a list of the Data Transfer instructions that access external Data Memory. Only indirect addressing can be used. The choice is whether to use a one-byte address, @Ri, where Ri can be either R0 or R1 of the selected register bank, or a two-byte address, @DPTR. All of these instructions execute in 2 mc.

Note that in all external Data RAM accesses, the Accumulator is always either the destination or source of the data.

The read and write strobes to external RAM are activated only during the execution of a MOVX instruction.



ADDRESS WIDTH	MNEMONIC	OPERATION	EXECUTION TIME (mc)
8 bits	MOVX A,@Ri	Read external RAM @Ri	2
8 bits	MOVX @Ri,A	Write external RAM @ Ri	2
16 bits	MOVX A,@DPTR	Read external RAM @ DPTR	2
16 bits	MOVX @DPTR,A	Write external RAM @ DPTR	2

Table 5-6. Data Transfer Instructions that access External Data Memory Space

Code Constants

Table 5-7 shows the two instructions that are available for reading code constants in Program Memory. Since these instructions access only Program Memory, the code constants can only be read, not updated.

The mnemonic is MOVC for "move constant". The first MOVC instruction in Table 5-7 can accommodate a table of up to 256 entries numbered 0 through 255. The number of the desired entry is loaded into the Accumulator, and the Data Pointer is set up to point to the beginning of the table. Then, *MOVC A*, *@A+DPTR* copies the desired table entry into the Accumulator.

The other MOVC instruction works the same way, excepting that the Program Counter (PC) is used as the base address: *MOVC A*, @A+PC

Table 5-7. Code Constants Read Instructions

MNEMONIC	OPERATION	EXECUTION TIME (mc)
MOVC A,@A+DPTR	Read program memory at (A + DPTR)	2
MOVC A,@A+PC	Read program memory at (A + PC)	2

5.4.6 Boolean Instructions

80C51 devices contain a complete Boolean (single-bit) processor. The internal RAM contains 128 addressable bits, and the SFR space can support up to 128 addressable bits as well. All of the port lines are bit-addressable, and each one can be treated as a separate single-bit port. The instructions that access these bits are not just conditional branches, but a complete menu of move, set, clear, complement, OR, and AND instructions. These kinds of bit operations are not easily obtained in other architectures with any amount of byte-oriented software. The instruction set for the Boolean processor is shown in Table 5-8. All bit accesses are by direct addressing.

Bit addresses 00H through 7FH are in the Lower 128, and bit addresses 80H through FFH are in SFR space.

The Carry bit in the PSW is used as the single-bit Accumulator of the Boolean processor. Bit instructions that refer to the Carry bit as C assemble as Carry-specific instructions (CLR C, etc.). The Carry bit also has a direct address, since it resides in the PSW register, which is bit-addressable.

Note that the Boolean instruction set includes ANL and ORL operations, but not the XRL (Exclusive OR) operation.

There is a series of bit-test instructions which execute a jump if the addressed bit is set (JC, JB, JBC) or if the addressed bit is not set (JNC, JNB). JBC executes the jump if the addressed bit is set, and also clears the bit. Thus a flag can be tested and cleared in one operation. All the PSW bits are directly addressable, so the Parity bit, or the general purpose flags, for example, are also available to the bit-test instructions.



Table 5-8. Boolean Instructions

MNEMONIC	OPERATION	EXECUTION TIME (mc)
ANL C,bit	C = C.AND.bit	1
ANL C,/bit	C = C.ANDNOT.bit	1
ORL C,bit	C = C.OR.bit	1
ORL C,/bit	C = C.ORNOT.bit	1
MOV C,bit	C = bit	1
MOV bit,C	bit = C	1
CLR C	C = 0	1
CLR bit	bit = 0	1
SETB C	C = 1	1
SETB bit	bit = 1	1
CPL C	C = .NOT.C	1
CPL bit	bit = .NOT.bit	1
JC rel	Jump if C = 1	2
JNC rel	Jump if C = 0	2
JB bit,rel	Jump if bit = 1	2
JNB bit,rel	Jump if bit = 0	2
JBC bit,rel	Jump if bit = 1; CLR bit	2

Relative Offset

The destination address for these jumps is specified to the assembler by a label or by an actual address in Program memory. However, the destination address assembles to a relative offset byte. This is a signed (two's complement) offset byte which is added to the PC in two's complement arithmetic if the jump is executed. The range of the jump is therefore -128 to +127 Program Memory bytes relative to the first byte following the instruction.

5.4.7 Jump Instructions

Table 5-9 shows the list of unconditional jumps and the execution time associated.

The table lists SJMP, LJMP, and AJMP, which differ in the format of the destination address. JMP is a generic mnemonic which can be used if the programmer does not care which way the jump is encoded.

The SJMP instruction encodes the destination address as a relative offset, as described above. The instruction is 2 bytes long, consisting of the opcode and the relative offset byte. The jump distance is limited to a range of -128 to +127 bytes relative to the instruction following the SJMP.

The LJMP instruction encodes the destination address as a 16-bit constant. The instruction is 3 bytes long, consisting of the opcode and two address bytes. The destination address can be anywhere in the 64k Program Memory space.



The AJMP instruction encodes the destination address as an 11-bit constant. The instruction is 2 bytes long, consisting of the opcode, which itself contains 3 of the 11 address bits, followed by another byte containing the low 8 bits of the destination address. When the instruction is executed, these 11 bits replace the low 11 bits in the PC. The high 5 bits stay the same. Hence the destination has to be within the same 2k block as the instruction following the AJMP.

The JMP @A+DPTR instruction supports case jumps. The destination address is computed at execution time as the sum of the 16-bit DPTR register and the Accumulator. Typically, DPTR is set up with the address of a jump table.

Table 5-9 shows two "CALL addr" instructions LCALL and ACALL, which differ in the format in which the subroutine address is given to the CPU. CALL is a generic mnemonic which can be used if the programmer does not care which way the address is encoded.

The LCALL instruction uses the 16-bit address format, and the subroutine can be anywhere in the 64k Program Memory space.

The ACALL instruction uses the 11-bit format, and the subroutine must be in the same 2k block as the instruction following the ACALL.

Subroutines should end with a RET instruction, which returns execution to the instruction following the CALL.

RETI is used to return from an interrupt service routine. The only difference between RET and RETI is that RETI tells the interrupt control system that the interrupt in progress is done. If there is no interrupt in progress at the time RETI is executed, then the RETI is functionally identical to RET.

Table 5-10 shows the list of conditional jumps available to the microcontroller user. All of these jumps specify the destination address by the relative offset method, and so are limited to a jump distance of -128 to +127 bytes from the instruction following the conditional jump instruction.

There is no Zero bit in the PSW. The JZ and JNZ instructions test the Accumulator data for that condition.

The DJNZ instruction (Decrement and Jump if Not Zero) is for loop control.

The CJNE instruction (Compare and Jump if Not Equal) can also be used for loop control. Two bytes are specified in the operand field of the instruction. The jump is executed only if the two bytes are not equal. Another application of this instruction is in "greater than, less than" comparisons. The two bytes in the operand field are taken as unsigned integers. If the first is less than the second, then the Carry bit is set (1). If the first is greater than or equal to the second, then the Carry bit is cleared.

MNEMONIC	OPERATION	EXECUTION TIME (mc)
(S)JMP addr	Jump to addr	2
(L)JMP addr	Jump to addr	2
(A)JMP addr	Jump to addr	2
JMP @A+DPTR	Jump to A + DPTR	2
(A,L)CALL addr	Call subroutine at addr	2
RET	Return from subroutine	2
RETI	Return from interrupt	2
NOP	No operation	1

Table 5-9. Unconditional Jump instructions



Table 5-10. Conditional Jump instructions

MNEMONIC	OPERATION	ADDRESSING MODES				EXECUTION
		DIR	IND	REG	ІММ	TIME(mc)
JZ rel	Jump if A = 0					2
JNZ rel	Jump if A /= 0			2		
DJNZ <byte>,rel</byte>	Decrement and jump if not zero	х		Х		2
CJNE A, <byte>,rel</byte>	Jump if A /= <byte></byte>	х			х	2
CJNE <byte>,#data,rel</byte>	Jump if <byte> /= #data</byte>		Х	Х		2

5.5 CPU Timing

Execution time is divided into machine cycles. A machine cycle consists of a sequence of 3 states, numbered S1 through S3. Each state time lasts for one clock period. Thus a machine cycle takes 3 clock periods (or near 0.15µs if the clock frequency is 20MHz).

Figure 5-10 shows fetch/execute sequences for various kinds of instructions; while an instruction is being executed code bytes for the next instruction are fetched. Normally three program fetches are generated during each machine cycle, even if the next instruction to be executed doesn't require it. If the instruction doesn't need the three code bytes, the CPU simply ignores the extra fetches, and the Program Fetch Counter is not incremented.

All the instructions, except MOVX, have the same fetch/execute sequence.

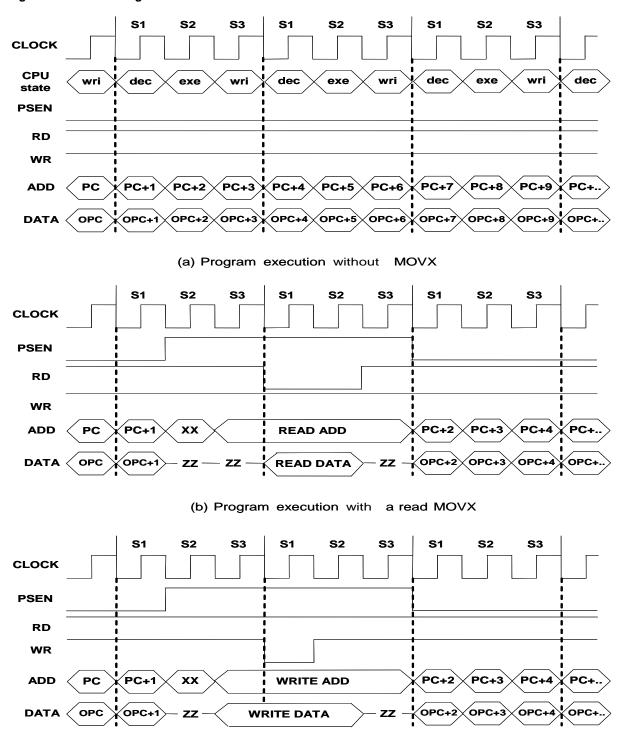
The MOVX instructions take two machine cycles to execute. During a MOVX instruction only one fetch is generated, S1 in machine cycle 1. This is the only time program fetches are skipped. The fetch/execute sequence for MOVX instructions is shown in Figure 5-10b and Figure 5-10c.

The fetch/execute sequences are the same whether the Program Memory is internal or external to the chip. Execution times do not depend on whether the Program Memory is internal or external.

Figure 5-10 shows the signals and timing involved in program fetches when the Program Memory is external. If Program Memory is external, then the Program Memory read strobe PSEN is normally activated ('0') three times per machine cycle, as shown in Figure 5-10a. If an access to external Data Memory occurs, as shown in Figure 5-10b and 15c, five PSENs are skipped, because the address and data bus are being used for the Data Memory access.



Figure 5-10. Bus timing in 8051C3A



(c) Program execution with a write MOVX



5.5.2 Reset Modes

ATPL210A can be reset in different ways, listed and explained below

- External reset. The external reset input signal is RSTA which is an asynchronous reset. Every time RSTA signal is asserted, the boot loader is triggered (thus, microcontroller is halted, target program from the serial flash is uploaded to the SRAM, and then microcontroller is released to begin program execution). In power-on, D_INIT must be set before RSTA in order to ensure proper system start up.
- Software reset → AUX1(5) = AUX1.SRST. This field in AUX1 SFR forces the microcontroller to start program execution from PC=0, and all the SFRs are initialized to their reset default values. This reset doesn't affect the boot loader reload.
- **Software reset** → CONF(2) = CONF.RLD. This field is directly connected to boot loader reload input. Every time the microcontroller sets this field, the boot loader is triggered.
- Watchdog reset → /EWDG signal and CONF(1)=CONF.RLT3.
 - CONF.RLT3 = '0' : After a watchdog timeout, microcontroller starts execution from PC=0, but boot loader is not triggered.
 - CONF.RLT3 = '1' : After a watchdog timeout, if field RLT3 in CONF SFR is '1', boot loader is triggered.

The microcontroller keeps track of the start-up point in the AUX2 register

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
AUX2						ST2	ST1	ST0

- --: Reserved bits
- ST(2:0): Start-up flags
 - "000":Asynchronous external reset
 - "001":Software reset by writing AUX1(5)='1'
 - "010":Watchdog T3 timeout
 - "100":Asynchronous external reset

SFR values after a reset are shown in Table 5-11. The internal RAM is not affected. Port registers are initialized to FFH to keep port pins in pull-up state.



Table 5-11. SFRs Reset values

SFR	Address	Reset Value
P0	80H	11111111
SP	81H	00000111
DPL	82H	00000000
DPH	83H	00000000
PCON	87H	00000000
TCON	88H	00000000
TMOD	89H	00000000
TL0	8AH	00000000
TL1	8BH	00000000
TH0	8CH	00000000
TH1	8DH	00000000
TPR	8EH	00111111
T1NC	8FH	00000000
P1	90H	11111111
IP2HL	91H	xxxx0000
T1NP	96H	00000000
SCON	98H	00000000
SBUF	99H	00000000
SCON1	9AH	00000000

SFR	Address	Reset Value
SBUF1	9BH	00000000
P3M	9CH	00000000
P4M	9DH	00000000
P5M	9EH	00000000
P2	A0H	11111111
IE2	A1H	xxxxxx00
AUX1	A2H	00000000
AUX2	A3H	xxxxxbbb
IE	A8H	00000000
P3	BOH	11111111
IPH	B7H	00000000
IP	B8H	00000000
P4	COH	11111111
P5	C4H	11111111
T2CON	C8H	00000000
T2MOD	C9H	xx001100
RCAP2L	CAH	00000000
RCAP2H	CBH	00000000
TL2	ССН	00000000

SFR	Address	Reset Value
TH2	CDH	00000000
PSW	D0H	00000000
SPDAT10	D8H	00000000
SPDAT11	D9H	0000000
SPDAT12	DAH	0000000
SPDAT13	DBH	00000000
SPDAT14	DCH	00000000
SPDAT15	DDH	00000000
SPDAT16	DEH	00000000
SPDAT17	DFH	0000000
ACC	EOH	00000000
SPSTAT	E1H	000xxxxx
SPCTL	E2H	00000100
SPDAT	E3H	00000000
SPSTAT1	E4H	000xx000
SPCTL1	E5H	00000100
В	FOH	00000000
CONF	F1H	00000000
Т3	FFH	00000000





5.5.3 Power Saving Modes

ADD8051C3A has two power-saving modes, Idle and Power Down.

In the Idle mode (IDL = 1), the CPU is stopped but the peripheral circuits continue running.

In Power Down (PD = 1), the CPU and the peripheral circuits are stopped, except external interrupts hardware that can be configured to run when in power down mode.

The Idle and Power Down Modes are activated by setting bits in Special Function Register PCON. If 1s are written to PD and IDL at the same time, PD takes precedence. When watchdog is enabled power down modes are automatically disabled.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
PCON	SMOD			WLE	GF1	GF0	PD	IDL

- SMOD: Double baud rate bit when timer 1 is used as time generator and serial port is in modes 1, 2 or 3
- --: Reserved bits
- WLE: Watchdog load enable. It must be set by software to enable T3 reload. It is reset by hardware after 13 machine cycles
- GF1: General purpose flag bit, user programmable
- GF0: General purpose flag bit, user programmable
- PD: Sets power-down mode
- IDL: Sets idle mode

5.5.4 Idle Mode

Setting PCON(0), CPU is stopped but peripheral circuits (Interrupts, Timers, Standard Serial Interface and Serial Peripheral Interface functions) continue running. The CPU status is entirely preserved; Stack Pointer, Program Counter, Program Status Word, Accumulator, and all other registers maintain their data during Idle. The port pins hold the logical states they had at the time Idle was activated. PSEN holds at logic high level.

There are two ways to terminate the Idle mode.

- Activation of any enabled interrupt will cause PCON(0) to be cleared by hardware, terminating the Idle mode. The interrupt will be served, and following RETI, the next instruction to be executed will be the one following the instruction that put the device into Idle.
- Reset signal. The signal at the RSTA pin redefines all the SFR values to their reset values and then the IDL bit is cleared. At this time, and with three clock cycle delay, the CPU restart program execution with program counter at initial value 0x0000, RAM content remains unchanged.

5.5.5 Power-Down Mode

An instruction that sets PCON(1) causes that to be the last instruction executed before going into the Power Down mode. In the Power Down mode all functions are stopped, except external interrupts hardware. The contents of the onchip RAM and Special Function Registers are maintained. The port pins output the values held by their respective SFRs. The PSEN output is held low.

There are two ways to exit Power Down mode.

- Hardware reset. Hardware reset redefines all the SFRs.
- External interrupt. External interrupts must be configured for low level activation. When they are enabled a low level at their inputs will cause a microcontroller restart and a jump to the corresponding interruption routine. As in idle mode, the interrupt will be served, and following RETI, the next instruction to be executed will be the one following the instruction that put the device into power down.



5.6 Interrupts

The microcontroller has 5 user interrupt sources:

- 1 timer interrupt (Timer2).
- 2 serial port interrupts (Standard Serial 0 & Standard Serial 1).
- 2 serial peripheral interface interrupts (SPI0 & SPI1).

And 4 internal system interrupt sources:

- 2 timer interrupts (Timer0 & Timer1)
- 2 generic interrupts (INT0 & INT1)

Interrupt priorities are user selectable to a priority level ranging from "000" to "011" (higher value correspond to higher priority level).

5.6.1 Interrupt Enabling

Each interrupt source can be individually enabled or disabled by setting or clearing a bit in the SFRs named IE and IE2 (Interrupt Enable). The IE register also contains a global disable bit, which can be cleared to disable all interrupts at once.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
IE	EA	ET2	ESPI	ES	ET1	EX1	ET0	EX0

- EA: '1' enables all interrupts; '0' disables all interrupts
- ET2: '1' enables timer 2 interrupt; '0' disables timer 2 interrupt
- ESPI: '1' enables serial peripheral interface0 interrupt; '0' disables serial peripheral interface0 interrupt
- ES: '1' enables serial port 0 interrupt; '0' disables serial port0 interrupt
- ET1: '1' enables timer 1 interrupt; '0' disables timer 1 interrupt
- EX1: '1' enables external interrupt 1; '0' disables external interrupt 1
- ET0: '1' enables timer 0 interrupt; '0' disables timer 0 interrupt
- EX0: '1' enables external interrupt 0; '0' disables external interrupt 0

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
IE2					ESPI1			ES1

- --: Reserved bit
- ESPI1: '1' enables serial peripheral interface1 interrupt; '0' disables serial peripheral interface1 interrupt.
- ES1: '1' enables serial port 1 interrupt; '0' disables serial port1 interrupt

5.6.2 Interrupt Priorities

Each interrupt source can also be individually programmed to four priority levels by setting or clearing two bits in the SFRs named IP, IPH and IP2HL (Interrupt Priority). Figure 5-11 shows how the IE, IE2, IP, IPH and IP2HL registers and the polling sequence work to determine the interrupt to be served. The IP, IPH, IP2HL and IE2 registers contain a number of unimplemented bits, these bits are reserved and user software should not write to these positions.



Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
IP		PT2	PSPI	PS0	PT1	PX1	PT0	PX0

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
IPH		PT2H	PSPIH	PSOH	PT1H	PX1H	РТОН	РХОН

-- : Reserved bit

- [IPH(6),IP(6)] : PT2 defines timer 2 interrupt priority level
- [IPH(5),IP(5)] : PSPI defines serial peripheral interface SPI0 interrupt priority level
- [IPH(4),IP(4)] : PS0 defines serial port 0 interrupt priority level
 - [IPH(3),IP(3)] : PT1 defines timer 1 interrupt priority level
 - [IPH(2), IP(2)] : PX1 defines external interrupt 1 priority level
 - [IPH(1),IP(1)] : PT0 defines timer 0 interrupt priority level
 - [IPH(0),IP(0)] : PX0 defines external interrupt 0 priority level

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
IP2HL	PSPI1_1	PSPI1_0	PS2_1	PS2_0			PS1_1	PS1_0

- -- : Reserved bit
- IP2HL(7:6) : PSPI1 defines serial peripheral port 1 SPI1 priority level
- IP2HL(1:0) : PS1 defines serial port 1 priority level

A low-priority interrupt can only be interrupted by a higher priority interrupt (but neither by another low-priority nor an equal-priority interrupt).

If two interrupt requests of different priority levels are received simultaneously, the request of higher priority is served. If interrupt requests of the same priority level are received simultaneously, an internal polling sequence determines which request is served. Thus within each priority level there is a second priority structure determined by the polling sequence as follows:

0.	IE0	(highest priority)
1.	TF0	
2.	IE1	ſ
3.	TF1	
4.	RI+TI	
5.	SPI0	
6.	TF2+EXF2	
7.	RI1+TI1	
9.	RI2+TI2	I
10.	SPI1	(lowest priority)



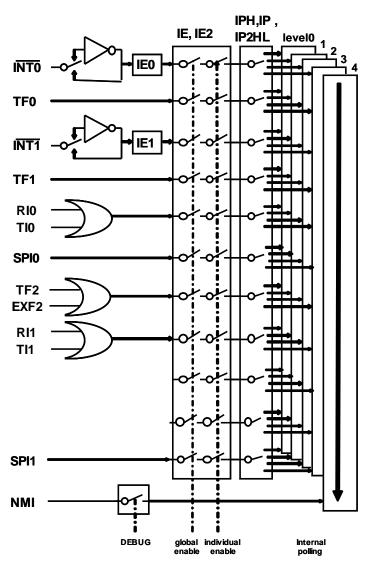
When in operation mode, all the interrupt flags are latched into the interrupt control system during state 2 of every machine cycle. The samples are polled during the following machine cycle. If the flag for an enabled interrupt is found to be set ('1'), the interrupt system generates an LCALL to the appropriate location in Program Memory, unless some other condition blocks the interrupt. Several conditions can block an interrupt, for example when an interrupt of equal or higher priority level is already in progress.

The hardware-generated LCALL causes the value in the Program Counter to be pushed into the stack, and reloads the PC with the beginning address of the service routine. As previously noted the service routine for each interrupt begins at a fixed location.

Only the Program Counter is automatically pushed onto the stack, not the PSW or any other register.

Having only the PC automatically saved it is a programmer task to save other registers if necessary

Figure 5-11. Interrupt Priorities diagram





5.6.3 Interrupt Handling

External Interrupts INT0 and INT1 can each be either level-activated or transition-activated, depending on bits IT0 and IT1 in Register TCON. When external interrupts are transition-activated the INTx pin must show a high level in one cycle and a low level in the next cycle to generate the interrupt, so that an input high or low should hold for at least 1 machine cycle to ensure sampling. The flags that actually generate these interrupts are bits IE0 and IE1 in TCON. The flags are cleared by hardware when the service routine is vectored only if the interrupts are transition-activated. When the interrupts are level-activated the interrupt flags are controlled by the external source, so that the external source has to deactivate the request before the interrupt service routine is completed, or else another interrupt will be generated. The response time is always more than 3 machine cycles and less than 7 machine cycles depending on program execution and interrupt status (see Figure 5-12).

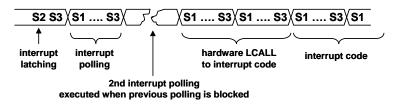
Timer 0 and Timer 1 Interrupts are generated by TF0 and TF1, which are set by a rollover. The flags are cleared by the on-chip hardware when the service routines are vectored.

Timer 2 interrupt (TF2) is user configurable, and depending on configuration flags will be cleared by hardware or by user software.

Serial Ports Interrupts are generated by the logical OR of RI and TI. They must be cleared by software.

SPI interrupts must be cleared by software.

Figure 5-12. Interrupt Handling



All of the bits that generate interrupts can be set or cleared by software. That is, interrupts can be generated or pending interrupts can be canceled by software. Each of these interrupt sources can be individually enabled or disabled by setting or clearing a bit in Special Function Register IE. IE also contains a global disable bit, EA, which disables all interrupts at once.

The interrupt flags are sampled at every machine cycle. The samples are polled during the following machine cycle. If one of the flags was in a set condition at the preceding cycle, the polling cycle will find it and the interrupt system will generate an LCALL to the appropriate service routine, provided this hardware-generated LCALL is not blocked by any of the following conditions:

- Condition 1. An interrupt of equal or higher priority level is already in progress.
- Condition 2. The current (polling) cycle is not the final cycle in the execution of the instruction in progress.
- Condition 3. The instruction in progress is RETI or any write to the IE or IP registers.

Any of these three conditions will block the generation of the LCALL to the interrupt service routine. Condition 2 ensures that the instruction in progress will be completed before vectoring to any service routine. Condition 3 ensures that if the instruction in progress is RETI or any access to IE or IP, then at least one more instruction will be executed before any interrupt is vectored to.

The polling cycle is repeated with each machine cycle, and the values polled are the values that were present at the previous machine cycle. The polling cycle/LCALL sequence is illustrated in Figure 5-12.

The hardware-generated LCALL executes a program jump to fixed program entries as shown below:



Source	Vector Address
IE0	0003H
TF0	000BH
IE1	0013H
TF1	001BH
RI0+TI0	0023H
SPI0	002BH
TF2	0033H
RI1+TI1	003BH
SPI1	0053H

Execution proceeds from that location until the RETI instruction is encountered then the interrupted program continues from where it left off.

5.7 I/O Ports

The ports P1, P3, P4 and P5 in the ADD8051C3A are bidirectional. Each port consists of a register (Special Function Registers P1, P3, P4, P5), an output driver, and an input buffer.

Pins in ports 3, 4 and 5 are multifunctional. They are port pins and also serve as input or output for the microcontroller peripherals:

Table 5-13. P	3, P4, P5	ports Alternate	Functions
---------------	-----------	-----------------	-----------

Pin	Alternate Function
P3.0	RxD (serial port 0 input)
P3.1	TxD (serial port 0 output)
P3.2	INT0 (external interrupt 0)
P3.3	INT1 (external interrupt 1)
P3.4	T0 (timer 0 external input)
P3.5	T1 (timer 1 external input)
P3.6	WR(external data memory write strobe)
P4.2	SS1 (SPI1 slave select input)
P4.3	SPICLK1 (SPI1 clock input/output)
P4.4	MOSI1 (SPI1 master out / slave in data)
P4.5	MISO1 (SPI1 master in / slave out data
P4.6	T2 (timer 2 input/output)
P5.0	SS (SPI0 slave select input)
P5.1	SPICLK (SPI0 clock input/output)
P5.2	MOSI (SPI0 master output / slave input data)
P5.3	MISO (SPI0 master input / slave output data)
P5.4	RxD (serial port 1 input)



P5.5 TxD (serial port 1 output)	ut)
---------------------------------	-----

Note: To use an Alternate Function the corresponding bit (or bits) in the SFR must contain a 1.

5.7.2 I/O Configurations

Figure 5-13 and Figure 5-14 show a functional diagram of bit register and I/O buffer in each of the four ports. The level of the port pin is placed on the internal bus and instructions can read the port pin value or the SFR register value.

If a bit in a register with AOF (Alternate Output Function) contains a 1, then the output level is controlled by the signal labeled "Alternate Output function".

Figure 5-13. Pins with AOF structure

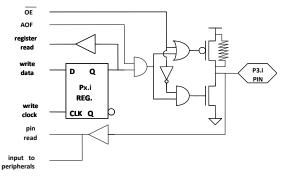
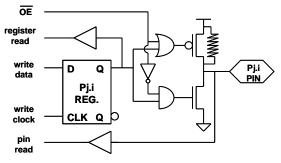


Figure 5-14. Pins without AOF structure



Pseudo-bidirectional mode

Ports 1, 3, 4 and 5 have internal pull-ups. Each I/O line can be independently used as an input or an output. To be used as an input, the port bit register must contain a '1'; When an SFR port bit is set to '1' the PMOS port driver is turned on one clock cycle and then is turned off. Then, the pin is pulled high by a weak internal pull-up, and can be pulled low by an external source. In the ADD8051C3A included in ATPL210A, the pull-up consists of 33Kohm resistors.

P1, P3, P4 and P5 registers are set to FFH by default (configured as inputs).

Push-Pull mode

Ports 3, 4 and 5 can be configured as Push-Pull output ports using P3M, P4M and P5M SFRs respectively.

Setting a bit in these SFRs to `1' configures the corresponding port as Push-Pull mode.

Setting a bit in these SFRs to '0' configures the corresponding port as Pseudo-bidirectional.

P3M, P4M and P5M are set to 00H by default (configured in Pseudo-bidirectional mode).



5.7.3 Read-Modify-Write Feature

Some instructions that read a port read the register and others read the pin. Only the read-modify-write instruction read the register. The instructions that read the register rather than the pin are the ones that read a value, possibly change it, and then rewrite it to the register. When the destination operand is a port, or a port bit, these instructions read the register rather than the pin: ANL, ORL, XRL, JBC, CPL, INC, DEC, DJNZ, MOV PX.Y,C, CLR PX.Y, SET PX.Y. The last three instructions are bit modify instructions. They read the port byte, all 8 bits, modify the target bit, then the modified byte is written to the register.

The-modify-write instructions are directed to the register to avoid a possible misinterpretation of the voltage level at the pin. When a port register is set to '1' to using the pin as input an external source can be pulling the pin to '0', if we need to know the configuration of the port we need to read the register and to know the state of the external source we need to read the pin.

5.7.4 Accessing External Memories

Accesses to external memory are of two types: accesses to external program code and accesses to external program data. Both share ports ADDRESS and DATA buses. Accesses to external program code use signal PSEN (program store enable) as the read strobe. Accesses to external program data use RD(P3.7) or WR(P3.6) to strobe the memory.

Fetches from external program memory always use a 16-bit logical address, fetches to external data can use 16-bit or 8-bit logical address.

5.8 Debug Mode

Debug mode is active when DBG pin is set to '1'. Debug mode is intended to implement software debugging tools.

DEBUG digital input is internally connected to DBG signal of the ADD8051C3A microcontroller

When in debug mode, a program can modify its own code. Thus, user **should not** work in debug mode unless explicitly stated otherwise.



6. Timers

The ADD8051C3A has a full set of timers, listed below:

- Two 16-bit configurable Timers/Counters: Timer 0 and Timer 1 that can be configured as timers or event counters (Figure 29). In ATPL210A this timers are internal devices, thus external signals counter configuration is not available. Nevertheless, the complete device operation is described below in order to know the suitable configuration registers values
- Two 8-bit auto reload counters Timers 11 and 12, used as baud rate generators for uarts 0 and 1 respectively (see Standard Serial Interfaces section).
- One 8-bit, with 18-bit preescaler, watchdog timer: Timer 3 (see Watchdog (timer 3) section).
- One 16-bit Timer/Counter Timer 2 with capture reload hardware (see 6.2 section).

6.1 Timer 0 and Timer 1

Timers 0 and 1 have 2-bit preescalers. These preescalers can be programmed to slow down count rate. The reset value set the preescaler values to the slowest count rate and the real count rate of the timers is the same than the timers in an "8051 legacy" microcontroller (12 clock cycles per machine cycle).

TPR, TMOD and TCON registers described below are involved in timer 0 and timer 1 management.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
TPR			T1P1	T1P0	T01P1	T01P0	T0P1	T0P0

- --: Reserved bit
- T1P(1:0): Rollover value for T1 preescaler when T1 is in modes 0, 1, 2 and 3
- T01P(1:0): Rollover value for TH0 preescaler when T0 is in mode 3
- T0P(1:0): Rollover value for T0 preescaler when T0 is in modes 0, 1, and 2, and for TL0 preescaler when T0 is in mode 3

When T0 or T1 are in the timer function, the registers are incremented every 1, 2, 3 or 4 machine cycles. Since a machine cycle consists of 3 oscillator periods, the count rate is 1/3, 1/6, 1/9 or 1/12 of the oscillator frequency.

In the counter function, the register is incremented in response to a 1-to-0 transition at its corresponding external input signal, T0(P3.4) or T1(P3.5). When in counter function, the external input is sampled once every machine cycle. When the samples show a high in one cycle and a low in the next cycle (1-to-0 transition), the count is incremented. The register value is updated during the cycle following the one in which the transition was detected, the maximum count rate is 1/2 of the machine cycle frequency.

Either the timer or counter function is selected by control bit C/T in the Special Function Register TMOD where T0 and T1 have four operating modes, which are selected by bit-pairs (M1, M0) in TMOD.

	TIMER 1			TIMER 0				
Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
TMOD	GATE	C/T	M1	MO	GATE	C/T	M1	MO

- GATE:
 - '1': Enables gate control. Device count is enabled only when both 'INTi' pin and "TCON.TRi" are high.
 - '0': Disables gate control. Device count is enabled when "TCON.TRi" is high.



- C/T:
 - '1': Counter operation mode. Input from Ti input pin
 - '0': Timer operation mode. Input from internal clock system
- M(1:0): Operation modes
 - "00": Mode 0: 8048 13bit timer, TLi used as 5 bit preescaler
 - "01": Mode 1: 16 bit timer/counter composed by concatenated THi and TLi
 - "10": Mode 2: 8 bit with auto reload (TLi <= THi) timer/counter
 - "11": Mode 3: timer 1 stopped
 - Timer 0 runs as two 8-bit independent timer/counters:

TL0 is a timer/counter controlled by timer 0 control bits

TH0 is a timer controlled by timer 1 control bits

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
TCON	TF1	TR1	TF0	TR0	IE1	IT1	IE0	IT0

- TF1: Timer 1 overflow flag, set by hardware overflow and cleared by software or by hardware when interrupt routine is vectored
- TR1: Timer 1 run control bit, when set turns on timer/counter 1
- TF0: Timer 0 overflow flag, set by hardware overflow and cleared by software or by hardware when interrupt routine is vectored
- TR0: Timer 0 run control bit, when set turns on timer/counter 0
- IE1: Interrupt 1 edge flag, set by hardware when an external interrupt edge is detected and cleared by software or by hardware when interrupt routine is vectored
- IT1: Interrupt 1 type control bit. "'1'/'0" to specify "falling edge/low level" trigger for external interruption 1
- IE0: Interrupt 0 edge flag, set by hardware when an external interrupt edge is detected and cleared by software or by hardware when interrupt routine is vectored
- IT0: Interrupt 0 type control bit. "'1'/'0'" to specify "falling edge/low level" trigger for external interruption 0

6.1.1 Timer Mode 0

The timer (0 or 1) into Mode 0 looks like a 13-bit timer: an 8-bit counter (all 8 bits of THx) with a divide-by-32 preescaler (the lower 5 bits of TLx), see Figure 6-1. As the count rolls over from all 1s to all 0s, it sets the Timer interrupt flag TFi. The counter input is enabled to the Timer when TRi=1 and either GATE=0 or INT1=1. When GATE=1 the Timer can be controlled by external input INTx, to facilitate pulse width measurements. TRx is a control bit in the Special Function Register TCON.

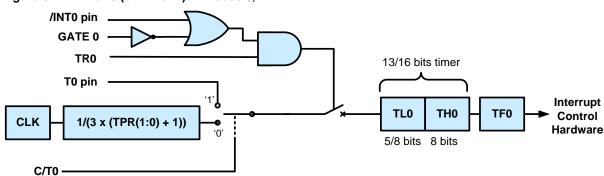


Figure 6-1. Timer 0 (or Timer 1) in modes 0,1

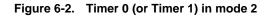


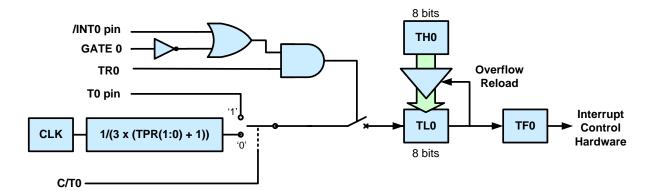
6.1.2 Timer Mode 1

In Mode 1 the timer (0 or 1) runs in the same way than Mode 0 but they are 16 bits-width and TLx is not used as a preescaler.

6.1.3 Timer Mode 2

Mode 2 configures the timer (0 or 1) register as an 8-bit Counter (TLx) with automatic reload of THx value. THx value is preset by software.





6.1.4 Timer Mode 3

In Mode 3 Timer 0 registers TL0 and TH0 are configured as two separate 8-bit counters, TL0 and TH0.

TL0 uses the Timer 0 control bits and can be configured either in timer or in counter mode.

TH0 is locked into a timer function (counting machine cycles) and is controlled by Timer 1 control bits TR1 and TF1.

TH0 controls the Timer 1 interrupt.

Timer 1 in Mode 3 holds its count. With Timer 0 in Mode 3, the number of timer/counters in the 80C51 is increased to three. When Timer 0 is in Mode 3, Timer 1 can be used by the serial port as a baud rate generator, or in any application not requiring an interrupt.



6.2 Timer 2

Timer 2 is a 16-bit timer/counter. The count is maintained by two eight-bit timer registers, TH2 and TL2, connected in cascade. T2MOD and T2CON registers described below control the operation of timer 2.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
T2MOD			TF2SC	EXF2SC	T2PR1	T2PR0	T2OE	DCEN

- --: Reserved bit
- TF2SC: Timer 2 overflow flag software control. When set to '1', TF2 flag must be cleared by software. When cleared to '0', TF2 flag is cleared by hardware when interrupt routine is vectored
- EXF2SC: External Interrupt 2 flag software control. When set to '1', EXF2 flag must be cleared by software. When cleared to '0', EXF2 flag is cleared by hardware when interrupt routine is vectored
- T2PR(1:0): Timer 2 preescaler
- T2OE: Timer 2 output enable bit. In the timer 2 clock-out mode connects the programmable clock input to the external pin T2
- DCEN: Down count enable bit, configures timer 2 as an up/down counter

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
T2CON	TF2	EXF2	SSI1	SSI0	EXEN2	TR2	C/T2	CP/RL2

- TF2: Timer 2 overflow flag, Set by timer 2 overflow. It must be cleared by software. TF2 is not set when timer 2 is in baud rate generator mode
- EXF2: Timer 2 external flag. If EXEN2='1' capture or reload caused by a negative transition on T2EX sets EXF2. EXF2 does not cause an interrupt in up/down counter mode (DCEN='1')
- SSI(1:0): Select T2 as baud rate generator for serial port
 - "00": none
 - "01": USART0
 - "10": USART1
 - "11": USART2
- EXEN2: Timer 2 external enable bit. Setting EXEN2 causes a capture or reload to occur as a result of a negative transition on T2EX unless timer 2 is being used as baud rate generator. Clearing EXEN2 causes timer 2 to ignore events at T2EX
- TR2: Timer 2 run control bit. Setting this bit starts the timer
- C/T2: Timer 2 counter/timer select. When set to '1', it selects counter operation then timer 2 counts negative transitions on system clock
- CP/RL2: When set to '1', captures occur on negative transitions at T2EX if EXEN='1'; When cleared to '0', auto reload occurs on timer 2 overflow or negative transitions at T2EX if EXEN2='1'. If RCLK='1' or TCLK='1' the CP/RL2 bit is ignored

Timer 2 provides the following operating modes: capture mode, auto-reload mode, baud rate generator mode, and programmable clock-out mode. Select the operating mode with T2MOD and T2CON register bits as shown in Table 6-1. Auto-reload is the default mode. Setting T2CON(5:4) selects the baud rate generator mode.



Table 6-1. Timer 2 operating modes

Mode	T2CON(5) OR T2CON(4)	CP/RL2	T2OE
Auto-reload	0	0	0
Programmable clock-out	0	0	1
Capture	0	1	0
stopped	0	1	1
Baud Rate generator	1	х	х

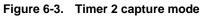
Timer 2 operation is similar to timer 0 and timer 1. C/T2 selects between timer operation mode (internal clock input) or counter operation mode (external pin T2 as the time register input). Setting TR2 allows TL2 to be incremented by the selected input.

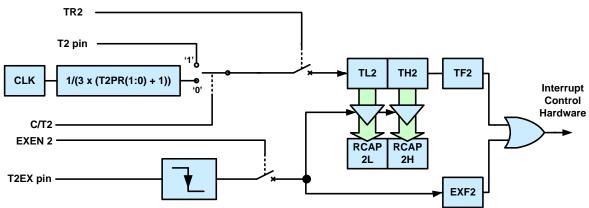
The timer 2 related interrupt flags (TF2, EXF2) can be configured to be cleared by hardware when the interrupt routine is vectored to, or by user software using TF2SC and EXF2SC.

The operating modes are described in the following paragraphs.

6.2.2 Capture mode

In the capture mode, timer 2 operates as a 16-bit timer or counter. An overflow condition sets bit TF2, which can be used to request an interrupt. Setting the external enable bit EXEN2 allows the RCAP2H and RCAP2L registers to capture the current value in timer registers TH2 and TL2 in response to a 1-to-0 transition at external input T2EX. The transition at T2EX also sets bit EXF2 in T2CON. The EXF2 bit, like TF2, can generate an interrupt.





6.2.3 Auto-Reload mode

The auto-reload mode configures timer 2 as a 16-bit timer or event counter with automatic reload. The timer operates as an up counter or as an up/down counter, as determined by the down counter enable bit (DCEN). When reset occurs, DCEN is cleared, so in the auto-reload mode, timer 2 in configured as an up counter by default.

• Up Counter Operation

When DCEN = 0, timer 2 operates as an up counter.



The external enable bit EXEN2 in the T2CON register provides two options. If EXEN2 = 0, timer 2 counts up to FFFFH and sets the TF2 overflow flag. The overflow condition loads the 16-bit value in the reload/capture registers (RCAP2H, RCAP2L) into the timer registers (TH2, TL2). The values in RCAP2H and RCAP2L are preset by software.

If EXEN2 = 1, the timer registers are reloaded by either a timer overflow or a high-to-low transition at external input T2EX. This transition also sets the EXF2 bit in the T2CON register.

Either TF2 or EXF2 bit can generate a timer 2 interrupt request.

Up/Down Counter Operation

When DCEN = 1, timer 2 operates as an up/down counter.

External pin T2EX controls the direction of the count.

When T2EX is high, timer 2 counts up. The timer overflow occurs at FFFFH which sets the timer 2 overflow flag (TF2) and generates an interrupt request. The overflow also causes the 16-bit value in RCAP2H and RCAP2L to be loaded into the timer registers TH2 and TL2.

When T2EX is low, timer 2 counts down. Timer underflow occurs when the count in the timer registers (TH2, TL2) equals the value stored in RCAP2H and RCAP2L. The underflow sets the TF2 bit and reloads FFFFH into the timer registers.

The EXF2 bit toggles when timer 2 overflows or underflows changing the direction of the count. When timer 2 operates as an up/down counter, EXF2 does not generate an interrupt. This bit can be used to provide 17-bit resolution.

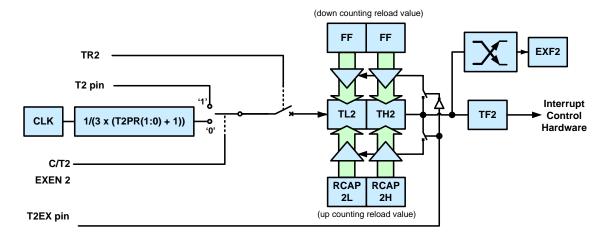


Table 6-2. Timer 2 Auto-reload mode

6.2.4 Clock-Out mode

In the clock-out mode, timer 2 operates as a 50%-duty-cycle, variable-frequency clock. The input clock increments TL2 at frequency FOSC/2 (where FOSC=20MHz). The timer repeatedly counts to overflow from a preloaded value. At overflow, the contents of the RCAP2H and RCAP2L registers are loaded into TH2/TL2. In this mode, timer 2 overflows do not generate interrupts. The formula gives the clock-out frequency as a function of the system oscillator frequency and the value in the RCAP2H and RCAP2L registers:

For a 20 MHz system clock, timer 2 has a programmable frequency range of 12.7 Hz to 3.3 MHz. The generated clock signal is brought out to the T2 pin (P4.6).

Timer 2 is programmed for the clock-out mode as follows:

1. Set the T2OE bit in T2MOD. This gates the timer register overflow to the \div 2 counter.



2. Clear the C/T2 bit in T2CON to select FOSC/3*(T2PRE(1:0)+1) as the timer input signal. This also gates the output of the ÷2 counter to pin T2.

3. Determine the 16-bit reload value from the formula and enter in the RCAP2H/RCAP2L registers.

4. Enter a 16-bit initial value in timer register TH2/TL2. This can be the same as the reload value, or different, depending on the application.

5. To start the timer, set the TR2 run control bit in T2CON.

Clock-out mode timer 2 operation is similar to timer 2 operating as a baud rate generator. Moreover, it is possible to use timer 2 as a baud rate generator and a clock generator simultaneously. For this configuration, the baud rates and clock frequencies are not independent since both functions use the values in the RCAP2H and RCAP2L registers

The baud rate is expressed by the following formula:

6.2.5 Baud rate Generator mode

This mode configures timer 2 as a baud rate generator for its use with the serial ports. Select this mode by setting the SSI(1:0) bits in T2CON.

Timer 2 may be selected as the baud rate generator for the transmitter and/or receiver in modes 1 and 3. The timer 2 baud rate generator mode is similar to the auto-reload mode. A rollover in the TH2 register reloads registers TH2 and TL2 with the 16-bit value in registers RCAP2H and RCAP2L, which are preset by software.

The timer 2 baud rate is expressed by the following formula:

Where "BAUD2" is the bit in the Special Function Register CONF and RCAP2(H,L) denotes the content of RCAP2H and RCAP2L taken as a 16-bit unsigned integer.

To select timer 2 as baud rate generator for transmission and reception, program the SSI(1:0) bits in the T2CON register as shown in Table 6-1. Setting SSI(1) and/or SSI(0) puts timer 2 into its baud rate generator mode. In this mode, a rollover in the TH2 register does not set the TF2 bit in the T2CON register. Also, a high-to-low transition at the T2EX pin sets the EXF2 bit in the T2CON register but does not cause a reload from (RCAP2H, RCAP2L) to (TH2, TL2).

User shall use the T2EX pin as an additional external interrupt by setting the EXEN2 bit in T2CON.

Turn the timer off (clear the TR2 bit in the T2CON register) before accessing registers TH2, TL2, RCAP2H, and RCAP2L

User shall configure timer 2 as a timer or a counter. In most applications, it is configured for timer operation (i.e., the C/T2 bit is clear in the T2CON register).

Note that timer 2 increments every 1, 2, 3 or 4 state times (3TOSC) when it is in the baud rate generator mode.

When timer 2 is configured as a timer and in baud rate generator mode, do not read or write the TH2 or TL2 registers. The timer is being incremented every state time, and the results of a read or write may not be accurate. In addition, user shall read, but not write to, the RCAP2 registers; a write may overlap a reload and cause write and/or reload errors.

Table 6-3 lists commonly used baud rates and shows how they are generated by timer 2.



Baud rate	T2PRE	RCAP2H	RCAP2L
230400	0	FF	FE
115200	0	FF	FC
57600	0	FF	F9
38400	0	FF	F5
19200	0	FF	EA
9600	0	FF	D5
4800	0	FF	A9
2400	0	FF	52
1200	0	FE	A5
600	0	FD	4A
300	0	FA	93
150	0	F5	26
110	0	F1	34

Table 6-3. Timer 2 generated baud rates with a 20MHz oscillator

6.3 Watchdog (timer 3)

The watchdog timer includes a 18-bit preescaler which is set to 0 when T3 is reloaded.

T3 is active only when /EWDG pin is tied to '0'.

T3 must be reloaded by software to avoid T3 rollover and program restart.

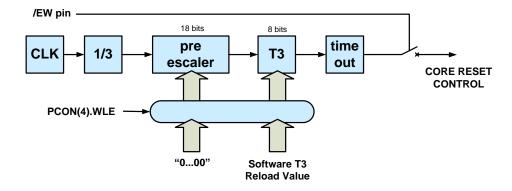
Reload of T3 is allowed only if WLE watchdog load enable (PCON(4)) is set.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
PCON	SMOD			WLE	GF1	GF0	PD	IDL

- SMOD: Doubles baud rate bit when timer 1 is used as time generator and serial port is in modes 1,2 or 3
- --: Reserved bit
- WLE: Watchdog load enable. It must be set by software to enable T3 reload. WLE is reset by hardware 13 machine cycles after been set by user software.
- GF1: General purpose flag bit, user programmable
- GF0: General purpose flag bit, user programmable
- PD: Sets power down mode
- IDL: Sets idle mode

Reload of T3 register automatically resets the preescaler and WLE







7. Standard Serial Interfaces

Two full duplex (simultaneous transmission/reception) serial ports are implemented in ATPL210. The serial port reception and transmission registers are both accessed at Special Function Register SBUF(0,1). Writing to SBUF loads the transmit register and starts transmission, while reading SBUF accesses a physically separate receive register. The reception register is buffered and it can start reception of a second byte before a previously received byte has been read from the register. When a second byte is fully received the first one is lost.

The serial port control and status register is the Special Function Register SCON(0,1).

This register contains the mode selection bits, the 9th data bit for transmit and receive (TB8 and RB8), and the serial port interrupt bits (TI and RI).

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
SCON	SM0	SM1	SM2	REN	TB8	RB8	TI	RI

• SM(0:1):

"00": Mode 0: Shift register mode; baud rate=1/12 fclk

"01": Mode 1: 8-bit UART mode; baud rate= T1 overflow

"10": Mode 2: 9-bit UART mode; baud rate= 1/32 or 1/64 fclk

"11": Mode 3: 9-bit UART mode; baud rate= T1 overflow

- SM2: Reception control. In modes 2 and 3 with SM2='1', RI will be '0' if the 9th received bit is '0'; In mode 1 with SM2='1', RI will be '0' if the received stop bit is not correct; In mode 0 SM2 keeps a '0'.
- REN: Serial port reception enable
- TB8: 9th transmitted bit in modes 2 and 3
- RB8: 9th received bit in modes 2 and 3, received stop bit in mode 1 when SM2='0'
- TI: Transmission interrupt flag. Set by hardware at the end of transmission, must be cleared by software
- RI: Reception interrupt flag. Set by hardware at the end of transmission, must be cleared by software

The serial port can operate in 4 modes. Transmission is always initiated by any instruction that writes SBUF, and TI flag is set after transmission. Reception is initiated in Mode 0 by the condition RI=0 and REN=1, and in the other modes by the incoming start bit if REN=1. In Modes 2 and 3, 9 data bits are received. The 9th one goes into RB8 and then comes a stop bit.

The port can be programmed such that when the stop bit is received, the serial port interrupt will be activated only if RB8 = 1. This feature is enabled by setting bit SM2 in SCON.

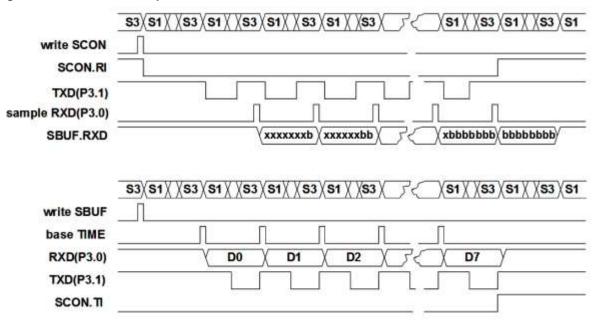
7.1 Serial Port modes

7.1.1 Mode 0 (shift register mode)

Serial data inputs and exits through RxD. TxD outputs the shift clock. Eight bits are transmitted/received (LSB first). The baud rate is fixed at 1/12 the oscillator frequency. The internal timing is such that four full machine cycles will elapse between "write to SBUF" command and activating the transmission; transmission ends in the 36th machine cycle after "write to SBUF". Reception is initiated by the conditions REN = 1 and RI = 0, and after 36 machine cycles reception finish and RI is set to '1'.



Figure 7-1. Serial Port reception/emission mode 0



7.1.2 Mode 1 (8-bit UART)

In this mode, 10 bits are transmitted (through TxD) or received (through RxD): a start bit (0), 8 data bits (LSB first), and a stop bit (1). In reception, the stop bit goes into RB8 in SFR SCON. The baud rate is controlled by T1 rollover. Transmission is initiated by any instruction that uses SBUF as a destination register and the bit times are synchronized to T1 rollover. Reception is initiated by a detected 1-to-0 transition at RxD, RxD is sampled at a rate of 16 times each bit time and the value accepted is the value that was seen in at least 2 of the 3 central samples (7th, 8th, and 9th).

A correct reception loads SBUF and RB8 and sets RI, this is achieved when at the end of the receive process the condition (RI='0' and (SM2='0' or received_stop_bit='1')) is met. Otherwise, if the condition is not met, the received frame is irretrievably lost.

7.1.3 Modes 2 and 3 (9-bit UART)

In this mode, 11 bits are transmitted (through TxD) or received (through RxD): start bit (0), 8 data bits (LSB first), a programmable 9th data bit (TB8), and a stop bit (1). The 9th transmitted data bit is loaded with TB8 value from SFR SCON. On reception, the 9th data bit goes into RB8 in SFR SCON. The stop bit is ignored. The baud rate is programmable to 1/32 or 1/64 the oscillator frequency. Mode 3 is the same as Mode 2 with programmable baud-rate controlled by T1 rollover.

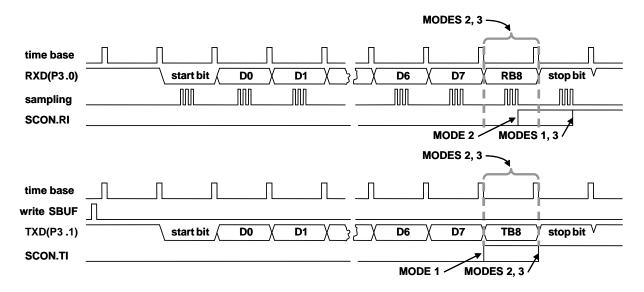
The transmission begins with a "write to SUBF" instruction and finish at the 11th rollover after "write to SUBF", TI is set when transmission ends.

Reception is initiated by detecting a high to low level transition at RxD. For this purpose RxD is sampled at a rate of 16 times each bit. The value accepted is the value that was seen in at least 2 out of 3 central samples. When reception ends SBUF and RB8 are loaded and RI is set. A frame is correctly received if, and only if, the following condition is met at the time the last bit is received (RI=0 and (SM2='0' or 9th_data_bit=1).

Otherwise, if the condition is not met the received frame is irretrievably lost and RI is not set.



Figure 7-2. Serial Port reception/emission modes 1,2,3.



7.2 Serial Port Timers (Timers 11 & 12)

Timers 11 and 12 are similar to T1 configured as 8-bit timer in auto reload mode (mode 2). These timers are used

as baud rate generators for uarts 0 and 1 respectively.

These timers are **not connected to the interruption system hardware**. They have not preescaler counter and their values are incremented every machine cycle (3 clock cycles).

They can be used only as baud rate generators for the standard serial interfaces. When the run control bit in T1NC register (T1iR) is set the corresponding timer (T1i) starts counting machine cycles and the related SSI is controlled by the overflow of T1i.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
T1NC	WT12	WT11	WT10				T12R	T11R

WT1(2:0): read/write access control to T1 timer registers

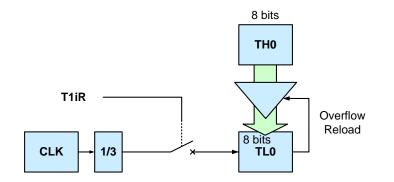
"000": TH1, TL1 "001": TH11, TL11 "010": TH12, TL12 "100": TH14, TL14

- --: Reserved bit.
- T12R: TT2 run control, when set T12 runs and T12 overflow controls UART1 variable baud rates.
- T11R: TT1 run control, when set T11 runs and T11 overflow controls UART0 variable baud rates.

Registers TH11, TH12 and TH14 are mapped in the same address than TH1, to select the specific register to be accessed (read or write) the bits WT1(2:0) in T1NC register must be set to the desired value. TL11, TL12, TL14 and TL1 are accessed in the same way.

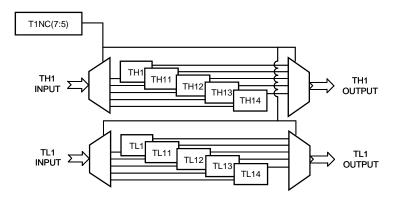


Figure 7-3. T1i timer



Timers T1i count machine cycles (3 clock cycles) in normal operation mode. They can be forced to count clock cycles. To enter this mode the corresponding bits in T1NP register must be set to '1' (T1NP(i-1) controls T1i). Unused bits in T1NP, bits 4 to 7, are not defined. Writing to these bits has no effect, reading these bits will always get "0000".

Figure 7-4. TH1 and TL1 registers



Mode 0

The baud rate in Mode 0 is fixed to Oscillator Frequency / 12.

Modes 1 and 3

The baud rates in Modes 1 and 3 are determined by the Timer 1 overflow rate and the value of SMOD and BAUD2 as follows:

Mode 2

The baud rate in Mode 2 depends on the value of bit SMOD in Special Function Register PCON and BAUD2 in Special Function Register CONF.

If [BAUD2,SMOD] = "00" (which is the value on reset), the baud rate is 1/64 the oscillator frequency.

If [BAUD2,SMOD] = "01" or "10", the baud rate is 1/32 the oscillator frequency.

If [BAUD2,SMOD] = "11", the baud rate is 1/16 the oscillator frequency



In the most typical applications, it is configured for "timer" operation in the auto-reload mode. In that case the baud rate is given by the formula:

When timers T11, T12, T13 or T14 are selected as timing sources the baud rate is given by the formula:

_

Timers T1i count machine cycles (3 clock cycles) in normal operation mode. They can be forced to count clock cycles. To enter this mode the corresponding bits in T1NP register must be set to '1'. When T1i are configured in this mode the baud rate is given by the formula:

In this mode 0xFF and 0xFE are not allowed as reload values. When these values are used the behavior of USART is unpredictable.

Timer 2 overflow can also be used as baud rate generator for serial ports in modes 1 and 3. See timer 2 section for more information.

Selection of T1i as timing source overrides T2 and T1 sources. Selection of T2 as timing source overrides T1 source.



8. Serial Peripheral Interfaces SPI0 and SPI1

Two Serial Peripheral Interfaces SPI0 and SPI1 are implemented in ATPL210A.

Both SPI0 and SPI1 are full-duplex, synchronous communication bus with two operation modes: master mode and slave mode. When acting as an SPI master, the microcontroller supports baud rates of up to ¹/₄ of the clock frequency.

The main difference between SPI0 and SPI1 is that while SPI0 is a byte oriented SPI, SPI1 is a byte-buffered SPI. SPI1 byte-buffering will be explained as a special SPI capability type in section 8.3. Ignoring this buffer capability, both SPI0 and SPI1 are identical internally.

8.1 SPI description

During a byte transfer over SPI, the master and the slave simultaneously transmit (shift out) and receive (shift in) data. Master is the responsible in SPI system for transfer initialization, proper shifting and sampling of data. Every activity in the SPI environment is synchronized with the clock signal which is under master's control. A slave select line allows the master to select the desired slaves to exchange data with; not selected slave devices cannot interfere ongoing bus activities. In systems with multiple masters, slave select line on the master device can be used to detect multiple master bus contention.

SPI0 is configured and controlled using three special function registers: SPCTL, SPSTAT and SPDAT. (Similarly, SPI1 is configured and controlled by means of SPCTL1, SPSTAT1 and SPDAT1N registers)

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
SPCTL	SSIG	SPEN	DORD	MSTR	CPOL	CPHA	PSC1	PSC0

- SSIG: /SS0 slave select ignore, if set to '1' MSTR decides whether the device is a master or a slave and /SS0 pin can be used as port pin; if set to '0' and MSTR is set to '1' then the /SS pin decides whether the device is a master or a slave
- SPEN: SPI0 enabled '1' or disabled '0', when the SPI0 is disabled the SPI0 pins can be used as general I/O pins
- DORD: SPI0 data order
 - '1' the lsb of the data byte is transmitted first
 - '0' the msb of the data byte is transmitted first
 - MSTR: master '1' / slave '0' mode select. See SSIG field.
- CPOL: SPI0 clock polarity

 $^{\prime}\text{1}^{\prime}-\text{SPICLK0}$ is high when idle, the leading edge of SPICLK0 is the falling edge and the trailing edge is the rising edge

 $^{\prime 0'}$ – SPICLK0 is low when idle, the leading edge of SPICLK0 is the rising edge and the trailing edge is the falling edge

CPHA: SPI0 clock phase select

 $^{\prime}\mbox{1'}$ - data is driven on the leading edge of SPICLK0 and is sampled on the trailing edge

'0' - the first data bit is on the MOSI0/MISO0 line before the first SPICLK0 leading edge, data is sampled on the leading edge of SPICLK0 and driven on the trailing edge of SPICLK

 PSC(1:0): SPI0 clock rate select, determines master clock output. 8 different baud rates can be selected using PSC2[SPSTAT(5)] & PSC(1:0)[SPCTL(1:0)].

When working as slave, PSC(2:0) value is not taken into account, nevertheless this value must be different from "000" $\,$



PSC(2:0)	clock rate
"000"	fosc/4
"001"	fosc/16
"010"	fosc/64
"011"	fosc/128
"100"	fosc/256
"101"	fosc/512
"110"	fosc/1024
"111"	fosc/2048

N	lame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
SF	PSTAT	SPIF	WCOL	PSC2					

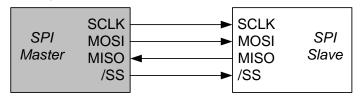
- SPIF: SPI0 transfer completion flag. When ESPI (IE.5) bit and EA (IE.7) bit are '1' and a SPI0 transfer finishes then the SPIF bit is set by hardware and an interrupt is generated. SPIF will also be set when SPI0 is in master mode with SSIG(SPCTL.7)='0' and /SS0(slave select) pin is driven low
- WCOL: SPI0 write collision flag. The WCOL bit is set by hardware when SPDAT is written during a data transfer. WCOL flag is cleared in software by writing '0' to this bit
- PSC2: SPI0 master clock rate select msb
- --: Reserved bit

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
SPDAT	B7	B6	B5	B4	B3	B2	B1	B0

• B(7:0): Data transferred. B7 msb; B0 lsb.

The SPI interface requires four pins: SPICLK, MOSI, MISO and SS:

Figure 8-1. SPI connection diagram



SPI0 pins are mapped to P5 port pins as alternate output functions

Table 8-1. SPI0 port map

ADD8051C3A Port	Alternate Function			
P5.0	/SS0			
P5.1	SPICLK0			
P5.2	MOSI0			
P5.3	MISO0			



SPICLK, MOSI and MISO are typically tied together between two or more SPI devices. Data flows from master to slave on the MOSI (Master Out Slave In) pin and flows from slave to master on the MISO (Master In Slave Out) pin. The SPICLK signal is output in the master mode and is input in the slave mode. If the SPI system is disabled, i.e. SPEN (SPCTL.6) = 0 (reset value), these pins can be used as general purpose I/O pins.

SS is an optional slave select pin. In a typical configuration, an SPI master asserts one of its port pins to select one SPI device as the current slave. An SPI slave device uses its SS pin to determine whether it is selected or not. The SS is ignored if any of the following conditions are true:

- The SPI0 system is disabled, i.e. SPEN (SPCTL.6) = 0 (reset value)
- SPI0 is enabled but SS pin is not needed in such system, i.e. SSIG(SPCTL.7) = 1; in this case SS pin can be used as general purpose pin on P5.

8.1.2 SPI clock phase, polarity and operation

There are four combinations (CPHA, CPOL) for sampling and shifting activities on data and clock lines in the SPI. Master can switch between any of them at any time thus having ability to communicate with slaves supporting data transfer using different modes.

In order to have successful data transfer between SPI devices, proper selection of the SPI clock phase and polarity is crucial. It is important to note that some of these configurations offer less capabilities than others. Clock Phase Bit CPHA allows user to specify the edges for sampling and shifting data. Clock Polarity bit CPOL allows user to set the clock polarity Figure 8-2 and Figure 8-3 show transfers with different values of CPHA and CPOL.

The SPI clock preescaler selection uses the PSC1-PSC0 bits in the SPCTL register and PSC2 in SPSTAT register. Master selects one of the available baud rates for the SPI communication. SPI Clock Preescaler bits do not have affect on the part acting as a slave, since it uses SPI clock supplied by the master.

When microcontroller operates as a slave with CPHA='0', some restrictions are present.

- SSIG must be '0' and the SS pin must be negated and reasserted between each successive byte transfer.
- If the SPDAT register is written while SS is active (low), a write collision error results.
- Microcontroller's behavior is undefined if CPHA is '0' and SSIG is '1'.

On the other hand, slave having CPHA='1' may set SSIG to '1'. If SSIG = 1, the SS pin may remain active low between successive transfers (can be tied low at all times). This format is sometimes preferred in systems having a single fixed master and a single slave driving the MISO data line. Microcontroller configured as a master with CPHA='0' does not need to negate and reassert slave's SS line in order to send and receive byte(s) of data.

In SPI, transfers are always initiated by the master. If the SPI is enabled (SPEN = 1) and microcontroller is configured as SPI master, writing to the SPI data register by the master will start the SPI clock generator and data transfer. The data will start to appear on MOSI about one half SPI bit-time to one SPI bit-time after data is written to SPDAT.





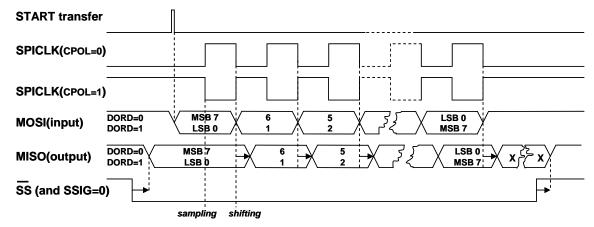
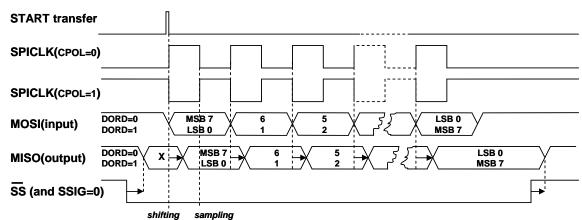


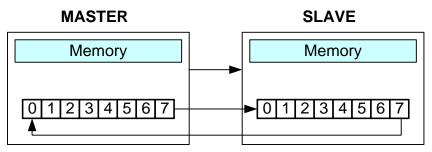
Figure 8-3. SPI0 transfer with CPHA=1



Note that the master selects a slave by driving the slave select pin of the corresponding slave device. Data written to the SPDAT register of the master is shifted out of the MOSI pin of the master to the MOSI pin of the slave, at the same time the data in SPDAT register on slave side is shifted out on its MISO pin to the MISO pin of the master.

After shifting one byte, the SPI clock generator stops, setting the transfer completion flag (SPIF) and an interrupt will be created if the SPI interrupt is enabled (ESPI, or IEN1.3 = 1). The two shift registers in the master CPU and slave CPU can be considered as one distributed 16-bit circular shift register. When data is shifted from the master to the slave, data is also shifted in the opposite direction simultaneously. This means that during one shift cycle, data in the master and the slave are interchanged.







If SPEN = 1, SSIG = 0 and MSTR = 1, the SPI0 is enabled in master mode. The SS pin is quasi-bidirectional. In this case, another master can drive this pin low to select this device as an SPI slave and start sending data to it. To avoid bus contention, the CPU becomes a slave. As a result of the observed SPI module becoming a slave, the MOSI and SPICLK pins are forced to be an input and MISO becomes an output.

The SPIF flag in SPSTAT is set, and if the SPI interrupt is enabled an SPI interrupt will occur.

User software should always check the MSTR bit. If this bit is cleared by a slave select and the user wants to continue to use the SPI as a master, the user must set the MSTR bit, otherwise it will stay in slave mode.

8.1.3 SPI0 write collision

The SPI0 is single buffered in the transmit direction and double buffered in the receive direction. New data for transmission cannot be written to the shift register until the previous transaction is completed. The WCOL (SPSTAT.6) bit is set to indicate data collision when the data register is written during transmission. In this case, the data currently being transmitted will continue to be transmitted, but the new data, i.e., the one causing the collision, will be lost.

While write collision is detected for either a master or a slave, it is uncommon for a master because the master has full control of the transfer in progress. The slave, however, has no control over transmission start by the master and therefore collision can occur.

Receiver transfers received data into a parallel read data buffer so that the shift register is free to accept a second character. However, the received character must be read from the Data Register before the next character has been completely shifted in. Otherwise. the previous data is lost. WCOL can be cleared in software by writing a '0' to this bit.

8.2 SPI Modes

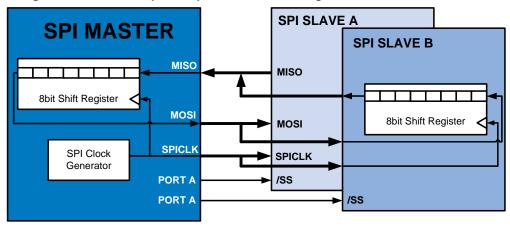
The communication using SPI in a single-master system is simple and usually works as described below:

- Both master and slave(s) configure their SPIs to operate in the same mode (one of four available modes) and turn them on; additionally, the master has to select baud rate for the communication; slave uses master's serial clock to sample and shift data during the transfer
- Master selects desired slave unit(s) using its/their SS line(s)
- As soon as the master writes to its SPI buffer register (assuming that slave(s) has/have already loaded data into its/their buffer(s)), transfer is initiated; master's SPI module generates serial clock and master's and slave's data are exchanged
- After the end of transfer, an indication is generated in all participating SPI units, both the master and the slave(s) read received data from their buffers
- If there is more data to be exchanged, all units taking part in this communication prepare the new set of data and master initiates another round of transfer; if there is no more data to be exchanged, master deselects used SS line(s) and previously active slaves are deselected; this is the end of SPI transfer

Typical connection in system using SPI is shown below.



Figure 8-5. SPI single master with multiple independent slaves configuration



In Figure 8-5, slave's SSIG = 0 (see 8.1), and SS is used to select the slave. The SPI0 master can use any port pin (including P5.0/SS with SSIG = 1) to drive the SS pin.

Multimaster configuration is the case where two devices are connected to each other and either device can be a master or a slave. When no SPI0 operation is occurring, both can be configured as masters (MSTR = 1) with SSIG = 0 and P5.0 (SS) being in quasi-bidirectional mode. Before device initiates a transfer, it must set SSIG = 1 in order to avoid its own mode change since it will drive P5.0 low, forcing a mode change in the other device to slave.

Table 8-2 shows configuration for the master/slave modes in various cases.

SPEN	SSIG	SS pin	MSTR	operation	MISO	MOSI	SPICLK	comment
0	х	P5.0 ⁽¹⁾	Х	disabled	P5.3 ⁽¹⁾	P5.2 ⁽¹⁾	P5.1 ⁽¹⁾	SPI0 disabled, P5(3:0) are used as port pins
1	0	0	0	slave	out	in	in	selected as slave
1	0	1	0	slave	Н	in	in	not selected, MISO is H to avoid bus contention
1	0	0	1(=>0) ⁽²⁾	slave	out	in	in	the uC is configured as a master, then it is selected by SS to be a slave, MSTR bit is cleared and the uC becomes a slave
1	0	1	1	master	in	out	out	
1	1	P5.0 ⁽¹⁾	0	slave	out	in	in	slave not needing SS
1	1	P5.0 ⁽¹⁾	1	master	in	Н	Н	master not needing SS

Table 8-2.	SPI modes
------------	-----------

Notes: 1. function as port or as alternate output function

2. the MSTR bit changes to '0' automatically when SS becomes low in input mode and SSIG is 0



8.3 SPI1 buffer operation

SPI1 is a byte buffered serial peripheral interface. Data transfers have a selectable additional number of bytes ranging from 0 to 7, SPSTAT1(2:0) is used to select the number of additional bytes. The buffer is composed of the registers SPDAT10 to SPDAT17. The content of SPDAT10 is the first to be transferred.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
SPCTL1	SSIG	SPEN	DORD	MSTR	CPOL	CPHA	PSC1	PSC0

 SSIG: /SS1 slave select ignore, if set to '1' MSTR decides whether the device is a master or a slave and /SS1 pin can be used as port pin; if set to '0' and MSTR is set to '1' then the /SS1 pin decides whether the device is a master or a slave

- SPEN: SPI1 enabled '1' or disabled '0', when the SPI1 is disabled the SPI1 pins can be used as general I/O pins
- DORD: SPI1 data order
 - '1' the lsb of the data byte is transmitted first
 - '0' the msb of the data byte is transmitted first
- MSTR: master '1' / slave '0' mode select. See SSIG field
- CPOL: SPI1 clock polarity

'1' – SPICLK1 is high when idle, the leading edge of SPICLK1 is the falling edge and the trailing edge is the rising edge

 $^{\prime 0'}$ – SPICLK1 is low when idle, the leading edge of SPICLK1 is the rising edge and the trailing edge is the falling edge

• CPHA: SPI1 clock phase select

'1' - data is driven on the leading edge of SPICLK1 and is sampled on the trailing edge

'0' - the first data bit is on the MOSI1/MISO1 line before the first SPICLK1 leading edge, data is sampled on the leading edge of SPICLK1 and driven on the trailing edge of SPICLK1

 PSC(1:0): SPI1 clock rate select, determines master clock output. 8 different baud rates can be selected using PSC2[SPSTAT1(5)] & PSC(1:0)[SPCTL1(1:0)]

When working as slave, PSC(2:0) value is not taken into account, nevertheless this value must be different from "000"

PSC(2:0)	clock rate
"000"	fosc/4
"001"	fosc/16
"010"	fosc/64
"011"	fosc/128
"100"	fosc/256
"101"	fosc/512
"110"	fosc/1024



Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
SPSTAT1	SPIF	WCOL	PSC2			N2	N1	N0

- --: Reserved bit
- SPIF: SPI1 transfer completion flag. When ESPI1 (IE2.3) bit and EA (IE.7) bit are '1' and a SPI1 transfer finishes then the SPIF bit is set by hardware and an interrupt is generated. SPIF will also be set when SPI1 is in master mode with SSIG(SPCTL1.7)='0' and /SS(slave select) pin is driven low
- WCOL: SPI1 write collision flag. The WCOL bit is set by hardware when SPDAT is written during a data transfer. WCOL flag is cleared in software by writing '0' to this bit.
- PSC2: SPI1 master clock rate select msb
- N(2:0): Number of additional bytes to transfer in SPI1

The data bytes to be transferred must be written in order. Transfer will start when SPDAT1N is written if N=SPSTAT(2:0).

A write collision will occur if any SPDAT1N (with N<=SPSTAT(2:0)) is written before the 1+N bytes transfer is finished.

Table 8-3. SPI1 port map

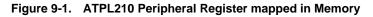
ADD8051C3A Port	Alternate Function
P4.2	/SS1
P4.3	SPICLK1
P4.4	MOSI1
P4.5	MISO1

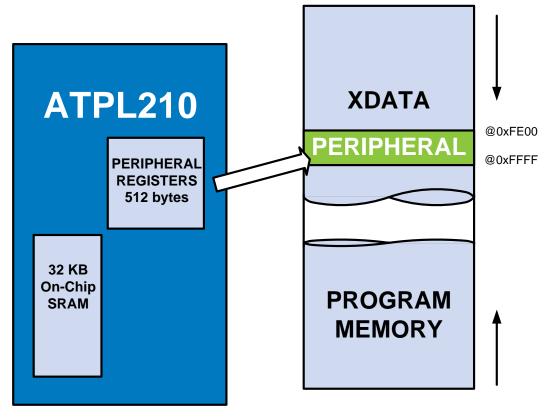


9. Peripheral Registers

A 512 bytes space is reserved on-chip to allocate the system peripheral registers. These registers are mapped in XDATA (External Data Memory space) in addresses from FE00 to FFFF, so these addresses cannot be used to store user data. This fact must be taken into account by the user when calculating the space needed to store the application data.

On-chip Peripheral Registers are mapped in External Data Memory. They can be accessed addressing FE00 to FFFF.



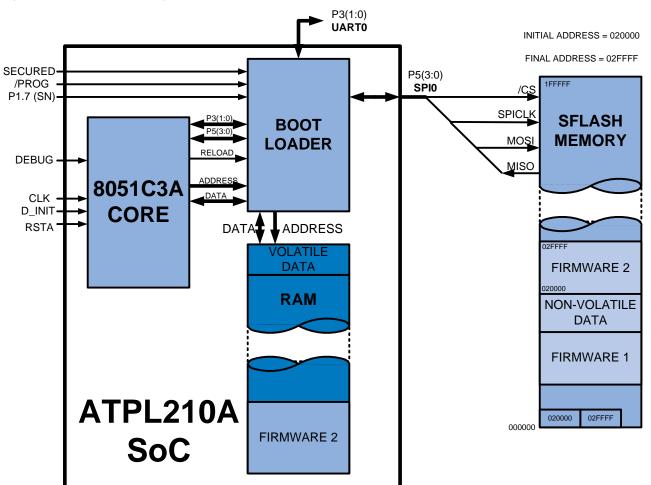




10. Boot Loader

ATPL210A needs an external SPI flash memory to store the firmware and all no-volatile data, while an SRAM (internal or external) is used for firmware execution (once a volatile firmware copy has been automatically generated by the boot loader in the system startup process) and to store all the volatile variables.

After power-up the boot loader forces a start-up cycle uploading the target program from the serial flash to the SRAM, the program is executed from the lower entries of the SRAM and the upper entries are used to store program data. A boot loader diagram is shown in Figure 10-1





Boot loader is executed in the following situations:

- Every time after a Power-On.
- Every time the system is reset by the asynchronous reset signal (RSTA).
- Every time the bit 2 in CONF SFR is set, forcing a reload.
- After a Watchdog timeout, only if bit 1 in CONF SFR is set.

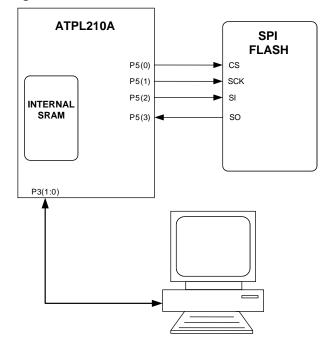
The boot loader circuitry also supports in-system programming. A user can update the firmware via a serial interface without additional logic and without disassembling the system.

When a SPI flash with serial number or/and a Silicon Serial Number IC are available in the system board, the system supports automatic firmware encryption in the flash.

In the ATPL210A the microcontroller can force program reloading by writing in CONF Specific Function Register.



Figure 10-2. Flash connection diagram



10.2 Pin Description

- **/PROG:** After power up, if this signal is tied low the in-system programming mode is enabled, else the system is in execution mode.
 - '0': In-system programming mode
 - '1': Execution mode
- **SECURED:** This pin enables encrypted firmware execution when the board configuration supports it (see Encrypted firmware requirements section).
 - '0': Cryptographic Storage disabled
 - '1': Cryptographic Storage enabled
- **P1.7 (SSN):** Input pin used to read a Serial Number if a valid Silicon Serial Number device is being used (see 10.4.2)

As shown in Figure 10-1 the boot loader acts as a bus switcher. It takes control over the Standard Serial port 0 (P3[1:0]) and pins of port P5 (P5[3:0]) that are used as Serial Peripheral Interface to connect the serial flash (SPI0). See table below:

ADD8051C3A Port	Alternate Function
P3.0	RxD0
P3.1	TxD0
P5.0	/SS0
P5.1	SPICLK0
P5.1	MOSI0
P5.3	MISO0

Table 10-1. SS0 & SPI0 buses



After program booting, the shared pins are fully software controlled. This allows the microcontroller to store no-volatile program data to the serial flash. The microcontroller has full access to all the serial flash content which has to be handled carefully to avoid unexpected overwriting of program code located in the flash memory.

10.3 Flash Programming

10.3.1 In-System programming

The boot loader supports in-system programming for SPI flash memories.

To enable in-system programming /PROG pin must be tied low.

When in-system programming is enabled, the boot loader is controlled via RS-232 commands.

The default RS-232 configuration is 57600 bauds, 8 data bits, 1 stop bit and no parity.

Atmel provides a set of tools that allows customers to carry on in-system programming. For further information, please contact Atmel support.

10.3.2 SPI Flash programming

After a RESET, if /PROG pin is tied low, SPI flash commands are allowed. This set of commands allows the user to send any standard SPI command to the flash memory through the serial port. The boot loader acts as a bridge from serial to SPI protocol.

The serial commands are composed by a header for the boot loader and the SPI command itself.

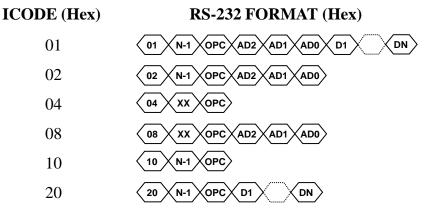
Figure 10-3. Serial commands

MSByte							LSByte
ICODE	BC	OPCODE	ADDR2	ADDR1	ADDR0	DATA1	DATAn
							/

BOOT LOADER HEADER

SPI COMMAND

• **ICODE:** This byte indicates to the boot loader the type of instruction to be executed



• BC: Byte Count indicates the number of data bytes to read or write. The value of this field must be the number of data bytes to read or write minus one, that is, the range of the BC field is [0,255] and the number of data bytes range is [1,256]. If the requested instruction is not a reading or writing instruction the BC field is not used and can take any value.



Note: If Atmel SPI flash memory (see 10.4.3) is being used, and it's configured with a page size of 264 bytes, BC field will have a length of 2 bytes.

ADDR(2:0): 24 bits address field. This field must point to the initial address of the instruction operation.
 DATA(1:n): Data bytes of the SPI instruction (n = BC + 1). The DATA1 byte is the first data byte in a write or read operation.

With the six command types listed in ICODE description, it is possible to execute any standard operation on the SPI flash memory:

Table 10-2. ICODE operations

ICODE (Hex)	Typical Use
01	Page programming
02	Read, Read id.
04	Write enable/disable
08	Erase
10	Read status register
20	Write status register

Example: Programming 256 bytes to flash memory starting at 010000 (Hex) address. (The page programming opcode of the memory is 02 (Hex) and the required erase procedure is not included in this example).

The RS-232 command in hexadecimal format is:

01.FF.02.01.00.00.[byte1 to byte256]

10.4 System startup

When the system is configured in execution mode (/PROG pin is set) and after a power up or reload process, the boot loader checks the system configuration and performs the required operations to ensure the correct execution of the firmware.

Boot loader transfers a volatile copy of the firmware from the flash memory to the SRAM memory starting at address 000000 (Hex). The maximum SPI flash size supported is 16Mbits. The part of the SRAM that is not used to store the firmware is used by the system to store the volatile data in the execution process.

The volatile data is stored in the SRAM memory from top to bottom. This arrangement is done automatically by the system, that is, when the microcontroller wants to save data at address A the system automatically converts the target address to (top address – A). Using this storage method, the amount of volatile data that can be stored in the RAM is determined by the firmware size.

The transfer process can last about 400ms depending on flash device and code size. The first time this process will be longer if auto encryption is activated. Auto encryption must be done once, and **care must be taken about respecting encryption times**.

When the transfer process is finished, the microcontroller begins to execute the program and the system is ready to use. The startup cycle is performed each time the system is powered up or when the microcontroller via software or watchdog forces a firmware reload process writing in a specific register.

Both the location of the firmware at the SPI flash and the transfer size are configurable. The lower six bytes of the SPI flash are used to store the initial and final position of the firmware (from address 0x000000 to 0x000005).

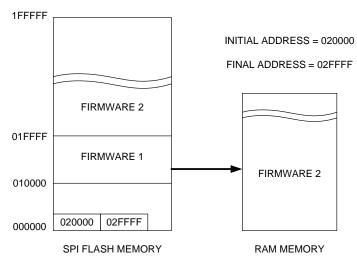


This feature allows working with a SPI flash device that has multiple firmware versions stored in different addresses. The version to be used can be selected by means of its initial and final addresses.

The transfer size is equal to the firmware size.

Bank-switching is supported, so firmware code bigger than 64KB can be managed. Extreme caution must be taken when using bank-switching in order to respect code and variables spaces and avoid overlapping.





10.4.2 Encrypted firmware requirements

In order to run Auto-encryption procedure, the conditions below must be met:

- SECURED = '1'
- DEBUG = '0'
- /PROG = '1'
- Valid Serial Number (see supported devices section)
- Correct flash memory START and END addresses
- Detection by hardware of no-encrypted software stored in a valid flash device
- Starting address of the stored firmware multiple of 4KB (0xXXX000)

First time auto-encryption procedure runs, the code is encrypted and rounded to the next multiple of 4kbytes, and the executions begins.

Next time the system boots, hardware will detect encrypted software in the flash, thus auto encryption procedure will not be started (that is, auto encryption is only necessary to be performed once).

Note1: Auto encryption procedure takes about 30 seconds to be completed, and during encryption no external signals or LEDs are asserted, so care must be taken by the user and encryption time must be respected.

Note2: If SECURED = '0' or some of the other conditions above are not met, the system may run without code encryption. Care must be taken by the user.

Note3: If SECURED = '0' and the code stored in the flash is encrypted (an unusual situation), the system will be expecting no-encrypted code and will fail.

Note 4: A helpful way to check that the firmware in the flash has been encrypted, is to read the start and end addresses allocated at the beginning of the code. If encryption has been done, these address values will make no sense because of the encryption.





10.4.3 Supported Devices

The boot loader supports any SPI flash memory compatible with SPI mode 0 (CLK signal is normally low) and with a value of read instruction operation code equals to 0x03 (Hex). The maximum supported flash size is 16Mbits

Atmel AT45DBxx SPI Flash family is strongly recommended and has been fully tested in ATPL210A development kits and in field configurations.

10.4.3.1 Serial Number Device

If the system is configured to use encrypted firmware it is necessary to include in the system board a device that provides a serial number to the ATPL210A SoC.

Atmel Flash AT45DBxx device also includes its own Serial Number, so the flash device provides the Serial Number itself and there is not necessary an external SN chip.

When a more robust encryption is desired, a SN device can be added to the system, for example:

• DS2401 (Dallas): IC Silicon Serial Number. Provides a valid serial number to be used for encryption purposes.

When both devices are used together (AT45DBxx+DS2401), a more robust encryption is achieved.



10.5 Boot loader registers

10.5.1 SN_BYTE registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SN_BYTE				SN_BYT	E(63:56)				@0xFE4B
				SN_BY	TE(7:0)				@0xFE52
Name:	SN_BYTE								
Address:	0xFE4B – 0	0xFE52							
Access:	Read only								
Reset:	0x00,, 0	x00							

These registers store a 64-bit Silicon Serial Number used in auto-encryption. This serial number is provided by an external device and is only necessary if the system uses auto-encryption features.



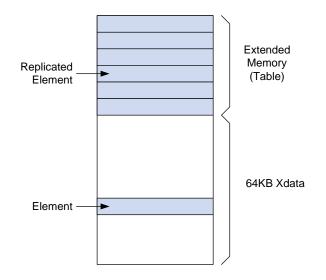
11. Data Memory Extension Block

In ADD8051C3A microcontroller data memory consists of 64 Kbytes of external memory. Due to data table handling in software development for PRIME MAC implementation, it could be useful to extend data memory above 64 Kbytes. For this extension a specific hardware peripheral block has been implemented, with the functionality of address regenerating when needed.

11.1 Description

This block consists of an address extension when address is located within specific limits. The new address is calculated using specific registers, providing capability of replicating a limited space in the data memory so many times as possible with the memory resources and addressing capacity. The space in data memory to be extended is defined by the initial address and the size of this space. This memory space could correspond to an element of a table (data structure), so it would be possible to define a table with a high number of these elements. In order to change from an element to another it is defined the index of the element, used for generating the corresponding address.

Figure 11-1. Basic Memory Extension



In Figure 11-1 it is shown the basic concept about memory extension, which it is thought as an element-replication at high positions of the data memory. The first replicated element starts just after 64 KB from xdata, and the next elements are located at following positions in memory. The element has a configurable size, with a maximum size of 1 KB. If the size of the element is set to 0, then this hardware block is disabled and no extension is applied.

Figure 11-2. Element Example

	0x012B
Element size= 40 bytes	0x012A=Final address 0x0103=Initial address
	0x0102



It is also possible to define several tables, so it is needed to define different initial addresses and sizes, from different spaces (elements of the tables). In this case the extended memory also requires an address offset to change from a table to another.

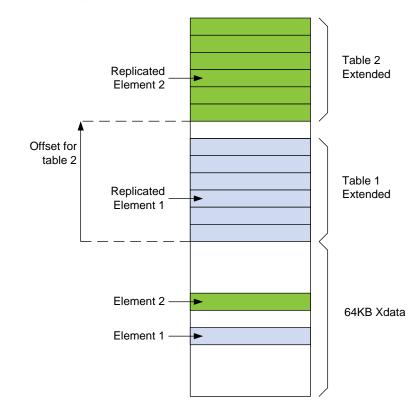


Figure 11-3. Two tables with Memory extension

When using several tables, it is needed to change the offset value, the initial address of the corresponding element, the element size if they are different and the index of the element.

If it is trying to access to the external memory at the corresponding positions to one of this elements, then the hardware block calculates a new address according to its parameters, as follows:

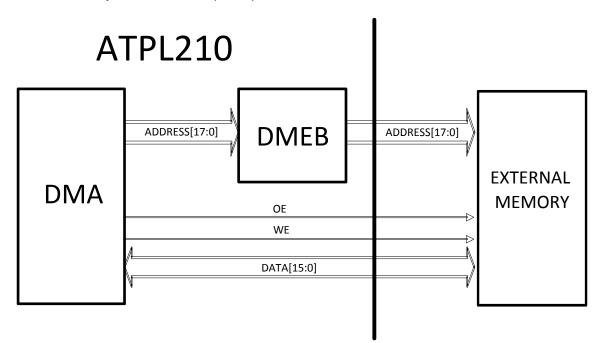
NEW_ADDRESS = 0x10000 + (ADDRESS - INITIAL_ADDRESS) + ELEMENT_SIZE * ELEMENT_INDEX + OFFSET

where ADDRESS is a valid address higher or equal than the initial address of the element and lower than the final address (initial address+element size). INITIAL_ADDRESS is obviously the initial address of the element, ELEMENT_SIZE is the size of the element, ELEMENT_INDEX is its index in the table, OFFSET is the address offset to change among tables, 0x10000 is the first address available above 64 KB from xdata and NEW_ADDRESS is the new address calculated and multiplexed by this block.

The function of this block is to multiplex the address when accessing external data memory, so the rest of ram data and control signals do not take part in data memory extension.



Figure 11-4. Data Memory Extension Block (DMEB) in ATPL210



It is needed to consider that program memory is also included in external memory, so both program and data memories of ADD8051C3A share the total RAM size, thus the data memory extension capability depends also on program memory size, in order to avoid any overwriting between these spaces.

11.2 DMEB Configuration Registers

Table 11-1 shows the configuration peripheral registers involved in Data Memory Extension block:

Table 11-1. Data Memory Extension Block registers address map

Address	Register Name	Description	
0xFE95		See 11.2.2	
0xFE96	TABLE_ELEMENT_SIZE	See 11.2.2	
0xFE97		Sec. 11.0.2	
0xFE98	TABLE_INDEX	See 11.2.3	
0xFE99		See 11.2.4	
0xFE9A	TABLE_ELEMENT_INIT		
0xFEA7			
0xFEA8		See 11.2.5	
0xFEA9	TABLE_OFFSET	See 11.2.5	
0xFEAA			

As explained in Peripheral Registers, Peripheral Registers are mapped in XDATA (External Data Memory space) in addresses from FE00 to FFFF.



11.2.2 TABLE_ELEMENT_SIZE registers

Ν	lame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
TABLE_ELEMENT_SIZE								TABLE_ELEME	@0xFE95	
			TABLE_ELEMENT_SIZE(7:0)							
Name:	TABLE_ELEN	E_ELEMENT_SIZE								
Address:	0xFE95 – 0xFE96									
Access:	Read/write									
Reset:	0x00, 0x00									

• **TABLE_ELEMENT_SIZE:** Size in bytes of the element of the table in extended data memory. The element with this size can be replicated so many times as needed within the available memory. It is important to know the number of elements in the table, in order to verify if the table exceeds memory resources. Note that RAM is used by both program and data memory, therefore carefulness must be taken by the user in order to not overwrite the program memory with the higher elements of the table.



11.2.3 TABLE_INDEX registers

Name		Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
TABLE_INDEX				-	TABLE_IN	@0xFE97				
			TABLE_INDEX(7:0)							@0xFE98
Name:	TABLE_INDEX									
Address:	0xFE97 – 0xFE98									
Access:	Read/write									
Reset:	0x00, 0x0	00								

• **TABLE_INDEX:** Index of the element of the table in extended data memory. The position memory to be read will be:

65536+TABLE_INDEX*TABLE_ELEMENT_SIZE+(ADDRESS-TABLE_ELEMENT_INIT) + TABLE_ OFFSET

where ADDRESS is the position at XDATA memory to be read of the microcontroller.



11.2.4 TABLE_ELEMENT_INIT registers

Na	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
TABLE_ELEMENT_INIT		TABLE_ELEMENT_INIT(15:8)								
				TA	BLE_ELEN	IENT_INIT	(7:0)			@0xFE9A
Name:	TABLE_ELE	EMENT_INIT								

Address: 0xFE99 – 0xFE9A

Access: Read/write

Reset: 0x00, 0x00

• **TABLE_ELEMENT_INIT:** Initial address of the table in data memory. This is the address of the element at the XDATA memory of the microcontroller (64KB-space). Positions from this initial address to TABLE_ELEMENT_INIT+TABLE_ELEMENT_SIZE-1 will never be accessed when data memory extension is used.



11.2.5 TABLE_OFFSET registers

Name	Bit 7 Bit 6 Bit 5	Bit 4 Bit 3 Bit 2 Bit 1 B it 0						
TABLE_OFFSET			@0xFEA7					
		TABLE_OFFSET(20:16)	@0xFEA8					
	TABLE_OFFSET(15:8)							
	TABLE_OFFSET(7:0)							

Name: TABLE_OFFSET

Address: 0xFEA7 – 0xFEAA

Access: Read/write

Reset: 0x00;...; 0x00

• TABLE_OFFSET: 21-bit offset of the table in extended data memory



11.3 Code Example

Following, a brief C code programming example of DMEBwill be explained.

Please note that the total number of elements is not defined. The maximum number of elements is fixed by memory total size. Thus, extra care must be taken by the user in order to not overwrite the program memory with the higher elements of the table.

Pointers must not be used as table elements, to avoid conflict between hardware and compiler.



ONE TABLE

 Table definition (struct) in include *.h Element size = 24 bytes; Uint8 = unsigned integer, 8 bits; Uint16 = unsigned integer, 16bits.

struct ExampleTable { Uint8 Element0; Uint8 Element1; Uint16 Element2; Uint8 Vector[20];

}

...

2. Element declaration as a variable in source file *.c, occupying 24 bytes in XDATA. "_at_" means the allocation of the variable in a concrete position in data memory.

struct ExampleTable TableElement _at_ 0x0200;

3. Initialization of DMEB to use a table with this element replicated out of the XDATA 64KB.

Variables TABLE_ELEMENT_SIZE, TABLE_ELEMENT_INIT and TABLE_INDEX are the DMEB peripheral registers.

```
TABLE_ELEMENT_SIZE = sizeof (TableElement);
TABLE_ELEMENT_INIT = (Uint16)&TableElement;
TABLE_INDEX = 1;
TABLE_OFFSET = 0;
```

4. Table management: reading/writing element 12. Only one DMEB register must be modified when changing between elements, TABLE_INDEX.

```
TABLE_INDEX = 12;

if (TableElement.Element0)

{

TableElement.Element1 = 0x46;

TableElement.Element2 = 0x1FB2;

}

TableElement.Vector[0] = 0x04;

TableElement.Vector[1] = 0x05;
```



1. Another struct is defined. New struct size is 3 bytes.

struct ExampleTableAdditional { Uint8 Element0; Uint16 Element1; }

2. Element declaration in an XDATA address. In this case, both elements are declared in consecutive memory locations. Thus the hardware will know that both elements are together, and it will set the element size equal to the addition of the different type elements.

struct ExampleTable TableElement_at_ 0x0200; struct ExampleTableAdditional TableElementAdditional_at_ 0x0200+ sizeof (TableElement);

3. Initialization of DMEB to use both tables.

TABLE_ELEMENT_SIZE = sizeof (TableElement) + sizeof (TableElementAdditional); TABLE_ELEMENT_INIT = (Uint16)&TableElement; TABLE_INDEX = 1; TABLE_OFFSET = 0;

4. Table management: Reading element 5 from TableElement, writing element to index 10 in TableElementAdditional. Only one DMEB register must be modified when changing between elements, TABLE_INDEX.

TABLE_INDEX = 5; aux = TableElement.Element0; TABLE_INDEX = 10; TableElementAdditional.Element0 = aux; ...



MORE THAN ONE TABLE. TABLES WITH DIFFERENT NUMBER OF ELEMENTS

- _____
 - With the structures defined above, also it is possible to use two tables with different number of elements. In this case, to realize an efficient memory usage, is useful to calculate the memory size of each table, and then work with an offset for each one. If the element differs from one table to another, then it will be necessary to declare variables in different locations and also set the element size (TABLE_ELEMENT_SIZE and TABLE_OFFSET) every time a table change occurs.

struct ExampleTable TableElement _at_ 0x0200; struct ExampleTableAdditional TableElementAdditional _at_ 0x0300;

2. Initialization of one of the tables.

TABLE_ELEMENT_SIZE = sizeof (TableElement); TABLE_ELEMENT_INIT = (Uint16)&TableElement; TABLE_INDEX = 1; TABLE_OFFSET = 0;

3. Table management: reading element 4 in TableElement and writing element 7 in TableElementAdditional. In this case, every time a table change occurs, the following values must be changed: element size, element initial address, table offset, element index.

TABLE_INDEX = 4; aux = TableElement.Element0; TABLE_ELEMENT_SIZE = sizeof (TableElementAdditional); TABLE_ELEMENT_INIT = (Uint16)&TableElementAdditional; TABLE_OFFSET = numElements * sizeof (TableElement); TABLE_INDEX = 7 TableElementAdditional.Element0 = aux;



12. Direct Memory Access (DMA)

ATPL210A SoC includes 4 DMA (Direct Memory Access) channels which provide an enhanced way for memory transference between positions from the same data memory location or between peripherals and the external data memory. The main goal is to increase memory transference performance, comparing with ADD8051C3A microcontroller using MOVX instructions. Two of these DMA channels are used by the physical layer and are not visible by the microcontroller or another interface. These two channels (PHY DMA channels) have a similar structure to the other two channels (generic DMA channels), but they are not configurable, except addresses –destination in reception and source in transmission-.

The other two DMA channels are generic and have some configurable parameters: source and destination addresses, source and destination sizes, time of each DMA access and time between consecutive accesses

12.1 DMA Management

DMA channels are intended to be used by means of a user interface.

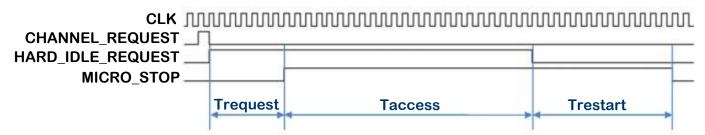
By using a software interface, the user can perform DMA write/read operations and DMA configuration, so low-level management becomes transparent. Please see "PRIME User Help" for further information about software management.

12.2 DMA Hardware

DMA requires to access data memory, so the microcontroller execution must be stopped. It is needed a Hard Idle ReQuest (HIRQ) to the microcontroller, in order to save the current state with the current instruction complete. There is a variable time from the DMA request to the micro stop, depending on its current execution state. When the microcontroller is correctly stopped a signal is provided to DMA to get full access to data memory. When transference is completed, the hard idle request is cleared and the microcontroller shall resume its execution. The signal provided by the micro to indicate the stop will be cleared too.

Figure 12-1 shows a diagram with the main DMA control signals and DMA times.

Figure 12-1. DMA control signals



- **CHANNEL_REQUEST:** One-cycle request signal from the current DMA channel. This signal can be generated simultaneously by several channels but the memory access will be assigned to the most prioritary one.
- **HARD_IDLE_REQUEST:** Request signal for a hard-idle state of the micro. It is a request to stop current micro execution to get access to the external memory, and it must be activated during the complete access.
- **MICRO_STOP:** Signal to indicate that the microcontroller is stopped. It is required by DMA to know if it is possible to access to the memory.



Table 12-1. DMA timing values

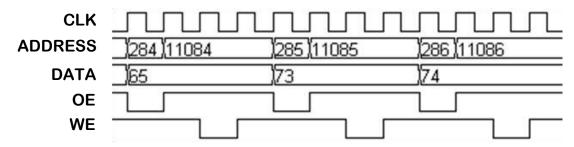
Time	Number of clock cycles
Trequest	1 - 18
Taccess	4*Nbytes
Trestart	1 - 9

- **Trequest:** Time from hard idle request to micro stop. After a channel request, the hard idle request signal has to keep activated till the end of the DMA access. DMA will get access to memory when the stop signal from micro is also activated. This time is variable from one to six machine cycles, depending on instruction execution.
- **Taccess:** Effective access time. This is the total time available for DMA to get full access to the external memory. When access ends, the hard idle request signal has to be deactivated, in order to restart micro execution.
- **Trestart:** Time of micro restart, from the deactivation of the hard idle request to the reactivation of micro execution.

When in DMA mode, only one clock cycle is used to read a byte from memory, and three clock cycles are used to write a byte; thus, a byte-transference consists of 4 clock cycles.

An example is shown in Figure 12-2.





12.3 DMA Channel Priority

Simultaneous DMA channels activation is possible, so it is needed to establish a priority for each DMA channel. It is possible for two or more channels to request memory access at the same time (low probability), so the channel with the highest priority will get the requested access. Because of hardware requirements, the priority of the channels has been fixed to optimize the physical layer operation. These are the priorities assigned:

Table 12-2. DMA channels priority

DMA Channel	Priority
PHY_RX	3
PHY_TX	2
GEN_1	1
GEN_2	0



Channel priority increases with "priority number".

DMA channel used in physical reception (PHY_RX) is the most critical, so it has the highest priority, followed by physical transmission DMA channel.

Generic DMA channels (GEN_1 and GEN_2) are identical, except by their priority, and both have lower priority than physical ones.

12.4 DMA Transfer Capabilities

The maximum amount of data which a DMA channel is able to move is 64 KB, corresponding to external data memory size of the ADD8051C3A.

The maximum data size moved in a single DMA access is 256 bytes, and the minimum is 1 byte.

When a DMA channel is active, it will access so many times as necessary to complete the transference, requesting accesses and microcontroller stops.

Transfer sizes are distributed in packets of 16 bytes, in accordance with AES128 data size (see PRIME specification). This is the reference size, and all DMA accesses are configured according to this 16-byte pattern.

Consequently all time accesses are established as well, owing to byte-transference duration (4 clock cycles). "Time access" or "time between accesses" are then referred as a number of these packets. The duration for this packet is 16 bytes * 4 clk/byte * 50 ns/clk = $3.2 \mu s$.

If the DMA access requirement is less than 16 bytes, then access ends when transfer is complete.

For physical DMA channels these times are established by real-time memory transfer requirements. Every physical DMA access for these channels is 16-byte long. Time between consecutive accesses is set up by symbol data of the physical layer when needed.

For generic DMA channels all times are fixed by specific registers. It is needed to consider for each DMA access an additional time to the total time access, which may be variable from 1 clock cycle to 18 clock cycles, depending on ADD8051C3A instruction execution (this variable time differs if the microcontroller is returning from an interruption subroutine or it is returning from Idle mode or Power Down mode).

12.5 DMA Interrupts

In order to inform the ADD8051C3A microcontroller about the completion of a DMA transference, two interrupts are generated by hardware when the transference or a part of it is completed. One of these interrupts is related to the physical layer, and it is activated in both emission and reception. The other interrupt is related to generic channels, and the activation is controlled by specific registers.

Interrupt type	Internal Bit	External Interrupt (ADD8051C3A port)		
Physical	INT_PHY PHY_SFR(0)	INT1 (P33)		
Generic	INT_DMA PHY_SFR(1)	INT0 (P32)		

See PHY_SFR Register in DMA Configuration Registers section.

The physical interruption field is the PHY_SFR(0) bit, named INT_PHY, and it is internally connected to the external interrupt INT1 of the ADD8051C3A microcontroller, which corresponds with P33 port.

The generic channel interruption is also located at the PHY_SFR register.



This interruption field is the PHY_SFR(1) bit, named INT_DMA, and it is internally connected to the external interrupt INT0 of the ADD8051C3A microcontroller, which corresponds to P32 port.

Both interrupts are "low-level active" and are activated by level. When complete or partial transference is finished then the bit will be set to '0' by hardware, generating the corresponding interruption. It the interrupt is active (bit EA -IE(7)– is set to '1' for all interrupts, and bit EX0 -IE(0)– is set to '1' for external interrupt 0 and bit EX1 -IE(2)– is set to '1' for external interrupt 1, see ADD8051 Specification Interrupt Enabling), then this bit must be cleared by writing '1' at the PHY_SFR bit in the interrupt subroutine.

12.6 DMA Flags

There are several flags activated simultaneously with DMA interruptions, that allow to know the type of DMA channel used in the transference, and if this transference is partial or complete. These flags are located at the DMA_SFR register, mapped at address 0xFE2B of the data memory.

See DMA_SFR Register in DMA Configuration Registers section

12.7 Generic DMA channels

Generic DMA channels are fully configurable with specific registers. Physical DMA channels can be understood as a particular case of generic channels, with some of their parameters fixed by physical requirements. There are 7 configurable parameters and 6 registers for each generic channel. Both generic channels are identical except in:

- The priority, which is higher for the generic 1.
- The flags activated at interruption.

The priority for both channels is lower than priority for physical channels.

The registers involved in DMA channels operation are described in Peripheral registers section.

ADDR_DMA1_SRC	ADDR_DMA2_SRC
ADDR_DMA1_DTN	ADDR_DMA2_DTN
SIZE_DMA1_SRC	SIZE_DMA2_SRC
SIZE_DMA1_DTN	SIZE_DMA2_DTN
SIZE_RQ_DMA1	SIZE_RQ_DMA2
DMA1_CHN_CTL	DMA2_CHN_CTL

As an example, below it is shown a memory-transference with channel 2, by establishing source and destination sizes equal to 48 bytes, the same as the request size, the number of packets for DMA access and between accesses equal to 15. It is possible to see the activation of DS_CHN2, DR_CHN2 and the interruption INT_DMA and memory control signals.

Figure 12-3. DMA channel 2 process

CLK	
CHANNEL_RQ	J
HARD_IDLE_RQ	
MICRO_STOP	
INT_DMA	
DS_CHN2	
DR_CHN2	
ADDRESS	
DATA	10111)-64 Jac 173 174 J20 Jac Jac 172 Jac 172 Jac 173 126 Jac
OE	
WE	



12.8 Physical DMA channels

Physical DMA channels can be understood as a particular case of generic channels in which some parameters are fixed. The main differences between PHY_RX (Reception physical DMA channel) and PHY_TX (Transmission physical DMA channel) channels are the direction of the transference, the interrupts generated, and the flags activated. Note that these DMA channels operate against internal registers mapped in data memory, so one of the buffers is always internal.

12.9 PHY_RX DMA channel

This is a dedicated DMA channel for physical data reception. The source is a 16-byte width buffer at the end of the reception branch (see PHY layer section), and the destination buffer is the data memory.

Size of the destination buffer is variable and corresponds to PPDU size (DMA destination size = PPDU length). Time between accesses depends on symbol duration, modulation schemes and physical layer processing capability.

There are two interrupts generated at reception, one of them when the physical header is correctly received (two first symbols), and the other one when the message is completely received. Then it is easy to see that DMA request size in this case correspond to the physical header size (11 bytes), which causes an interruption when the current transference size is equal to PHY header size and also activates an internal flag. For the other interruption another internal flag is activated to indicate the completion of reception.

The priority for this channel is the maximum available, because of critical reception requirements: physical data must be moved as fast as possible after symbol processing.

12.10 PHY_TX DMA channel

This is a dedicated DMA channel for physical data transmission. The source buffer is data memory and the destination is a 16-byte width buffer at the beginning of the transmission branch.

Size of the source buffer is also variable and corresponds to PPDU size (DMA source size = PPDU length). Time between accesses NPUM also depends on physical requirements.

For this channel there is only one interrupt available, at the end of the transmission, so it is easy to see that DMA request size in this case is equal to DMA source size. Note that this is the only DMA channel in which the end of memory transference is caused by the completion of the source buffer size, instead of the rest of the channels, in which the end of DMA is provoked by the destination buffer size. At the end of transmission, a physical interruption is generated and an internal flag is set.

The priority for this channel is lower than PHY_RX channel, but higher than generic channels, due to transmission requirements: physical data must be prepared before symbol processing.





12.11 DMA Configuration Registers

12.11.1 PHY_SFR Register

Name	Register Bit 7	Bit 6	Bit 5	Bit 4	Dit 2	Bit 2	Bit 1	B it 0
IY_SFR	BCH_ERR	CD	UMD	ыц 4 	Bit 3	TXRX	INT_DMA	INT_PHY
ne:	PHY_SFR	0.5	OND	1		17404		
dress:	0xFE2A							
cess:	Read/write							
set:	0x87							
	UNUT							
•	:	Reserved	bit					
٠	BCH_ERR:		nel Error Fla	ag.				
			cleared to on instant. O	-		-	esence of a	n OFDM sig
			used for ret E specification	-	sult of "Busy	Channel" i	in the PHY_I	DATA.confirr
٠	CD:	Carrier De	tect bit.					
		This bit is a whole rece	-	hardware wl	hen an OFD	M signal is	detected, an	d it is active
		This bit is u	used in chan	nel access (CSMA-CA a	lgorithm) fo	or performing	channel-sen
٠	UMD:	Unsupport	ed Modulatio	on Scheme	flag.			
		-		-	-		th correct CF	
٠	TXRX:	Transmiss	ion order.					
		at TX_TIM		nd then the	emission lev	el is specifi	data memory ed at ATTEN	
		If this bit is '1'.	read, only r	eturns '0' wl	nen physical	transmissio	on has starte	d. Otherwise
		The transm	nission will b	egin when T	IMER_BEAG	CON_REF i	s equal to TX	CTIME.
٠	INT_DMA:	DMA interr	ruption					
			nternally cor esponds with		ne ADD8051	C3A microc	controller exte	ernal interrup
		It is low-le		is set to '0	' by DMA ha	ardware an	d is cleared	by writing '1
٠	INT_PHY:	Physical L	ayer interrup	otion				
			nternally cor esponds with		ne ADD8051	C3A microc	controller exte	ernal interrup
		It is low-le PHY_SFR		t is set to 'C)' by physica	al layer and	d is cleared	by writing '1
		PHY_SFR	(0).					



12.11.2 DMA_SFR Register

	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0				
	DMA_SFR	DR_CHN2	DS_CHN2	DR_CHN1	DS_CHN1	DR_TX	DS_TX	DR_RX	DS_RX				
I	Name:	DMA_SFR											
	Address:	0xFE2B											
	Access:	Read/write											
I	Reset:	0x04											
	٠	DR_CHN2:	Data Request flag for generic DMA channel 2										
			This flag is a	This flag is activated when the requested size of the destination buffer is already transferred.									
			The request	The requested size is specified at SIZE_RQ_DMA2.									
	٠	DS_CHN2:	Data Size fla	ag for gener	ic DMA chan	nel 2							
			This flag is a	activated whe	en the compl	ete size of t	he destinatio	on buffer is a	lready transferred.				
			The total tra	nsference si	ze is specifie	d at SIZE_E	DMA2_DTN.						
			The source	buffer is data	a memory at	ADDR_DM/	A2_SRC.						
			The destinat	tion buffer is	data memor	y at ADDR_	DMA2_DTN						
	٠	DR_CHN1:	Data Reque	st flag for ge	eneric DMA o	channel 1							
			This flag is a	activated whe	en the reques	sted size of	the destinati	on buffer is a	already transferred.				
			The request	ed size is sp	ecified at SIZ	E_RQ_DM	A1.						
	٠	DS_CHN1:	Data Size fl										
			This flag is a	activated whe	en the compl	ete size of t	he destinatio	n buffer is a	lready transferred.				
			The total tra	nsference si	ze is specifie	d at SIZE_E	DMA1_DTN.						
			The source	buffer is data	a memory at	ADDR_DM/	A1_SRC.						
					data memor		DMA1_DTN						
	٠	DR_TX:	•	•	HY_TX DMA								
				•		_	-	-	d at the end of the when the header is				
			received and	d the other o	ne when the	complete m	lessage is re	ceive. In tra	nsmission there is no				
			•		ader is sent)								
	٠	DS_TX:		-	TX DMA cha								
			-		en the compl				ferred.				
			That means	that a comp	lete message	e has been o	correctly tran	ismitted.					
					a memory at		′_INI_TX.						
			The destinat	tion buffer is	PHY_TX buf	fer.							
					the complete		Y Protocol D	ata Unit) len	igth.				
	٠	DR_RX:		-	HY_RX DMA								
			-						already transferred.				
			That means	that an inco	ming messag	ge physical	header has b	been correct	ly received.				

The requested size is fixed to 11 bytes, in accordance with the physical layer header.



• **DS_RX:** Data Size flag for PHY_RX DMA channel

This flag is activated when the complete size of the destination buffer is transferred.

That means that a complete message has been correctly received.

The source buffer is PHY_RX and destination buffer is data memory at ADDR_PHY_INI_RX.

The transference size is the complete PPDU (PHY protocol Data Unit) length.





12.11.3 PHY_TX Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
PHY_TX	PHY_TX(127:120)									
	PHY_TX(7: 0)									
Name:	PHY_TX									
Address:	0xFE00 – 0	xFE0F								
Access:	Read only									

Reset: 0x00, ..., 0x00;

• **PHY_TX:** Physical layer input buffer for transmission. The buffer size is 16 bytes in order to fit in with standard AES128 data size.

Specific transfer from SRAM to this buffer is executed by DMA when necessary in transmission. DMA TX channel is dedicated for this type of transference.

DMA source address	=>	ADDR_PHY_INI_TX
DMA destination address	=>	PHY_TX(0)
DMA source size	=>	variable
DMA destination size	=>	16



12.11.4 PHY_RX Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0			
PHY_RX	PHY_RX(127:120)										
				PHY_R	RX(7: 0)				@FE1F		
Name:	PHY_RX										
Address:	0xFE10 - 0	xFE1F									
Access:	Read only										

Reset: 0x00, ..., 0x00

• **PHY_RX:** Physical layer output buffer for reception. The buffer size is 16 bytes in order to fit in with standard AES128 data size.

Specific transfer from SRAM to this buffer is executed by DMA when necessary in reception. DMA RX channel is dedicated for this type of transference.

DMA source address	=>	PHY_RX(0)
DMA destination address	=>	ADDR_PHY_INI_RX
DMA source size	=>	16
DMA destination size	=>	variable



12.11.5 ADDR_PHY_INI_RX Registers

Nar	ne	e Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 Bit 0								
ADDR_PH	Y_INI_RX			A	DDR_PHY_	INI_RX(15:8)			@0xFE20
		ADDR_PHY_INI_RX(7:0)								
Name:	ADDR_P	HY_INI_R>	K							
Address:	0xFE20 -	- 0xFE21								
Access:	Read/write									
Reset:	0x02, 0x10									
•	ADDR_P	HY_INI_R	K: I	Data memo	ory address	for reception	n			
					destination n the physic		o which	DMA RX	channel	initiates data



12.11.6 ADDR_PHY_INI_TX Registers

Nar	ne Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 Bit 0									
ADDR_PH	Y_INI_TX			A	DDR_PHY_	INI_TX(15:8	3)			@0xFE22
		ADDR_PHY_INI_TX(7:0)								
Name:	ADDR_PHY_INI_TX									
Address:	0xFE22 – 0xFE23									
Access:	Read/write									
Reset:	0x00, 0x10									
•	ADDR_P	HY_INI_T	K: [Data memo	ry address	for transmi	ssion			
					source a a source a	ddress fron Iayer	n which I	DMA TX c	hannel ini	tiates data



12.11.7 ADDR_DMA1_SRC Registers

•

Nar	ne	Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 B it 0								
ADDR_DN	IA1_SRC			A	DDR_DMA	1_SRC(15:8	5)			@0xFEE0
		ADDR_DMA1_SRC (7:0)								@0xFEE1
Name:	ADDR_D	MA1_SRC								
Address:	0xFEE0 – 0xFEE1									
Access:	Read/write									
Reset:	0x03, 0x0	00								

ADDR_DMA1_SRC: The 16-bit value stored in these registers is the source address for DMA channel 1. Start address at data memory from which source data is copied to the destination buffer. DMA begins reading from this address and increments it one by one, after each byte reading



12.11.8 ADDR_DMA1_DTN Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ADDR_DN	IA1_DTN			A	DDR_DMA [,]	1_DTN(15:8	5)			@0xFEE2
				A	DDR_DMA	1_DTN (7:0)			@0xFEE3
Name:	ADDR_D	MA1_DTN								
Address:	0xFEE2 -	– 0xFEE3								
Access:	Read/wri	te								
Reset:	0x05, 0x0	00								

ADDR_DMA1_DTN: The 16-bit value stored in these registers is the destination address for DMA channel 1.

Start address at data memory to which data is copied from the source buffer. DMA begins writing to this address and increments it one by one, after each byte-writing.



12.11.9 SIZE_DMA1_SRC Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SIZE_DM	A1_SRC			S	SIZE_DMA1	_SRC(15:8)				@0xFEE4
				5	SIZE_DMA1	_SRC (7:0)				@0xFEE5
Name:	SIZE_DM	IA1_SRC								
Address:	0xFEE4 -	- 0xFEE5								
Access:	Read/writ	te								
Reset:	0x00, 0x2	10								

• SIZE_DMA1_SRC: The 16-bit value stored in these registers is the source size for DMA channel 1

Note that source buffer size and destination buffer size can be different:

- If source buffer size is smaller than destination buffer size, then DMA will restart the source address at ADDR_DMA1_SRC when source buffer size is reached, in order to full fill the destination buffer with data from the source buffer.
- If source buffer size is greater than destination buffer size, then only SIZE_DMA1_DTN bytes will be copied and the rest will be ignored.

Nevertheless, in normal operation both source and destination buffers have the same size.



12.11.10 SIZE_DMA1_DTN Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SIZE_DM	A1_DTN			S	SIZE_DMA1	_DTN(15:8))			@0xFEE6
				:	SIZE_DMA1	_DTN (7:0)				@0xFEE7
Name:	SIZE_DN	IA1_DTN								
Address:	0xFEE6 -	– 0xFEE7								
Access:	Read/wri	te								
Reset:	0x01, 0x0	00								

• **SIZE_DMA1_DTN:** The 16-bit value stored in these registers is the destination size for DMA channel 1 This is the required size to complete the DMA transference. The channel will be active until this size is transferred.

The completion of this transference causes an interrupt by clearing to '0' the INT_DMA bit (PHY_SFR(1)), and the activation of the DS_CHN1 flag to '1' (when using GEN_1 channel) or DS_CHN2 (when using GEN_2 channel)





12.11.11 SIZE_RQ_DMA1 Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SIZE_RC	2_DMA1				SIZE_RQ_I	DMA1(15:8)				@0xFEE8
					SIZE_RQ_	DMA1(7:0)				@0xFEE9
Name:	SIZE_RC	2_DMA1								
Address:	0xFEE8 -	– 0xFEE9								
Access:	Read/wri	te								
Reset:	0x01, 0x0	00								
٠	SIZE_RC	Q_DMA1:	The 16-b	it value sto	ored in thes	e registers	is the requ	est size fo	r DMA cha	innel 1
						•				to detect a an interrupt

partial completion of the transference. The completion of this size causes an interrupt by clearing to '0' the INT_DMA bit (PHY_SFR(1)) and the activation of DR_CHN1 flag to '1' (when using GEN_1) or DR_CHN2 (when using GEN_2).

Note that one memory transfer can cause two interrupts at different time. If SIZE_RQ_DMA1 is equal to SIZE_DMA1_DTN, then only one interrupt will be generated and both DS_CHN1 and DR_CHN1 flags will be activated when using GEN_1 (DS_CHN2 and DR_CHN2 when using GEN_2)



12.11.12 DMA1_CHN_CTL Register

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0				
DMA1_C	HN_CTL		NPDM	A(3:0)			NPUN	NPUM(3:0)					
Name:	DMA1_C	HN_CTL											
Address:	0xFEEA												
Access:	Read/wri	te											
Reset:	0x00												

This register is used in DMA channel 1 control. It is possible to configure the time of each access or the time between consecutive accesses. These times are referred to 16-bytes packets, with a duration of 3.2μ s.

• NPDMA: Number of Packets for DMA

Time (in 3.2μ s steps) of every DMA access when the DMA channel 1 is active. Nevertheless, the real time spent depends on the microcontroller execution.

When DMA channel 1 is not active, this value is cleared to "0". Modifying this value causes the activation of the channel. If NPDMA>0 then the channel will be activated. When the channel completes the transference, NPDMA value is cleared by hardware (NPDMA=0)

• **NPUM:** Number of Packets for UM (microcontroller).

This field is used to set the time (in $3.2\mu s$ steps) between DMA accesses when the DMA channel 1 is active.

This time cannot be null, so the number of packets used is NPUM+1.



12.11.13 ADDR_DMA2_SRC Registers

٠

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ADDR_DN	IA2_SRC			A	DDR_DMA	2_SRC(15:8	3)			@0xFEEB
				A	DDR_DMA	2_SRC (7:0)			@0xFEEC
Name:	ADDR_D	MA2_SRC								
Address:	0xFEEB	– 0xFEEC								
Access:	Read/wri	te								
Reset:	0x07, 0x0	00								

ADDR_DMA2_SRC: The 16-bit value stored in these registers is the source address for DMA channel 2. Start address at data memory from which source data is copied to the destination buffer. DMA begins reading from this address and increments it one by one, after each byte reading



12.11.14 ADDR_DMA2_DTN Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ADDR_DN	IA2_DTN			А	DDR_DMA	2_DTN(15:8	3)			@0xFEED
				A	DDR_DMA	2_DTN (7:0)			@0xFEEE
Name:	ADDR_D	MA2_DTN								
Address:	0xFEED	– 0xFEEE								
Access:	Read/wri	te								
Reset:	0x09, 0x0	00								

ADDR_DMA2_DTN: The 16-bit value stored in these registers is the destination address for DMA channel 2.

Start address at data memory to which data is copied from the source buffer. DMA begins writing to this address and increments it one by one, after each byte-writing.



12.11.15 SIZE_DMA2_SRC Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SIZE_DM	A2_SRC			S	IZE_DMA2	_SRC(15:8)				@0xFEEF
				S	SIZE_DMA2	2_SRC (7:0)				@0xFEF0
Name:	SIZE_DM	IA2_SRC								
Address:	0xFEEF -	– 0xFEF0								
Access:	Read/writ	te								
Reset:	0x00, 0x ²	10								

• SIZE_DMA2_SRC: The 16-bit value stored in these registers is the source size for DMA channel 2

Note that source buffer size and destination buffer size can be different:

- If source buffer size is smaller than destination buffer size, then DMA will restart the source address at ADDR_DMA2_SRC when source buffer size is reached, in order to full fill the destination buffer with data from the source buffer.
- If source buffer size is greater than destination buffer size, then only SIZE_DMA2_DTN bytes will be copied and the rest will be ignored.

Nevertheless, in normal operation both source and destination buffers have the same size.



12.11.16 SIZE_DMA2_DTN Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SIZE_DM	A2_DTN			S	SIZE_DMA2	_DTN(15:8)				@0xFEF1
				:	SIZE_DMA2	_DTN (7:0)				@0xFEF2
Name:	SIZE_DM	1A2_DTN								
Address:	0xFEF1 -	- 0xFEF2								
Access:	Read/writ	te								
Reset:	0x01, 0x0	00								

• **SIZE_DMA2_DTN:** The 16-bit value stored in these registers is the destination size for DMA channel 2 This is the required size to complete the DMA transference. The channel will be active until this size is transferred.

The completion of this transference causes an interrupt by clearing to '0' the INT_DMA bit (PHY_SFR(1)), and the activation of the DS_CHN1 flag to '1' (when using GEN_1 channel) or DS_CHN2 (when using GEN_2 channel)



12.11.17 SIZE_RQ_DMA2 Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SIZE_RC	_DMA2				SIZE_RQ_E	DMA2(15:8)				@0xFEF3
					SIZE_RQ_	DMA2(7:0)				@0xFEF4
Name:	SIZE_RC	2_DMA2								
Address:	0xFEF3 -	- 0xFEF4								
Access:	Read/writ	te								
Reset:	0x01, 0x0	00								
٠	SIZE_RC	Q_DMA2:	The 16-b	it value sto	red in these	e registers	is the requ	est size for	r DMA cha	nnel 2
			This size	has to he s	maller or e	aual than S		DTN and	haeu ei ti k	to detect a

This size has to be smaller or equal than SIZE_DMA2_DTN, and it is used to detect a partial completion of the transference. The completion of this size causes an interrupt by clearing to '0' the INT_DMA bit (PHY_SFR(1)) and the activation of DR_CHN1 flag to '1' (when using GEN_1) or DR_CHN2 (when using GEN_2).

Note that one memory transfer can cause two interrupts at different time. If SIZE_RQ_DMA2 is equal to SIZE_DMA1_DTN, then only one interrupt will be generated and both DS_CHN1 and DR_CHN1 flags will be activated when using GEN_1 (DS_CHN2 and DR_CHN2 when using GEN_2)



AMEL

12.11.18 DMA2_CHN_CTL Register

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
DMA2_C	HN_CTL		NPDMA	A(3:0)			NPUN	/(3:0)	
Name:	DMA2_C	HN_CTL							
Address:	0xFEF5								
Access:	Read/wri	te							
Reset:	0x00								

This register is used in DMA channel 2 control. It is possible to configure the time of each access or the time between consecutive accesses. These times are referred to 16-bytes packets, with a duration of 3.2μ s.

• NPDMA: Number of Packets for DMA

Time (in 3.2μ s steps) of every DMA access when the DMA channel 2 is active. Nevertheless, the real time spent depends on the microcontroller execution.

When DMA channel 2 is not active, this value is cleared to "0". Modifying this value causes the activation of the channel. If NPDMA>0 then the channel will be activated. When the channel completes the transference, NPDMA value is cleared by hardware (NPDMA=0)

• **NPUM:** Number of Packets for UM (microcontroller).

This field is used to set the time (in $3.2\mu s$ steps) between DMA accesses when the DMA channel 2 is active.

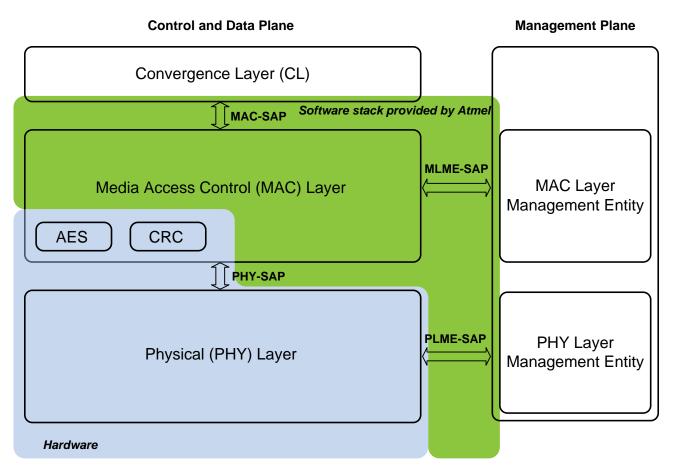
This time cannot be null, so the number of packets used is NPUM+1.



13. ATPL210 MAC Coprocessor

The ATPL210 hardware MAC layer consists of a hardware implementation of some functionalities of the MAC Layer Entity specified in PRIME specification. These features are CRC calculation and AES128 block. So it is possible to consider this hardware as a MAC coprocessor to improve the ADD8051C3A performance.





Atmel PRIME stack implements by software the rest of the MAC layer requirements and capabilities. Furthermore, the software package allows the communication with the Management Plane by means of the two Access points described by PRIME (PHY Layer Management Entity SAP and MAC Layer Management Entity SAP) and the interface to communicate MAC layer with the upper layer (Convergence Layer).

Please check the "Atmel PRIME Stack User Manual" for software package detailed description and functionality.



13.2 Cyclic Redundancy Check (CRC)

There are three types of MAC PDUs (generic, promotion and beacon) for different purposes, and each one has its own specific CRC. In ATPL210A there is a hardware implementation of every CRC type calculated by the MAC layer. This CRC hardware-calculation is enabled by default. Note that the CRC included at the physical layer is also a hardware implementation available in ATPL210A and it is also enabled by default.

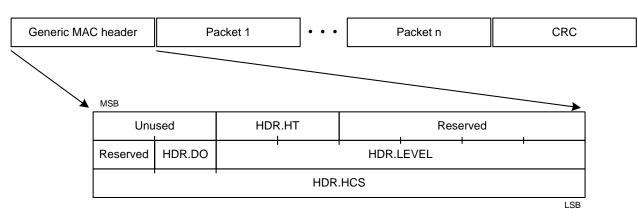


Figure 13-2. Generic MAC PDU format and generic MAC header detail

In transmission all CRC bytes are real-time calculated and the last bytes of the MAC PDU are overwritten with these values, (provided that the field HT in the first byte of the MAC header in transmission data is equal to the corresponding MAC PDU type).

In reception the CRC bytes are also real-time calculated and these bytes are checked with the last bytes of the MAC PDU. If the CRC is not correct, then an error flag is activated, the complete frame is discarded, and the corresponding error counter is increased. These counters allow the MAC layer to take decisions according to error ratio.

For the Generic MAC PDU, there is an 8-bit CRC in the Generic MAC header, which corresponds to PRIME HDR.HCS. In reception if this CRC doesn't check successfully, the current frame is discarded and no interruption is generated.

This works in the same way as CRC for the PHY layer (CRC Ctrl, located in the PHY header, see PRIME specification for further information).

There is another CRC for the Generic MAC PDU which is the last field of the GPDU. It is 32 bits long and it is used to detect transmission errors. The CRC shall cover the concatenation of the SNA with the GPDU except for the CRC field itself. In reception, if the CRC is not successful then an internal flag is set and the error counter is increased.

For the Promotion Needed PDU there is an 8-bit CRC, calculated with the first 13 bytes of the header. In reception, if this CRC is not correct, then an internal flag is set and the corresponding error counter is increased.

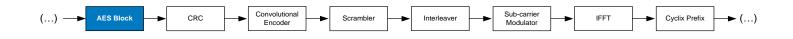
For the Beacon PDU there is a 32-bit CRC calculated with the same algorithm as the one defined for the CRC of the Generic MAC PDU. This CRC shall be calculated over the complete BPDU except for the CRC field itself. In reception, if this CRC is not successful, then an internal flag is set and the same error counter as for GPDU is increased. The hardware used for this CRC is the same as the one used for GPDU.



13.3 Advanced Encryption Standard

One of the security functionalities in PRIME is the 128-bit AES encryption of data and its associated CRC. ATPL210A includes a hardware implementation of this block, and it is used by the physical layer in real-time transmission/reception. It is possible to use this block externally as a peripheral unit, by accessing the specific registers designed to control it. Therefore there are some configurable parameters and input/output buffers to the block.

Figure 13-3. PHY Layer transmitter block diagram



There are two basic operation ways in ATPL210A when using PRIME Security Profile 1. The first one is real-time encryption and the second one is independent encryption from the PHY layer.

Real-Time Encryption: the AES128 core is integrated in the physical chain, and data is encrypted and decrypted in real-time when needed. In transmission, data is transferred to the emission buffer by means of the DMA TX channel. Then the 128 bits located in the buffer are encrypted before starting transmission (Note that Beacon PDU, Promotion PDU and Generic MAC header, as well as several control packets, are not encrypted). Data is extracted when required from this buffer until it is empty, and then a new DMA transfer is requested to fill the 16 bytes and a new encryption is executed. The key used for encryption must be set at the corresponding register, and it can vary from a packet to another.

In reception, data is obtained from the PHY layer and it is passed to the AES128 block. When the reception buffer is full with incoming data, the 128 bits are decrypted and transferred to external memory through DMA RX channel. Then the reception buffer is available again to fill with processed data.

The header is always real-time analyzed in order to know if encryption process must be applied.

Independent Encryption: the AES128 core is used as a peripheral unit, accessible with several registers mapped in external memory. In this mode, when in transmission, data must be encrypted previously to the use of the PHY_DATA.request primitive (see PRIME specification), in an independent way. In reception, data passed by the PHY layer is already encrypted and must be decrypted in a subsequent process.

When working with AES block as a peripheral unit, automatic CRC calculation by hardware is disabled.

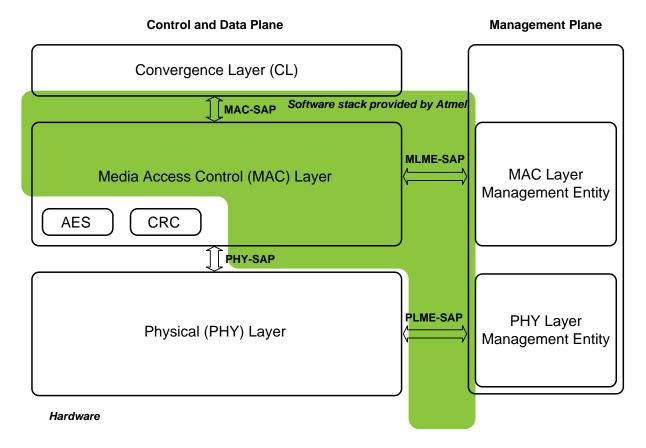


13.4 Atmel PRIME Software Stack

Looking for the best compromise between efficiency and versatility, Atmel has developed a PRIME software stack that implements and fulfills the capabilities of the MAC layer not carried out by ATPL210 hardware.

Atmel PRIME stack allows the final user to control the system by means of four sets of primitives, making the low level functions and the PHY-MAC interface transparent to the user.





By using Atmel PRIME stack, ATPL210 PHY and MAC layers can be controlled by means of the three access points shown in the figure: MAC-SAP (MAC Service Access Point), MLME-SAP (MAC Layer Management Entity Service Access Point) and PLME-SAP (PHY Layer Management Entity Service Access point). Furthermore, there are general functions to manage the software package.

An example of high level functions that manage the software stack is shown in Figure 13-5.

Please check the "Atmel PRIME Stack User Manual" for software package detailed and updated description and functionality.



CONTROL PRIMITIVES

- prime_Init
- prime_Start
- prime_Process
- prime Upd1msInt
- prime Upd618msInt
- prime_GetVersionNum
- prime_GetVersionStr
- prime_GetIdentificationStr

MAC PRIMITIVES

prime_MAC_ESTABLISH_request
prime_MAC_ESTABLISH_response
prime_MAC_RELEASE_request
prime_MAC_RELEASE_response
prime_MAC_JOIN_request
prime_MAC_JOIN_response_base
prime_MAC_LEAVE_request_base
prime_MAC_LEAVE_request_service
prime_MAC_REDIRECT_response
prime_MAC_Callback
prime_MAC_serial_process

PLME PRIMITIVES

- prime_PLME_RESET_request
- prime_PLME_SLEEP_request
- prime_PLME_RESUME_request
- prime_PLME_TESTMODE_request
- prime_PLME_GET_request
- prime_PLME_callback
- prime_PLME_SET_request
- prime_PLME_SET_BUF_request
- prime_PLME_receivedCmd
- prime_PLME_serial_process

MLME PRIMITIVES

prime_MLME_REGISTER_request
prime_MLME_UNREGISTER_request_base
prime_MLME_UNREGISTER_request_service
prime_MLME_PROMOTE_request_base
prime_MLME_DEMOTE_request_base
prime_MLME_DEMOTE_request_service
prime_MLME_GET_request
prime_MLME_LIST_GET_request
prime_MLME_SET_request
prime_MLME_SET_request
prime_MLME_SET_request
prime_MLME_SET_request
prime_MLME_SET_request
prime_MLME_SET_request
prime_MLME_SET_request
prime_MLME_SET_request
prime_MLME_SET_request
prime_MLME_SET_BUF_request
prime_MLME_callback
prime_MLME_receivedCmd
prime_MLME_serial_process

Atmel PRIME Stack



13.5 MAC Coprocessor Registers

13.5.1 SNA Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SNA				SNA(47:40)				@F
				SNA	.(7: 0)				@
Name:	SNA								
Address:	0xFE62 – 0	xFE67							
Access:	Read/write								
Reset:	0x00,, 0	x00							

• SNA: Sub Network Address

These registers store the 48-bit Sub Network Address. When the system Sub Network Address is available, the microcontroller must write it down so the Physical layer will be able to correctly calculate the CRC's, which depend on this parameter

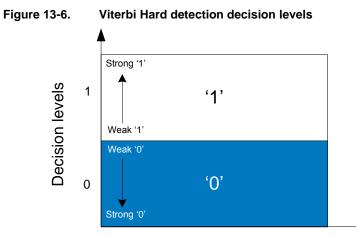


13.5.2 VITERBI_BER_HARD Register

Na	me	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
VITERBI_B	BER_HARD			VIT	ERBI_BE	R_HARD(7	:0)		
Name:	VITERBI_E	BER_HARD)						
Address:	0xFE36								
Access:	Read only								
Reset:	0x00								

• VITERBI_BER_HARD: This register stores the number of errors accumulated in a message reception using Viterbi hard* decision. The value is cleared by hardware each time a new message is received.

*Hard Decision: in "hard" detection there are only two decision levels. If the received value is different than the corrected one, the error value taken is "1". Otherwise, the error value taken is "0".



From the value in VITERBI_BER_HARD register it is possible to calculate de Bit Error Rate according to the following formula:



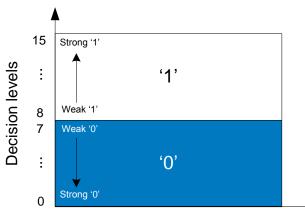
13.5.3 VITERBI_BER_SOFT Register

Na	me	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
VITERBI_E	BER_SOFT			VI	TERBI_BEI	R_SOFT(7:	0)		
Name:	VITERBI_E	BER_SOFT	-						
Address:	0xFE37								
Access:	Read only								
Reset:	0x00								

• VITERBI_BER_SOFT: This register stores a value proportional to the number of errors accumulated in a message reception using Viterbi soft* decision. The value is cleared by hardware each time a new message is received.

*Soft Decision: in "soft" decision there are fifteen decision levels. A strong '0' is represented by a value of "0", while a strong '1' is represented by a value of "15". The rest of values are intermediate, so "7" is used to represent a weak '0' and "8" represents a weak '1'. Soft decision calculates the error in one bit received as the distance in decision levels between the value received (a value in the range 0 to 15) and the corrected one (0 or 15).

Figure 13-7. Viterbi Hard detection decision levels





13.5.4 ERR_CRC32_MAC Registers

Nan	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ERR_CRC	32_MAC			E	RR_CRC32	_MAC(15:8	8)			@0xFEBA
				E	ERR_CRC3	2_MAC(7:0))			@0xFEBB
Name:	ERR_CR	C32_MAC								
Address:	0xFEBA	– 0xFEBB								
Access:	Read/wri	te								
Reset:	0x00, 0x00									

• ERR_CRC32_MAC: 16-bit value that stores the number of received messages that have been discarded by an error in the MAC layer CRC32.

Note: to clear this value, these registers must be reset by the microcontroller.



13.5.5 ERR_CRC8_MAC Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
ERR_CR	C8_MAC			E	RR_CRC8	_MAC(15:8))			@0xFEBC	
					ERR_CRC8	B_MAC(7:0)				@0xFEBD	
Name:	ERR_CR	C8_MAC									
Address:	0xFEBC	– 0xFEBD									
Access:	Read/wri	te									
Reset:	0x00, 0x0	0x00, 0x00									

• ERR_CRC8_MAC: 16-bit value that stores the number of received messages that have been discarded by an error in the payload MAC layer CRC8. Note: to clear this value, these registers must be reset by the microcontroller.



13.5.6 ERR_CRC8_AES Registers

Nai	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ERR_CR	C8_AES			ł	ERR_CRC8	_AES(15:8)				@0xFEBE
					ERR_CRC8	B_AES(7:0)				@0xFEBF
Name:	ERR_CR	C8_AES								
Address:	0xFEBE	– 0xFEBF								
Access:	Read/wri	te								
Reset:	0x00, 0x0	0, 0x00								

• ERR_CRC8_AES: 16-bit value that stores the number of received messages that have been discarded by an error in the payload AES CRC8. Note: to clear this value, these registers must be reset by the microcontroller.

Atmel ATPL210A [Datasheet] 130 43005A-ATPL-06/12



13.5.7 ERR_CRC8_MAC_HD Registers

Na	ime	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ERR_CRC	8_MAC_HD			ER	R_CRC8_N	IAC_HD(15	5:8)			@0xFEC0
				ER	R_CRC8_I	MAC_HD(7	:0)			@0xFEC1
Name:	ERR_CRC	B_MAC_H	C							
Address:	0xFEC0 – 0	0xFEC1								
Access:	Read/write									
Reset:	0x00, 0x00									

• ERR_CRC8_MAC_HD:16-bit value that stores the number of received messages that have been discarded by an error in the header MAC layer.

Note: to clear this value, these registers must be reset by the microcontroller.



13.5.8 ERR_CRC8_PHY Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ERR_CR	C8_PHY			I	ERR_CRC8	_PHY(15:8)				@0xFEC2
					ERR_CRC8	8_PHY(7:0)				@0xFEC3
Name:	ERR_CR	C8_PHY								
Address:	0xFEC2 ·	– 0xFEC3								
Access:	Read/wri	te								
Reset:	0x00, 0x0	00								

• ERR_CRC8_PHY: 16-bit value that stores the number of received messages that have been discarded by an error in the PHY layer CRC8. Note: to clear this value, these registers must be reset by the microcontroller.

Atmel ATPL210A [Datasheet] 132 43005A-ATPL-06/12



13.5.9 FALSE_DET_CONFIG Register

Name	- -	Bit 7 Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
FALSE_I _CONF			ERR_CRC8 _MAC	INVALID _PROTOCOL	ERROR _LEN	ERROR_PAD _LEN	UNKNOWN _PDU	UNKNOWN _SP
Name:	FALS	E_DET_CON	IFIG					
Address:	0xFE	C4						
Access:	Read	/write						
Reset:	0x10							
٠	:		Reserved b	oits				
٠	ERR_	_CRC8_MAC		essage has a c		-		error counter if yer CRC8 prese
٠	INVAL	ID_PROTOCO	received n		correct	PHY layer CR		error counter if PROTOCOL fie
٠	ERF	ROR_LEN:		essage has a c		•		error counter if eld indicates a r
٠	ERF	ROR_PAD_LI	received m		correct PH	•		error counter if EN field indicat
٠	UNP	(NOWN_PDI	received m	nessage has a	correct P	-	8 but the HT	error counter if field indicates
٠	UNM	(NOWN_SP:		essage has a c		•		error counter if RITY_PROTOCO



13.5.10 FALSE_DET Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
FALSE	_DET				FALSE_C	DET(15:8)				@0xFEC5
					FALSE_I	DET(7:0)				@0xFEC6
Name:	FALSE_[DET								
Address:	0xFEC5	– 0xFEC6								
Access:	Read/wri	te								
Reset:	0x00, 0x0	00								
٠	FALSE_	DET:	Erroneou	s non-disc	arded mes	sages.				

16-bit value that stores the number of received messages that have not been discarded since its PHY layer CRC8 is correct, but in which there are other/others incorrect field/fields. The fields to be taken into account to increase the counter in case they were wrong can be selected by FALSE_DET_CONFIG register.

Note: to clear this value, these registers must be reset by the microcontroller



13.5.11 MAX_LEN_DBPSK Register

Na	me	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
MAX_LEN	I_DBPSK	-	-		Γ	MAX_LEN_	DBPSK(5:0)	
Name:	MAX_LE	N_DBPSK							
Address:	0xFEC8								
Access:	Read/wri	te							
Reset:	0xFF								
٠	:		Reserved	lbits					
٠	MAX_L	EN_DBPS	•	llows to re		•			M symbols ation and

If a message in such modulation/encoding is received and its LEN field indicates a length above the threshold defined by MAX_LEN_DBPSK value, the message will be discarded.



13.5.12 MAX_LEN_DBPSK_VTB Register

N	lame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
MAX_LEN	_DBPSK_VTB				MA	X_LEN_D	BPSK_VT	B(5:0)]
Name:	MAX_LEN_D	BPSK_VTB								
Address:	0xFEC9									
Access:	Read/write									
Reset:	0xFF									
•	:		Res	erved bits	6					
•	MAX_LEN_D	BPSK_VTE	the	•	llows to re		•		OFDM sym PSK modu	
			indic	cates a le		e the three	shold defin	•	red and its X_LEN_DB	



13.5.13 MAX_LEN_DQPSK Register

Nan	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
MAX_LEN	DQPSK	-	-		Ν	MAX_LEN_I	DQPSK(5:0)	
Name:	MAX_LE	N_DQPSK							
Address:	0xFECA								
Access:	Read/writ	te							
Reset:	0xFF								
٠	:		Reserved	l bits					
٠	MAX_LI	EN_DBPS	•	llows to re		•			M symbols lation and

If a message in such modulation/encoding is received and its LEN field indicates a length above the threshold defined by MAX_LEN_DQPSK value, the message will be discarded.



13.5.14 MAX_LEN_DQPSK_VTB Registers

Ν	lame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
MAX_LEN	_DQPSK_VTB				MAX	X_LEN_DQ	PSK_VTB	5:0)		
Name:	MAX_LEN_D	QPSK_VTB								
Address:	0xFECB									
Access:	Read/write									
Reset:	0xFF									
٠	:		Rese	erved bits						
٠	MAX_LEN_	DQPSK_VT	the	•	lows to re		•		•	nbols, that lation and
			indic	ates a len		the thresh	old define			LEN field PSK_VTB



13.5.15 MAX_LEN_D8PSK Registers

Nai	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
MAX_LEN	I_D8PSK	-	-			MAX_LEN_	D8PSK(5:0))	
Name:	MAX_LE	N_D8PSK							
Address:	0xFECC								
Access:	Read/writ	te							
Reset:	0xFF								
•	:		Reserved	l bits					
٠	MAX_L	EN_D8PSI	•	llows to r		num length en working			•

If a message in such modulation/encoding is received and its LEN field indicates a length above the threshold defined by MAX_LEN_D8PSK value, the message will be discarded.



13.5.16 MAX_LEN_D8PSK_VTB Register

N	lame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
MAX_LEN	MAX_LEN_D8PSK_VTB				MA	X_LEN_D8	PSK_VTB	(5:0)		
Name:	MAX_LEN_D	8PSK_VTB								
Address:	0xFECD									
Access:	Read/write									
Reset:	0xFF									
٠	:		Res	erved bits						
•	MAX_LEN_	D8PSK_VT	the	-	llows to re		-			nbols, that lation and
			indi	cates a lei	in such n ngth above ssage will l	the thres	hold defin			



13.5.17 AES_PAD_LEN Register

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
AES_PA	D_LEN		-	-			AES_PAD	_LEN(3:0)	
Name:	AES_PA	D_LEN							
Address:	0xFE25								
Access:	Read/wri	te							
Reset:	0x00								
٠	:		Reserved	bits					
٠	AES_P	AD_LEN:	AES protocol works over 16-bytes-lenght blocks. When a block is not 16-by long, this register indicates the number of padding bytes to append.						
			This regis	ster takes	values betw	veen 0 and	15.		
			In transmission, if encryption is being used, microcontroller must write the Al padding length in this register.						
			In no-enc	rypted trai	nsmission a	and in rece	ption, the v	alue in this	register is i



13.5.18 AES_DATA_IN Registers

Name		Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0			
AES_DATA	AES_DATA_IN			AES_DATA_IN(127:120)								
	ļ											
	ĺ											
					AES_DAT	A_IN(7: 0)				@FFAF		
Name:	AES	_DATA_IN										
Address:	0xFF	FA0 – 0xFFAF										
Access:	Rea	d/write										
Reset:	0x00),, 0x00										
٠	AES	_DATA_IN	l: Input	buffer for A	ES128 bloc	k.						
		This buffer can be written to be encrypted/decrypted by the key in KEY_ (see 13.5.20) register.										
			The r	esulting dat	a could be	read at AES	_DATA_OL	JT (see 13.	5.19) regist	ers.		

13.5.19 AES_DATA_OUT Registers

Name	e	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0			
AES_DATA	A_OUT			AES_DATA_OUT(127:120)							₽FFB0	
]		
										Ì		
										İ		
			AES_DATA_OUT(7:0)									
Name:	AES_C	DATA_OUT								_		
Address:	0xFFB	0 – 0xFFBF										
Access:	Read c	only										
Reset:	0x00, .	, 0x00										
٠	AES_D	DATA_OUT	: Output	buffer for A	ES128 bloc	:k.						
			This bu	uffer stores	the result	of the e	encryption/de	cryption	processing	of	data ir	۱

AES_DATA_IN (see 13.5.18) register with the key in KEY_PERIPH (see 13.5.20) register.

13.5.20 KEY_PERIPH Registers

Name		Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
KEY_PERI	PH				KEY_PERIF	PH(127:120)				@FFC0
	<u> </u>									
					KEY_PER	IPH (7: 0)				@FFCF
Name:	KEY_	PERIPH								
Address:	0xFF	C0 – 0xF	FCF							
Access:	Read	/write								
Reset:	KEY_	PERIPH(127:120) :	0x00;	KEY_PERI	PH(119:112) :	0x01;	KEY_PERIP	H(111:104) :	0x02;
	KEY_	PERIPH(103:96) :	0x03;	KEY_PERIF	PH(95:88) :	0x04;	KEY_PERIP	H(87:80) :	0x05;
	KEY_	PERIPH(79:72) :	0x06;	KEY_PERI	PH(71:64) :	0x07;	KEY_PERIP	H(63:56) :	0x08;
	KEY_	PERIPH(55:48) :	0x09;	KEY_PERIF	PH(47:40) :	0x0A;	KEY_PERIP	H(39:32) :	0x0B;
	KEY_	PERIPH(31:24) :	0x0C;	KEY_PERI	PH(23:16) :	0x0D;	KEY_PERIP	H(15:8) :	0x0E;
	KEY_	PERIPH(7:0) :	0x0F;						
٠	KEY_	PERIPH	: Key	for AES128	block when	used as pe	ripheral p	art.		

This key is used for encrypting/decrypting data in AES_DATA_IN registers.



13.5.21 KEY_PHY Registers

Name	Bit 7 Bit 6	E	Sit 5	Bit 4	Bit	3	Bit 2	Bit 1	B it 0	
KEY_PHY	KEY_PHY(127:120)									@FFD0
				KEY_PH	HY(7: 0)					@FFDF
Name:	KEY_PHY									
Address:	0xFFD0 – 0xFFDF									
Access:	Read/write									
Reset:	KEY_PHY(127:120):	0x00;	KEY_P	HY(119:112	2):	0x01;	KEY_PI	HY(111:104) :	0x02;	
	KEY_PHY(103:96) :	0x03;	KEY_P	HY(95:88) :		0x04;	KEY_PI	HY(87:80) :	0x05;	
	KEY_PHY(79:72) :	0x06;	KEY_P	HY(71:64) :		0x07;	KEY_PI	HY(63:56) :	0x08;	
	KEY_PHY(55:48) :	0x09;	KEY_P	HY(47:40) :		0x0A;	KEY_PI	HY(39:32) :	0x0B;	
	KEY_PHY(31:24) :	0x0C;	KEY_P	HY(23:16) :		0x0D;	KEY_PI	HY(15:8) :	0x0E;	
	KEY_PHY(7:0) :	0x0F;								

• **KEY_PHY:** Key for AES128 block when used by the physical layer

This key is used in real time encryption/decryption for Security Profile 1. When any of the DMA channels of the physical layer accesses to the memory, then this key and the input data are multiplexed to the AES128-core. Also output data is multiplexed in order to provide encrypted/decrypted data to the physical buffer.





13.5.22 AES_SFR Register

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
AES_SFR							READY	START	CIPHER	
Name:	AES_SFR									
Address:	0xFFE0									
Access:	Read/write									
Reset:	0x00									
٠	:	Reserv	ed bits.							
٠	READY:	Flag to	indicate e	ncryption/d	ecryption p	process cor	npletion.			
		When t	When the encryption/decryption has been completed, this flag is set to '1'.							
		This fla	This flag is automatically cleared when an encryption/decryption process begins.							

- **START:** When this bit is set to '1', the encryption/decryption process is triggered.
 - If encryption/decryption starts successfully, then this bit is automatically cleared to '0'.
- **CIPHER:** This field indicates if data must be encrypted or decrypted.
 - '0' Decryption mode
 - '1' Encryption mode



14. ATPL210 PRIME PHY Layer

14.1 ATPL210 PHY Layer

The physical layer of ATPL210 consists of a hardware implementation of the PRIME Physical Layer Entity, which is an Orthogonal Frequency Division Multiplexing (OFDM) system in the CENELEC A-band. This PHY layer transmits and receives MPDUs (MAC Protocol Data Unit) between neighbor nodes.

From the transmission point of view, the PHY layer receives its inputs from the MAC (Medium Access Control) layer, via DMA. At the end of transmission branch, data is output to the physical channel.

On the reception side, the PHY layer receives its inputs from the physical channel, and at the end of reception branch, the data flows to the MAC layer, via DMA.

A PHY layer block diagram is shown below:

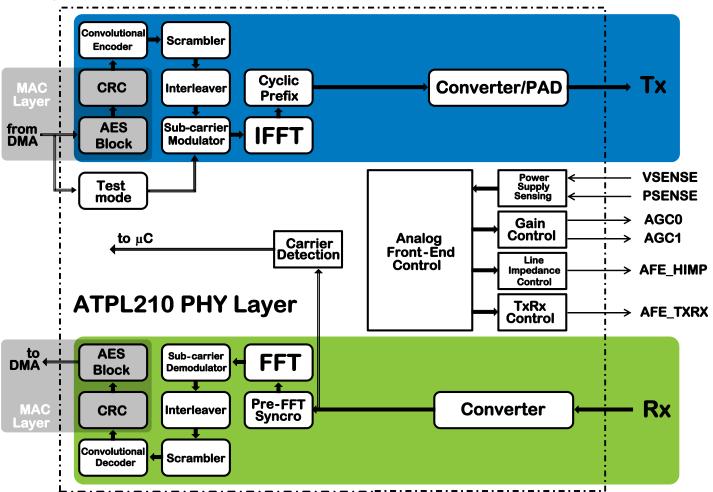


Figure 14-1. ATPL210 PHY Layer Block Diagram

The diagram can be divided in four sub-blocks: Transmission branch, Emission branch, Analog Front End control and Carrier Detection.

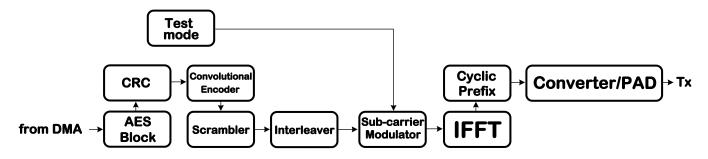


14.1.2 Transmission and Reception branches

Phy layer takes data to be sent from dedicated DMA channel (PHY_TX). 128-bit AES encryption is done "on the fly", and the Cyclic Redundancy Check (CRC) fields are hardware-generated in real time. These CRCs are properly appended to the transmission data. The rest of the chain is hardware-wired, and performs automatically all the tasks needed to send data according to PRIME specifications.

In Figure 14-2, the block diagram of the transmission branch Is shown.





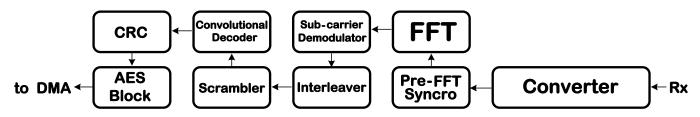
The output is differentially modulated using a BPSK/DQPSK/D8PSK scheme. After modulation, IFFT (Inverse Fourier Transform) block and cyclic prefix block allows to implement an OFDM scheme.

A Converter and a Power Amplifier Driver is the last block in the transmission branch. This block is responsible for adjusting the signal to reach the best transmission efficiency, thus reducing consumption and power dissipation.

Test mode: When selected, test mode injects data directly to Sub-carrier modulation block. When in test mode, data can be injected continuously to the line using only a set of selected frequencies, in order to test channel behavior.

The reception branch performs automatically all the tasks needed to process received data. Phy layer delivers data to MAC layer through the dedicated DMA channel (PHY_RX).





14.1.3 Carrier Detection

Looking for an easy detection of incoming messages, PRIME specification defines a chirp signal located at the beginning of the PRIME frames devised to ease synchronization in the receptor. By means of detection techniques, the receiver can know accurately when the chirp has been completely received and then the correct instant when the frame begins.

Before starting a transmission, it is also necessary to use carrier detection in order to check if another device is already emitting, thus avoiding collisions. If any device is emitting, the carrier detection triggers a microcontroller interruption and sets an internal flag, thus the transmission will be stopped.

The main drawback of this process is that chirp signal length (2.4 milliseconds) is not short enough to guarantee very low collision ratio.



To improve this drawback, the ADD1320 PLC modem implements two different algorithms to detect the carrier as soon as possible, aiming to reduce collisions and improving the medium access behavior. By these early detection techniques, the system achieves low collision ratio, and the communication throughput increases significantly.

14.1.4 Analog Front End control

The Phy layer controls the Analog Front End by means of four sub-blocks:

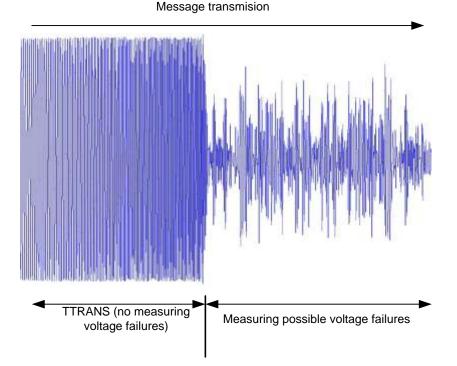
- Power Supply sensing
- Gain control
- Line Impedance control
- TxRx control

14.1.4.1 Power Supply Sensing: VSENSE and PSENSE

The power supply is continuously monitored to avoid power supply failures that could damage the supply device. This block senses the power channel using two different inputs:

• VSENSE: VSENSE detects whether voltage falls below 3.3v during a number of cycles while a message is being transmitted. This measurement is done after a transitory guard time (TTRANS in figure below). If a Voltage failure occurs, the transmission is shut down and sending messages again will be not possible if an internal flag (VFAILURE) is not previously cleared.

Figure 14-4. Transitory guard time in message transmission



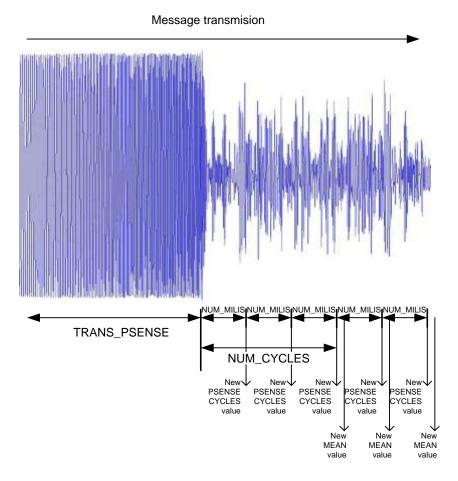
 PSENSE: PSENSE measures the power source current consumption, shutting down the transmission if the consumption exceeds a defined threshold (stored in MAXPOT phy layer registers, see 14.5.34).
 This measurement is done after a transitory guard time. As the current measurement varies over time, an averaging is done taking into account an average parameter (Alpha), a configurable number of cycles (NUMCYCLES, see 14.5.35) and a configurable length of each cycle (A_NUMMILIS, see 14.5.36).



If a power failure occurs, the transmission is shut down and sending messages again will be not possible if an internal flag (PFAILURE, see 14.5.22) is not previously cleared.

The system considers that a power failure has occurs when the value read from MEAN registers (see 14.5.30) is above the user-definable value stored in MAXPOT registers.





Psense and Vsense configurations parameters are automatically set by the Phy layer.

See related peripheral registers for more information about Psense and Vsense.

14.1.4.2 Gain Control

This block implements two Automatic Gain Control outputs to adjust the received signal level to a suitable range. Both of them are set to '1' when the received signal is above two system thresholds in order to activate external attenuators placed in the external analog front end.

The value of these outputs is set during the beginning of a received message and is hold until the end of the message.

AGC0 and AGC1 follow different algorithms, thus using both of them ensures a more accurate gain control.

See AGC_CONFIG Register for information about AGC configuration.

14.1.4.3 Line Impedance Control

This block modifies the configuration of the Analog Front End by means of AFE-HIMP output. When working with a suitable external configuration, the system can change the filter conditions in order to adjust its behavior to the line impedance values. See last ATPL210 reference design for further information about Line Impedance topologies.



14.1.4.4 TxRx Control

This block modifies the configuration of the Analog Front End by means of AFE-TXRX output. Thus is possible to change filter conditions between transmission/reception.

See reference design for further information about TxRx control.

14.2 PHY parameters

As described below, the PHY layer is specified by certain main parameters, which are fixed for each specific constellation/coding combination. These parameters have to be identical in a network in order to achieve compatibility.

Table 14-1.	PRIME Phy main parameters
-------------	---------------------------

PRIME Phy parameter	Value
Base Band Clock (Hz)	250000
Subcarrier spacing (Hz)	488,28125
Number of data subcarriers	84 (header), 96 (payload)
Number of pilot subcarriers	13 (header), 1 (payload)
FFT interval (samples)	512
FFT interval (μs)	2048
Cyclic Prefix (samples)	48
Cyclic Prefix (µs)	192
Symbol interval (samples)	560
Symbol interval (µs)	2240
Preamble period (µs)	2048

Table 14-2 shows the PHY data rate during payload transmission, and maximum MSDU length for various modulation and coding combinations

Table 14-2. Phy parameters depending on the modulation

	DBPSK		DQPSK		D8PSK	
Convolutional Code (1/2)	On	Off	On	Off	On	Off
Information bits per subcarrier	0,5	1	1	2	1,5	3
Information bits per OFDM symbol	48	96	96	192	144	288
Raw data rate (kbps approx)	21,4	42,9	42,9	85,7	64,3	128,6
MAX MSDU length with 63 symbols (bits)	3016	6048	6040	12096	9064	18144

Table 14-3 shows the modulation and coding scheme and the size of the header portion of the PHY frame



Table 14-3. Header parameters

	DBPSK
Convolutional Code (1/2)	On
Information bits per subcarrier	0,5
Information bits per OFDM symbol	42

All the parameters of the physical layer such as the base band clock, subcarrier spacing, number of subcarriers...; are defined in PRIME Specification, and have to be identical in a network in order to achieve compatibility.

14.3 PHY Protocol Data Unit (PPDU) Format

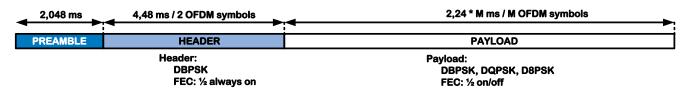
Figure 14-6 shows how OFDM symbols are transmitted in a PPDU (Physical layer Protocol Data Unit). The preamble is used at the beginning of every PPDU for synchronization purposes.

Figure 14-6. PHY layer transmitter block diagram



Phy layer adaptively modifies attenuation values applied to the whole signal. Also, additional attenuations are applied to the chirp section of the signal (preamble) and to the rest of the signal itself (header+payload), to smoothly adapt amplitude values and transitions.

Figure 14-7. PPDU OFDM symbols and duration



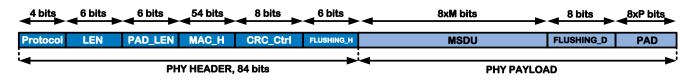
14.4 PHY Service Specification

There is an interface specified in PRIME for the PHY layer, with several primitives relative to both data and control planes.

PHY layer has a single 20-bit free-running clock measured in 10µs steps. Time measured by this clock is the one to be used in some PHY primitives to indicate a specific instant in time.

ATPL210A includes a hardware implementation of this clock, which consists of a 20-bit register. This register is readonly and it can be accessed as a 32-bit variable by the ADD8051C3A microcontroller.

Figure 14-8. Header and payload structure





Prime specifies a complete set of primitives to manage the PHY Layer, and the PHY-SAP (PHY Service Access Point) from MAC layer. Atmel PRIME stack integrates all this functions, making them transparent to the final user and simplifying the management.





14.5 PHY Layer registers

14.5.1 PHY_SFR Register

This register is described in 12.11.1, DMA Configuration Registers section.



14.5.2 SYS_CONFIG Register

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
SYS_CONF	FIG			CONV_PD	PHY_PD	PHY_ERR_EN	PHY_ERR	PHY_RST	
Name:	SYS_CONFIG								
Address:	0xFE2C	JxFE2C							
Access:	Read/write								
Reset:	0x04								

- --: Reserved bits
- CONV_PD: Converter Power Down

Microcontroller can activate internal converter power down mode by setting this bit. When internal converter is in power down mode, the system is unable to receive.

This bit is high-level active.

PHY_PD: PHY Power Down

This bit shuts down Physical Layer clock. When in PHY power down mode, all the system blocks involved in communication remain inactive. Thus, the system will be unable to transmit or receive. The next sequence must be respected to ensure proper power down:

Setting PHY power down mode

1-Set Physical Layer reset (SYS_CONFIG(0)), PHY_RST='1'

2-Set CONV_PD and PHY_PD fields

Exiting PHY power down mode

1-Clear CONV_PD and PHY_PD fields

2-Clear Physical Layer reset (SYS_CONFIG(0)), PHY_RST='0'

This bit is high-level active.

PHY_ERR_EN:Physical Layer Watchdog enable

This bit enables or disables Physical layer watchdog. Physical layer watchdog is enabled by default.

This bit is high-level active.

PHY_EN: Physical Layer Error Flag

This flag indicates if a Physical layer error has occurred. Physical layer watchdog has a 200milliseconds sampling period. When Physical layer detects an error, it activates the Physical layer interrupt and this flag is set.

To restore situation, microcontroller must reset Physical layer by means of PHY_RST bit (SYS_CONFIG(0)).

PHY_RST: Physical Layer Reset

This bit resets the Physical layer. To perform a Physical layer reset cycle, microcontroller must set this bit to '1' and then must clear it to '0'.



14.5.3 PHY_CONFIG Register

Name	Bit 7 Bit 6		Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
PHY_CON	NFIG		CINR_MODE	PAD_LEN_AC	AES_EN	CD_MOD1_EN	CD_MOD2_DET	MAC_EN		
Name:	PHY_CON	FIG								
Address:	0xFE68	0xFE68								
Access:	Read/write									
Reset:	0x1F									

- --: Reserved bits
- **CINR_MODE:**Carrier to Interference + Noise Ratio mode
 - This bit enables/disables CINR mode when set to '1'.
 - '0': CINR mode disabled.
 - '1': CINR mode enabled.
- PAD_LEN_AC: This field allows the system to work with two different representations of the Phy header PAD_LEN field (PAD_LEN represented before coding or PAD_LEN represented after coding).
 - '0': PAD_LEN field in PHY header is represented before coding. This is the suitable value to fulfill PRIME specification.
 - '1': PAD_LEN field in PHY header is represented after coding.
- **AES_EN:** This field enables/disables "on the fly" AES encryption and decryption by hardware.
 - '0': "on the fly" AES encryption/decryption disabled.
 - '1': "on the fly" AES encryption/decryption enabled.
- **CD_MOD1_EN:** This field enables/disables Carrier Detection mode 1.
 - '0': Carrier Detection mode 1 disabled.
 - '1': Carrier Detection mode 1 enabled.
- **CD_MOD2_DET:** This field enables/disables Carrier Detection mode 2.

•

- '0': Carrier Detection mode 2 disabled.
- '1': Carrier Detection mode 2 enabled.
- MAC_EN: This field enables/disables CRC processing by hardware.
 - '0': CRC processing disabled.
 - '1': CRC processing enabled.



14.5.4 ATTENUATION Register

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
ATTENUA	ATTENUATION ATTENUATION(7:0)							
Name:	ATTENUATION							
Address:	0xFE24							
Access:	Read/write							
Reset:	0xFF							
٠	ATTENUATION:	Global	attenuation fo	or the transm	nitted signal (chirp+signal). T	he 16-bit sig	nal level is

ON: Global attenuation for the transmitted signal (chirp+signal). The 16-bit signal level is multiplied by this 8-bit value and the result is truncated to 16 bits.
 Attenuation value = 0xFF → the transmitted signal amplitude is not attenuated.
 Attenuation value = 0x00 → the transmitted signal amplitude is nullified.





14.5.5 ATT_CHIRP Register

Name	Name Bit 7		Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ATT_CHIR	P	ATT_CHIRP(7:0)							
Name:	ATT_CHIRP								
Address:	0xFE9B								
Access:	Read/write								
Reset:	0xFF								
٠	ATT_CHIRP:	This re	gister stores	the attenuati	on value for t	he chirp. The 1	6-bit chirp da	ata is	

This register stores the attenuation value for the chirp. The 16-bit chirp data is multiplied with this 8-bit value and the 24-bit result is truncated to 16 bits. Attenuation value = $0xFF \rightarrow$ the chirp amplitude is not attenuated Attenuation value = $0x00 \rightarrow$ the chirp amplitude is nullified



14.5.6 ATT_SIGNAL Register

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
ATT_SIGN	AL	ATT_SIGNAL(7:0)								
Name:	ATT_SIGNAL									
Address:	0xFE9C									
Access:	Read/write									
Reset:	0xFF									
٠	ATT_SIGNAL:		•			ne signal witho it value and th	•			

truncated to 16 bits. Attenuation value = $0xFF \rightarrow$ the signal amplitude is not attenuated Attenuation value = $0x00 \rightarrow$ the signal amplitude is nullified



14.5.7 TX_TIME Registers

Name	Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 B it 0									
TX_TIME	TX_TIME(19:12)	@0xFE26								
	TX_TIME(11:4)									
	TX_TIME(3:0) "0000"									
	"0000000"									

Name: TX_TIME

Address: 0xFE26 - 0xFE29

Access: Read/write

Reset: 0x00, ..., 0x00;

• **TX_TIME:** This 20-bit value sets the time instant when the MPDU (MAC Protocol Data Unit) has to be transmitted. The time is expressed in 10μs steps.

When writing a new value to TX_TIME register, a specific writing order must be taken, always from the most significant byte (TX_TIME(19:12) at address 0xFE26) to the least significant byte (TX_TIME(3:0) at address 0xFE28), and it is required to write the 3 bytes to avoid wrong time comparisons in transmission.

The 20-bit TX_TIME value is managed by the microcontroller as a 4-byte variable. The TX_TIME value is aligned to the 20 most significant bits, being the 12 least significant bits padded with zeros.

This register is used by the physical layer for being in accordance with PRIME specifications about transmission time (see PRIME spec.)

Note: TXRX bit (PHY_SFR(2)) has to be cleared to '0' in order to init transmission. Once this bit has been cleared, the transmission will start when TIMER_BEACON_REF value is equal to TX_TIME.





14.5.8 TIMER_FRAME Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0				
TIMER_FRAME				TIMER_FR	AME(19:12)				@0xFE2D			
				TIMER_FR	AME(11:4)				@0xFE2E			
		TIMER_FRAME(3:0) "0000"										
		"0000000"										
Name: TIME	R_FRAME											

Address: 0xFE2D – 0xFE30

Access: Read only

Reset: 0x00, ..., 0x00;

TIMER_FRAME: Time of receipt of the preamble associated with the PSDU (PHY Service Data Unit).

It is expressed in $10\mu s$ steps and is taken from the physical layer timer TIMER_BEACON_REF.

It is set by hardware and is a read-only register.

This register is used by the physical layer for being in accordance with PRIME specification about reception time (see PRIME specification).

The 20-bit TIMER_FRAME value is managed by the microcontroller as a 4-byte variable. The TIMER_FRAME value is aligned to the 20 most significant bits, being the 12 least significant bits padded with zeros. This simplifies arithmetic calculations with time values.



14.5.9 TIMER_BEACON_REF Registers

Na	ame	Bit 7	Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 B it 0									
TIMER_BE	ACON_REF		TIMER_BEACON_REF(19:12)									
			TIMER_BEACON_REF (11:4)									
			TIMER_BEACON_REF (3:0) "0000"									
			"0000000"									
Name:	TIMER_BE	ACON_RE	F									
Address:	0xFE47 – 0	0xFE4A										
Access:	Read only	ad only										

Reset: 0x00, ..., 0x00;

• **TIMER_BEACON_REF:** Timer for the physical layer, which consists of a single 20-bit free-running clock measured in 10µs steps.

It indefinitely increases a unit each 10 microseconds from 0 to 1048575, overflowing back to 0.

It is set by hardware and is a read-only register.

This register is used by the physical layer for being in accordance with PRIME specification. It is reserved 32-bit in data memory to be able to declare as 32-bit variable. The 20-bit register MSB is aligned to the 32-bit variable MSB, in order to simplify arithmetic calculations with time values.



14.5.10 RX_LEVEL Registers

Ν	lame	Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 B it 0								
RX_	LEVEL				RX_LE∖	′EL(15:8)				@0xFE31
			RX_LEVEL(7:0)							
Name:	RX_LEVEL									
Address:	0xFE31 – 0x	FE32								
Access:	Read only									
Reset:	0x00; 0x00	; 0×00								

• **RX_LEVEL**:These registers store the autocorrelation level of the chirp signal.

When the reception process has started, these registers are set by hardware.



14.5.11 RSSI_MIN Register

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
RSS	SI_MIN	RSSI_MIN(7:0)							
Name:	RSSI_MIN								
Address:	0xFE33								
Access:	Read only								
Reset:	0xFF								

• RSSI_MIN: Received Signal Strength Indication Min

This register stores the minimum RSSI value measured in the last message received. The measurement is done at symbol level. The value is stored in ¼dB steps

14.5.12 RSSI_AVG Register

N	lame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
RSS	SI_AVG				RSSI_A	VG(7:0)			
Name:	RSSI_AVG								
Address:	0xFE34								
Access:	Read only								
Reset:	0x00								

• **RSSI_AVG:** Received Signal Strength Indication Average

This register stores the average RSSI value measured in the last message received. The measurement is done at symbol level. The value is stored in ¼dB steps

14.5.13 RSSI_MAX Register

N	ame	Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 Bit 0								
RSS	I_MAX	RSSI_MAX(7:0)							0xFE35	
Name:	RSSI_MAX									
Address:	0xFE35									
Access:	Read only									
Reset:	0x00									

• RSSI_MAX: Received Signal Strength Indication Max

This register stores the maximum RSSI value measured in the last message received. The measurement is done at symbol level. The value is stored in ¼dB steps

14.5.14 CINR_MIN Register

N	lame	Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 Bit 0								
CIN	R_MIN	CINR_MIN(7:0)							0xFE38	
Name:	CINR_MIN									
Address:	0xFE38									
Access:	Read only									
Reset:	0xFF									

CINR_MIN: Carrier to Interference + Noise ratio Min

This register stores the minimum CINR value measured in the last message received.

In order to calculate CINR properly, the algorithm takes beacon-type messages as a reference, since this message type allows knowing its content beforehand.

The system uses a table that must be loaded with the beacon data to be received, so CINR mode must be activated (see PHY_CONFIG register) and the same procedure used to send beacons must be followed. As CINR mode is activated, physical layer will load the message in the table instead of sending it (table load time is in the order of microseconds, and is much shorter than the one used to send the message).

Once the table is loaded, CINR must be disabled, and next messages CINR will be calculated taken the beacon loaded in the table as reference.

The measurement is done at symbol level.

The value is stored in ¼dB steps.



14.5.15 CINR_AVG Register

N	lame	Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 Bit 0							B it 0	
CIN	R_AVG	CINR_AVG(7:0)							0xFE39	
Name:	CINR_AVG									
Address:	0xFE39									
Access:	Read only									
Reset:	0x00									

CINR_AVG:Carrier to Interference + Noise ratio Average

This register stores the average CINR measured in the last message received.

In order to calculate CINR properly, the algorithm takes beacon-type messages as a reference, since this message type allows knowing its content beforehand.

The system uses a table that must be loaded with the beacon data to be received, so CINR mode must be activated (see PHY_CONFIG register) and the same procedure used to send beacons must be followed. As CINR mode is activated, physical layer will load the message in the table instead of sending it (table load time is in the order of microseconds, and is much shorter than the one used to send the message).

Once the table is loaded, CINR must be disabled, and next messages CINR will be calculated taken the beacon loaded in the table as reference

The measurement is done at symbol level.

The value is stored in ¼dB steps.





14.5.16 CINR_MAX Register

N	lame	Bit 7 Bit 6 Bit 5 Bit 4 Bit 3 Bit 2 Bit 1 B it 0								
CIN	R_MAX	CINR_MAX(7:0)							0xFE3A	
Name:	CINR_MAX									
Address:	0xFE3A									
Access:	Read only									
Reset:	0x00									

CINR_MAX:Carrier to Interference + Noise ratio Max

This register stores the maximum CINR value measured in the last message received.

In order to calculate CINR properly, the algorithm takes beacon-type messages as a reference, since this message type allows knowing its content beforehand.

The system uses a table that must be loaded with the beacon data to be received, so CINR mode must be activated (see PHY_CONFIG register) and the same procedure used to send beacons must be followed. As CINR mode is activated, physical layer will load the message in the table instead of sending it (table load time is in the order of microseconds, and is much shorter than the one used to send the message).

Once the table is loaded, CINR must be disabled, and next messages CINR will be calculated taken the beacon loaded in the table as reference

The measurement is done at symbol level.

The value is stored in ¼dB steps.



14.5.17 EVM_HEADER Registers

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
EVM_	HEADER				EVM_HEA	DER(15:8)				@0xFE3B
					EVM_HEA	DER (7:0)				@0xFE3C
Name:	EVM_HEADE	ĒR								
Address:	0xFE3B – 0xl	FE3C								
Access:	Read only									
Reset:	0x00; 0x00									
٠	EVM_HEAD	ER:	Header Erro	r Vector N	/lagnitude-					
			These regis measured in					imum erro	r vector	magnitude
			The 7 msb EVM_HEAD	• –	•	<i>,,</i> .		• .		0
			This registe specification variable. Th simplify arith	. It is res ne 20-bit r	erved 32-b egister MS	it in data B is aligne	memory t d to the 3	to be able	to declare	e as 32-bit



14.5.18 EVM_PAYLOAD Registers

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
EVM_F	PAYLOAD				EVM_PAY	LOAD(15:8)			@0xFE3D	
					EVM_PAY	/LOAD(7:0)				@0xFE3E	
Name:	EVM_PAYLC	DAD	D								
Address:	0xFE3D – 0x	FE3E									
Access:	Read only										
Reset:	0x00; 0x00										
٠	EVM_PAYLO	DAD: P	AD: Payload Error Vector Magnitude-								
		These registers store in a 16-bit value the maximum error vector r measured in the reception of a message payload.								magnitude	

The 7 msb (EVM_PAYLOAD(15:9)) represent the integer part in %, being the EVM_PAYLOAD(8:0) bits the fractional part if more precision were required.



14.5.19 EVM_HEADER_ACUM Registers

Name		Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0				
EVM_HEADER	R_ACUM		EVM_HEADER_ACUM(19:12)										
			EVM_HEADER_ACUM (11:4)										
		EVI	EVM_HEADER_ACUM (3:0) "0000"										
			"0000000"										
Name: E\	/M_HEAD	DER_ACUN	Л										

Address: 0xFE3F – 0xFE42

Access: Read only

•

Reset: 0x00, ..., 0x00;

EVM_HEADER_ACUM: Header Total Error Vector Magnitude Accumulator

When receiving an OFDM symbol, the sum total of all its individual carriers EVMs is calculated in order to further calculate the average EVM value. These registers store the maximum sum total between the two OFDM symbols received in a message header.

14.5.20 EVM_PAYLOAD_ACUM Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0				
EVM_PAYLOAD_ACUM			EVM_	_PAYLOAD	_ACUM(19	9:12)			@0xFE43			
		EVM_PAYLOAD_ACUM(11:4)										
	EVM_PAYLOAD_ACUM(3:0) "0000"											
		"0000000"										

Name: EVM_PAYLOAD_ACUM

Address: 0xFE43 – 0xFE46

Access: Read only

•

Reset: 0x00, ..., 0x00;

EVM_PAYLOAD_ACUM: Payload Total Error Vector Magnitude Accumulator

When receiving an OFDM symbol, the sum total of all its individual carriers EVMs is calculated in order to further calculate the average EVM value. These registers store the maximum sum total between all the OFDM symbols received in a message payload.



14.5.21 RMS_CALC Register

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
RMS	_CALC				RMS_C	ALC(7:0)				
Name:	RMS_CALC									
Address:	0xFE58									
Access:	Read only									
Reset:	0x00									

RMS_CALC: This register stores an 8-bit value which magnitude is proportional to the emitted signal amplitude.

By measuring the amplitude of the emitted signal, the hardware can estimate the power line input impedance. Thus hardware can adjust emission configuration appropriately.





14.5.22 VSENSE_CONFIG Register

Name		Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
VSENSE_CONFIG						PFAILURE	PSENSE_SOFT	VFAILURE	VSENSE_EN
Name:	VSENSE_CONFIG								
Address:	0xFE59								
Access:	Read only								
Reset:	0x00								

• **PFAILURE:** Power Failure Flag

This flag is set when a power failure occurs. The transmission is stopped and a new transmission in not possible if this flag is not cleared previously.

When a power failure occurs, a consideration about decreasing voltage amplitude in the source should be taken.

This flag must be cleared by software.

- PSENSE_SOFT: Current measurement is done every time a transmission takes place. With PSENSE_SOFT the system can force a continuous current measurement, including both idle and transmission states.
 - '0': Current consumption is measured every time a transmission begins (after a guard time defined by TRANS_PSENSE). NUMMILIS, NUMCYCLES and TRANS_PSENSE values must be taken into account to accurate PSENSE measurements. This is the default mode and it is the expected one when ATPL210 is working.
 - '1': Current consumption is measured both in idle and transmission states. This mode is useful for design purposes, in order to find suitable values for the current threshold (MAXPOT registers) depending on the external net requirements.
- VFAILURE: Voltage Failure Flag

This flag is set to '1' when a voltage failure occurs. The transmission is stopped and a new transmission is not possible if this flag is not cleared previously.

When a voltage failure occurs, a consideration about decreasing voltage amplitude in the source should be taken.

This flag must be cleared by software.

VSENSE_EN:VSENSE enable

This bit enables VSENSE.

- '0': VSENSE disabled (default).
- '1': VSENSE enabled.



14.5.23 NUM_FAILS Register

N	ame	Bit 7	Bit 6	Bit 5	Bit 4 Bit 3	Bit 2	Bit 1	B it 0	
NUM	_FAILS				NUM_FAILS(7:0)				
Name:	NUM_FAILS								
Address:	0xFE5A								
Access:	Read/write								
Reset:	0x02								

 NUM_FAILS: This register stores the number of 50 ns cycles (clk=20MHz) during which a voltage failure must be detected before shutting off the transmission and setting VFAILURE flag. This detection shall be done after a guard period set by TTRANS from the beginning of the transmission.

Default value: $0x02 \rightarrow 2 * 50 = 100$ ns



14.5.24 TTRANS Register

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0	
ТТ	RANS				TTRA	NS(7:0)				
Name:	TTRANS									
Address:	0xFE5B									
Access:	Read/write									
Reset:	0x2D									

TTRANS: This register stores the number of 50 μs cycles (clk=20MHz) to wait from the beginning of the transmission before looking for a possible voltage failure.
 Default value: 0x2D → 45 * 50 - 2.25mc (Thus, values are not expected until the

Default value: $0x2D \rightarrow 45 * 50 = 2.25ms$ (Thus, voltage failures are not expected until the end of chirp signal period)





14.5.25 AGC0_KRSSI Register

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
AGCO	_KRSSI				AGC0_K	RSSI(7:0)			
Name:	AGC0_KRSS	C0_KRSSI							
Address:	0xFE5C								
Access:	Read/write								
Reset:	0x00								

• AGC0_KRSSI: This register is used to correct RSSI (Received Signal Strength Indication) computation when Automatic Gain Control 0 (AGC0) is active.



14.5.26 AGC1 KRSSI Register

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
AGC1	_KRSSI				AGC1_K	RSSI(7:0)			
Name:	AGC1_KRSS	AGC1_KRSSI							
Address:	0xFE5D								
Access:	Read/write	Read/write							
Reset:	0x00								

• AGC1_KRSSI: This register is used to correct RSSI (Received Signal Strength Indication) computation when Automatic Gain Control 1 (AGC1) is active.



14.5.27 ZERO_CROSS_TIME Registers

Na	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0				
ZERO_CR	ROSS_TIME		ZERO_CROSS_TIME(19:12)										
		ZERO_CROSS_TIME (11:4)											
	ZERO_CROSS_TIME (3:0)								"0000"				
			"0000000"										
Name:	ZERO_CRO	DSS_TIME	SS_TIME										
Address:	0xFE69 – 0	0xFE69 – 0xFE6C											

Access: Read only

Reset: 0x00, ..., 0x00;

ZERO_CROSS_TIME: Instant in time at which the last zero-cross event took place. It is expressed in 10μs steps and may take values from 0 to 1e6 (20-bit effective). It is set by hardware and is a read-only register.

This register is used by the physical layer for being in accordance with PRIME specification. It is reserved 32-bit in data memory to be able to declare as 32-bit variable. The 20-bit register MSB is aligned to the 32-bit variable MSB, in order to simplify arithmetic calculations with time values.



14.5.28 ZERO_CROSS_CONFIG Register

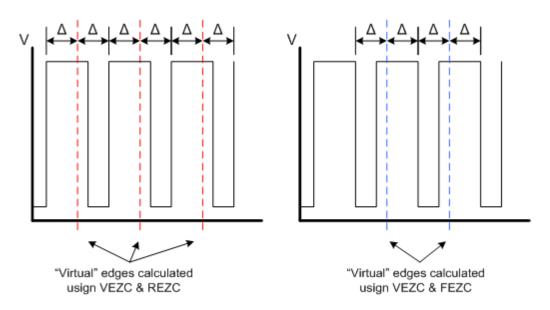
N	lame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
ZERO_CR	OSS_CONFIG						VEZC	REZC	FEZC		
Name:	ZERO_CROS	S_CONFI									
Address:	0xFE6D)xFE6D									
Access:	Read/write										
Reset:	0x06										
٠	: F	Reserved I	bits								

• VEZC:

Virtual Edge for Zero Crossing

In this bit is equal to one, the hardware calculates the middle point between two VNR edges to calculate de zero crossing.

This mode is used when the VNR signal duty cycle is different from 50%:



VECZ can be used simultaneously with RECZ or FECZ. Using the three of them at a time is not recommended.

- REZC: Rising Edge for Zero Crossing If this bit is set to '1', the hardware uses the VNR rising edges to calculate zero-crossing. FEZC and REZC can be used simultaneously.
- FEZC: Falling Edge for Zero Crossing If this bit is set to '1', the hardware uses the VNR falling edges to calculate zero-crossing. FEZC and REZC can be used simultaneously.



14.5.29 PSENSECYCLES Registers

PSENSEC	TULES Rec	Jisters										
Na	me	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0			
PSENSE	CYCLES			•		FLAG_PSENSE		D(18:16)		@0xFE7D		
						D(15:8)				@0xFE7E		
		D(7:0)										
Name:	PSENSEC	YCLES										
Address:	0xFE7D –	0xFE7F										
Access:	Read/write											
Reset:	0x00,, 0	x00;										
٠	:		Reserve	d bits								
٠	FLAG_SE	NSE:	Whenever a new power value is written in PSENSECYCLES, FLAG_PSENSE is set to'1'.									
			This flag	must be	cleared b	y software						
٠	D(17:0):		Power su	upply con	sumptior	n measurement						
			•	during N	UMMILI	s sampled (=2 S milliseconds is s				•		
			Note: The first valid value is written after NUMMILIS, and then a new valid value is written every NUMMILIS milliseconds.									
						only active when re (see Name:14.5.2		sage tran	smission	begins or		



14.5.30 MEAN Registers

Na	ime	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0			
ME	AN					FLAG_MEAN	/EAN D(18:16)			@0xFE80		
					@0xFE81							
						D(7:0)				@0xFE82		
Name:	MEAN											
Address:	0xFE80 – 0)xFE82	FE82									
Access:	Read/write											
Reset:	0x00,, 0	x00;	00;									
٠	:		Reserve	d bits								
٠	FLAG_ME	AN:	Wheneve	er a new	value is	written in MEAN, F	LAG_ME	AN is set	to '1'			
			This flag	must be	cleared b	oy software						
٠	D(17:0):		PSENSE A_NUMN Note: Th	CYCLES /ILIS reg e first val	8 and h ister in 1 lid value	erage power cons aving into accou 4.5.36). is written after NUI v NUMMILIS millise	nt the o	convergen	ce facto	r "A" (see		



14.5.31 PMAX Registers

Na	ime	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0			
PN	IAX			-		FLAG_PMAX	FLAG_PMAX D(18:16)					
						D(15:8)				@0xFE84		
			D(7:0)									
Name:	PMAX											
Address:	0xFE83 – 0	0xFE85	FE85									
Access:	Read/write											
Reset:	0x00,, 0	x00;)0;									
٠	:		Reserve	d bits								
•	FLAG_PM	AX:	Whenev	er a new	value is v	written in PMAX, F	LAG_PM	AX is set	to '1'.			
			This flag	must be	cleared b	y software						
٠	D(17:0):		As described in MAXPOT register (see 14.5.34), every time the average por consumption exceeds a user defined threshold value, the current transmission cancelled.									
			PMAX register stores the average power consumption value that has ris MAXPOT threshold.									



14.5.32 TRANS_PSENSE Register

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
TRANS	_PSENSE				TRANS_P	SENSE(7:0))		
Name:	TRANS_PSENSE								
Address:	0xFE86								
Access:	Read/write								
Reset:	0x2B								

TRANS_PSENSE: This register stores the number of 50 μs cycles to wait from the beginning of a transmission before looking for a possible power failure. This guard time is taken to avoid transient period where the measurement would be inaccurate
 Default value: 0x2B → 43 * 50 = 2.15ms



14.5.33 P_TH Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0				
P_TH						F	P_TH(18:16	5)	@0xFE87			
		P_TH(15:8)										
	P_TH(7:0)											

Name: P_TH

Address: 0xFE87 – 0xFE89

Access: Read/write

Reset: 0x07, 0xFF, 0xFF.

- --: Reserved bits
- **P_TH:** These registers contain a user defined power threshold. When the threshold value is exceeded, a low power consumption mode is automatically activated. In this low power consumption mode, the power dissipated in the transistors decreases at the expense of distortion increasing.



14.5.34 MAXPOT Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0				
ΜΑΧΡΟΤ						MA	XPOT(18:	16)	@0xFE8A			
		MAXPOT (15:8)										
				MAXPC	DT (7:0)				@0xFE8C			
									-			

Name: MAXPOT

Address: 0xFE8A - 0xFE8C

Access: Read/write

Reset: 0x07, 0xFF, 0xFF.

- --: Reserved bits
- **MAXPOT:** These registers contain a user defined power consumption threshold. When this threshold is exceeded, current transmission is cancelled.

When the threshold is exceeded, two flags are activated:

- POTFAILURE flag (see VSENSE_CONFIG in 14.5.22). This flag indicates that a power failure has occurred.
- FLAG_PMAX flag (see PMAX in 14.5.31). This flag indicates that, after a power failure, the last mean power value measured has been stored in PMAX register.

To reset both flags is enough to reset either of them, the other will be automatically reset. This will enable to start new transmissions.





14.5.35 NUMCYCLES Register

N	lame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
NUM	CYCLES				NUMCY	CLES(7:0)			
Name:	NUMCYCLES	S							
Address:	0xFE8D								
Access:	Read/write								
Reset:	0x05								
٠	NUMCYCLE		umber of o ken as va	-	neasuring	power bef	ore obtain	ing a mear	n value that c

Example1: If NUMCYCLES=5(cycles) and NUMMILIS=1(milliseconds), 5 power measurements will be taken during 1 millisecond each one . The first valid power measurement value will be output in the fifth millisecond.

Example2: If NUMCYCLES=3(cycles) and NUMMILIS=20(milliseconds), 3 power measurements will be taken during 20 milliseconds each one. The first valid power measurement value will be output after 60 milliseconds.



14.5.36 A_NUMMILIS Register

N	lame	Bit 7	Bit 6 Bit 5	Bit 4 Bit 3 Bit 2 Bit 1 B it 0
NUM	MILIS		A(1:0)	NUMMILIS(4:0)
Name:	A_NUMMILIS	i		
Address:	0xFE8E			
Access:	Read/write			
Reset:	0x21			
•	:	R	eserved bits	
٠	A(1:0):	С	onvergence Factor	
		al A	gorithm.	at sets the convergence speed of the mean calculation onvergence, while A=11 sets the slowest one. A=01,10 are
		ta m ດາ	ke into account in o ean value can be c ver a long period c	resents high dispersion values, so NUMMILIS value must be rder to select a suitable value for A. If NUMMILIS is high, the alculated slowly, because the averaging in being calculated of time. When NUMMILIS is low, the mean value must be rder to obtain more accurate values.
٠	NUMMILIS(4	•	•	ion time in milliseconds
				ent acquisition time in milliseconds.
		m	easurements will be	YCLES=5(cycles) and NUMMILIS=1(milliseconds), 5 power e taken during 1 millisecond each one .The first valid power vill be output in the fifth millisecond.
		m	easurements will be	(CLES=3(cycles) and NUMMILIS=20(milliseconds), 3 power taken during 20 milliseconds each one. The first valid power till be output after 60 milliseconds.



14.5.37 EMIT_CONFIG Register

Na	Name		Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
EMIT_C	CONFIG							TR_EMIT	TWO_H_BRIDGES
Name:	EMIT_CON	IFIG							
Address:	0xFE8F								
Access:	Read/write								
Reset:	0x03								
٠	:		Reserve	ed bits					

• TR_EMIT: Emission mode

This bit selects the emission mode (Internal Drive or External transistors bridge).

- '0': Emission is done by means of internal ATPL210 driver.
- '1': Emission is done by means of external transistors (Default).
- TWO_H_BRIDGES: This bit selects the number of semi-H-bridges in the external interface.
 - '0': There is only one semi-H-bridge in the external interface.
 - '1': There are two semi-H-bridges in the external interface and the field HIMP (AFE_CTL register) determines which one is active (Default).

Semi-H-Bridges must be connected following the table below

	TWO_H_BRIDGES='0'	TWO_H_BRIDGES='1'
EMIT1	Р	P1
EMIT2	Р	P1
EMIT3	Р	P1
EMIT4	Р	N1
EMIT5	Р	N1
EMIT6	Р	N1
EMIT7	Ν	P2
EMIT8	Ν	P2
EMIT9	Ν	P2
EMIT10	Ν	N2
EMIT11	Ν	N2
EMIT12	Ν	N2



14.5.38 AFE_CTL Register

Na	me	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0					
AFE	_CTL				HIMP	HIMP_INV	TXRX	TXRX_HARD	TXRX_INV)				
Name:	AFE_CTL													
Address:	0xFE90													
Access:	Read/write	е												
Reset:	0x10													
			Poson	/ed bits										
	: HIMP:													
·		This bit selects which branch is active when working with a two half-H-bri branches analog front end.												
			 • '0': "Low impedance" half-H-bridge is active (P2-N2). 											
			• '1': "	High imp	edance" l	nalf-H-bridge	is active	(P1-N1).						
٠	HIMP_IN	V :	HIMP	pin polari	ty control									
			This fi	eld invert	s the pola	arity of the HI	MP pin ou	utput.						
					•	ct to the pola E_CTL(4)) re	•	e external pin H changed	IMP output, t	he value				
٠	TXRX:		The va pin lev		d in this t	oit is taken by	the micr	ocontroller in or	der to set the	e TXRX				
			• '0': 1	XRX pin	output =	'0'.								
			• '1': 7	XRX pin	output =	'1'.								
٠	TXRX_HA	ARD:	TXRX	pin contr	ol									
			This fi	eld select	ts if the T	XRX pin is so	oftware/ha	ardware controll	ed.					
				TXRX pin E_CTL(2)		re controlled	. TXRX va	alue is set by T>	KRX bit field					
			• '1': 7	XRX pin	is hardwa	are controlled	ł.							
٠	TXRX_IN	V :	TXRX	pin polar	ity contro	I								
			This fi	eld invert	s the pola	arity of the TX	(RX pin o	utput						



14.5.39 R Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
R1		R1(7:0)								
R2		R2(7:0)								
R3				R3	(7:0)				0xFEA1	
R4		R4(7:0)								
R5		R5(7:0)								
R6		R6(7:0)								
R7		R7(7:0)								
R8				R8	(7:0)				0xFEA6	

Name: R1 – R8

Address: 0xFE9F - 0xFEA6

Access: Read/write

Reset: 0x60; 0x60; 0x60; 0x60; 0xFF; 0xFF; 0xFF; 0xFF.

R: The value in these registers strongly depends on the external circuit configuration.
 Atmel provides values to be used according with the design recommended in ATPL210 kits
 Please contact Atmel Power Line if different external configurations are going to be used

Recommended values (according to the configuration recommended in ATPL210 kits) R1(7:0): 0x21 R2(7:0): 0x20 R3(7:0): 0x12 R4(7:0): 0x02 R5(7:0): 0x37 R6(7:0): 0x77 R7(7:0): 0x37 R8(7:0): 0x77

Order of precedence: In the event of a conflict between the Ri(7:0) values above and Ri(7:0) values specified in the latest documentation in an ATPL210 kit, the values in the kit documentation shall take precedence.



14.5.40 PHY_ERRORS Registers

N	ame	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
PHY_	ERRORS					PHY	_ERROR	S(4:0)	
Name:	PHY_ERRO	RS							
Address:	0xFE94								
Access:	Read/write								
Reset:	0x00								
٠	:	Res	served bi	ts					
٠	PHY_ERRO	RS: Phy	sical Lay	/er Error C	Counter				
			•			the numbe ar this cou			hysical layer
		The	value st	ored in this	s register i	s cleared e	every time	the registe	er is read.



14.5.41 FFT_MODE Registers

Nar	ne	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
FFT_N	IODE			NSY	M(5:0)			CONTINUOUS	TEST_MODE_EN
Name:	FFT_MOD	Ξ							
Address:	0xFEB0								
Access:	Read/write								
Reset:	0x00								
•	NSYM:		Numb	per of sy	mbols to	transmit			
									ets as a free-running buffer,
				Ū				owing back to syn	
٠	CONTINUC	DUS:	This f	ield ena	ables/disa	ables cor	itinuous t	ransmission mod	e.
			• '0':	Continu	ious tran	smission	mode di	sabled.	
			• '1':	Continu	ious tran	smission	mode er	nabled.	
٠	TEST_MO	DE_EN:	This f	ield ena	ables/disa	ables tes	t mode		
			• '0':	Test mo	ode disat	oled.			
			• '1':	Test mo	ode enab	led.			
				-			•	•	e physical layer to fulfill with
			PRIM	_	specifica		· –	TESTMODE.requ	•
						•		•	fication). In this mode data
			•					•	HY_INI_TX, codifying each
						-			. The msb of the value is to
			indica	te an ir	nput of ze	ro when	set to '1'	. Each byte in da	ta memory contains 2 input

the next symbol data.

values for FFT, with the first value located at high bits. There are 97 input values for FFT, so many as the number of subcarriers, so there are 48 bytes and a half of the next byte used for codifying them. The other half of this byte (low bits) will be used for



14.5.42 AGC_CONFIG Register

Name	;	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
AGC_COM	NFIG		-	AGC0_POL	AGC0_VALUE	AGC0_MODE	AGC1_POL	AGC1_VALUE	AGC1_MODE
Name:	AGC_	_CONFI	G						
Address:	0xFE	B1							
Access:	Read	/write							
Reset:	0x24								

ATPL210 has implemented two Automatic Gain Control outputs in order to adjust the received signal level to a suitable range. When in "automatic" mode, both of them are set to '1' when the received signal is above a 16-bit-user-definable thresholds (AGC1_TH and AGC0_TH) in order to activate external attenuators placed in the external analog front end.

The value of these outputs is set during the beginning of a received message and is hold until the end of the message.

AGC0 and AGC1 follow different algorithms, thus using both of them ensures more accurate gain control

٠	:	Reserved bits
٠	AGC0_POL:	AGC0 polarity
		This bit sets the polarity of the AGC0 output.
		 '0': Polarity is inverted.
		 '1': Polarity is not inverted (default).
٠	AGC0_VALUE:	AGC0 output value-
		This bit stores the value wrote by the user to be the AGC0 output.
		This bit is only taken into account when AGC0 "forced" mode is active (AGC0_MODE='1').
		AGC0_POL field can invert this value.
٠	AGC0_MODE:	AGC0 mode
		This bit selects which AGC0 mode is being used
		 '0': "Automatic" Mode. AGC0 output will be managed by the MAC, depending on saturation detected in received signal. If saturation is detected, AGC0 output will be '1'. Else, AGC0 output will be '0'. AGC0_POL field can invert this value.
		(See SAT_TH registers in 14.5.43)
		 '1': "Forced" Mode. AGC0 output will be managed by the user, according to the value wrote in AGC0_VALUE field (AGC_CONFIG(4)).
٠	AGC1_POL:	AGC1 polarity
		This bit sets the polarity of the AGC1 output.
		 '0': Polarity is inverted.
		 '1': Polarity is not inverted (default).
٠	AGC1_VALUE:	AGC1 output value-
		This bit stores the value wrote by the user to be the AGC1 output.
		This bit is only taken into account when AGC1 "forced" mode is active (AGC1_MODE='1').
		AGC1_POL field can invert this value.



AGC1_MODE:

AGC1 mode

This bit selects which AGC1 mode is being used

• '0': "Automatic" Mode. AGC1 output will be managed by the MAC, depending on saturation detected in received signal. If saturation is detected, AGC1 output will be '1'. Else, AGC1 output will be '0'. AGC1_POL field can invert this value.

(See SAT_TH registers in 14.5.43)

• '1': "Forced" Mode. AGC1 output will be managed by the user, according to the value wrote in AGC1_VALUE field (AGC_CONFIG(4)).



14.5.43 SAT_TH Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
SAT_TH		SAT_TH(15:8)								
		SAT_TH(7:0)								

Name: SAT_TH

Address: 0xFEB7 – 0xFEB8

Access: Read/write

Reset: 0x40; 0x00

• **SAT_TH:** These registers store a threshold for the PLC input-signal amplitude.

If this threshold is exceeded, AGC thresholds (AGC0_TH and AGC1_TH) will be taken into account.

If this threshold is not exceeded, AGC0_TH and AGC1_TH thresholds will be ignored, thus the AGC algorithm will be never triggered.

Recommended value for Atmel reference design = 0x37AA.



14.5.44 AGC1_TH Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
AGC1_TH		AGC1_TH(15:8)								
-		AGC1_TH(7:0)								

Name: AGC1_TH

Address: 0xFE5F – 0xFE60

Access: Read/write

Reset: 0x40; 0x00

• AGC1_TH: AGC1 Threshold

These registers store the 16-bit upper threshold used by the AGC1 algorithm to determine that the input signal must be attenuated.

This threshold is only taken into account in AGC1 "automatic" mode (AGC_CONFIG.AGC1_MODE='0').

This threshold is only taken into account if SAT_TH value is exceeded.

Recommended value for Atmel reference design = 0x4A00.



14.5.45 AGC0_TH Registers

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0		
AGC0_TH		AGC0_TH(15:8)								
		AGC0_TH(7:0)								

Name: AGC0_TH

Address: 0xFEB2 – 0xFEB3

Access: Read/write

Reset: 0x10; 0x00

• AGC0_TH: AGC0 Threshold

These registers store the 16-bit upper threshold used by the AGC0 algorithm to determine that the input signal must be attenuated.

This threshold is only taken into account in AGC0 "automatic" mode (AGC_CONFIG.AGC0_MODE='0').

This threshold is only taken into account if SAT_TH value is exceeded.

Recommended value for Atmel reference design = 0x1000.



14.5.46 AGC_PADS Register

Na	ime	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	B it 0
AGC_	PADS							P46_MODE	SWITCH_AGC
Name:	AGC_PAD	S							
Address:	0xFE61								
Access:	Read/write								
Reset:	0x00								
٠	:		Reserve	ed bits					
٠	P46_MOD	E:	This fiel	d control	s the P4.6	6/T2/AGC	C1 output	pin (pin no.94).
			• '0': Pi	n no.94 v	vorks as I	P4.6/T2 c	output pin		
			● '1': Pi	n no.94 v	vorks as <i>i</i>	AGC1 ou	tput pin.		
٠	SWITCH_A	AGC:	This bit	switches	the AGC	0 and AG	GC1 outpu	uts.	
			• '0': No	ot switche	ed AGC o	utputs.			
			• '1': S\	witched A	GC outpu	uts.			



15. Electrical Characteristics

15.1 Absolute Maximum Ratings

Permanent device damage may occur if Absolute Maximum Ratings are exceeded. Functional operation should be restricted to the conditions given in the Recommended Operating Conditions section. Exposure to the Absolute Maximum Conditions for extended periods may affect device reliability.

Table 15-1. ATPL210A Absolute Maximum Ratings	Table 15-1.	ATPL210A	Absolute	Maximum	Ratings
---	-------------	----------	----------	---------	---------

Parameter	Symbol	Rating	Unit
Supply Voltage	VCC	-0.5 to 4.0	V
Input Voltage	VI	-0.5 to VCC+0.5(≤4.0V)	V
Output Voltage	vo	-0.5 to VCC+0.5(<4.0V)	V
Storage Temperature	TST	-55 to 125	°C
Junction Temperature	TJ	-40 to 125	°C
Output Current ⁽¹⁾	ю	±10 ⁽²⁾	mA

Notes: 1. DC current that continuously flows for 10ms or more, or average DC current.

2. Applies to all the pins except EMIT pins. EMIT pins should be only used according to circuit configurations recommended by Atmel Power Line.

ATTENTION observe EDS precautions



Precautions for handling electrostatic sensitive devices should be taken into account to avoid malfunction. Charged devices and circuit boards can discharge without detection



15.2 Recommended Operating Conditions

Devenueter	Cumahal		Unit		
Parameter	Symbol	Min	Тур	Мах	Unit
	VCC	3.00	3.30	3.60	
Supply Voltage	VDE0	3.00	3.30	3.60	V
	AVD	3.00	3.30	3.60	
Junction Temperature	TJ	-40	25	125	20
Ambient Temperature	ТА	-40		85	°C

Table 15-2.	ATPL210A Recommended Operating Conditions
-------------	---



15.3 DC Characteristics

Parameter	Condition	Symbol		Rating			
Parameter	Condition	Symbol	Min.	Тур.	Max.	Unit	
Supply Voltage		VCC	3.00	3.30	3.60		
H-level Input Voltage (3.3v CMOS)		VIH	2.0	-	VCC+0.3		
L-level Input Voltage (3.3v CMOS)		VIL	-0.3	-	0.8	Ň	
H-level Output Voltage	3.3v Ι/Ο ΙΟΗ=-100μΑ	VOH	VCC-0.2	-	VCC	V	
L-level Output Voltage	3.3v Ι/Ο ΙΟL=100μΑ	VOL	0	-	0.2		
H-level Output V-I Characteristics	3.3v I/O VCC=3.3±0.3	ЮН	See 15.3.2		mA		
L-level Output V-I Characteristics	3.3v I/O VCC=3.3±0.3	IOL	See 15.3.2		mA		
Internal Pull-up Resistor ⁽¹⁾	3.3v I/O	Rpu	10	33	80	kΩ	
Internal Pull-down Resistor ⁽¹⁾	3.3v I/O	Rpd	10	33	80	kΩ	
Junction Temperature		TJ	-40	-	125	°C	

Notes: 1. Only applicable to pins with internal pulling. See Table 2-1.



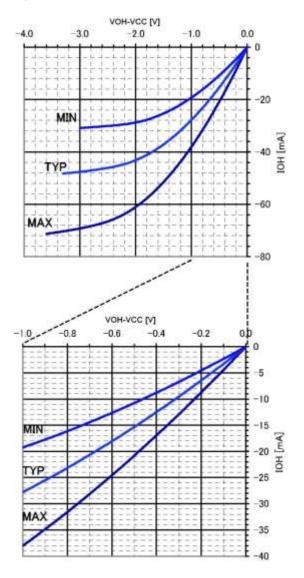
15.3.2 V-I curves

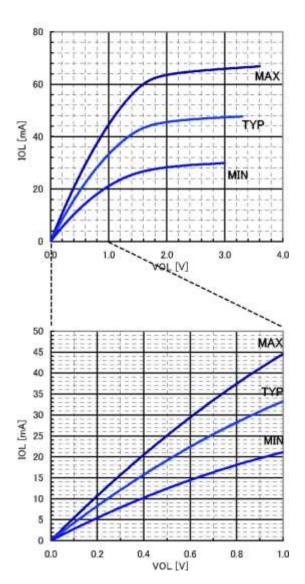
V-I Characteristics 3.3 Vstandard CMOS IO L, M type

Pins marked in Table 2-1 with Nominal Current I(mA)=±5

Condition:	MIN	Process=	Slow	Tj=	125°C	VCC=	3.0 V
	TYP	Process=Ty	pical	Tj=	25°C	VCC=	3.3 V
	MAX	Process=Fa	ast	Tj=	-40°C	VCC=	3.6 V







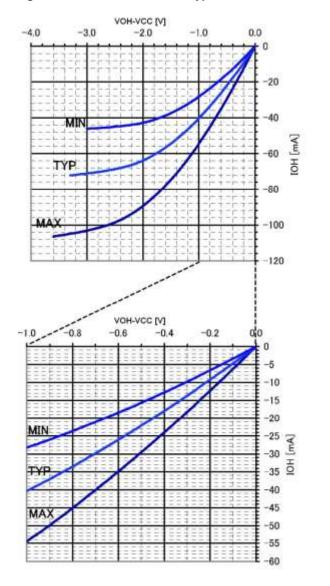


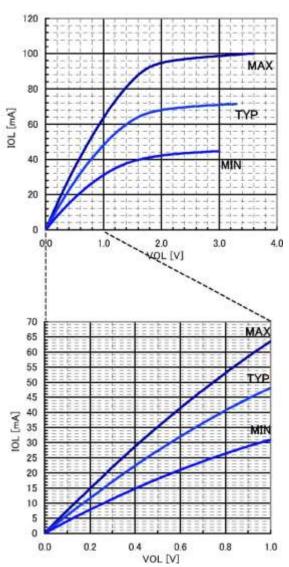
V-I Characteristics 3.3 Vstandard CMOS IO H, V type

Pins marked in Table 2-1 with Nominal Current I(mA)=±10

Condition:	MIN	Process=	Slow	Tj=	125°C	VCC=	3.0 V
	TYP	Process=Ty	pical	Tj=	25°C	VCC=	3.3 V
	MAX	Process=Fa	ast	Tj=	-40°C	VCC=	3.6 V

Figure 15-2. CMOS IO H and V type, V-I curves





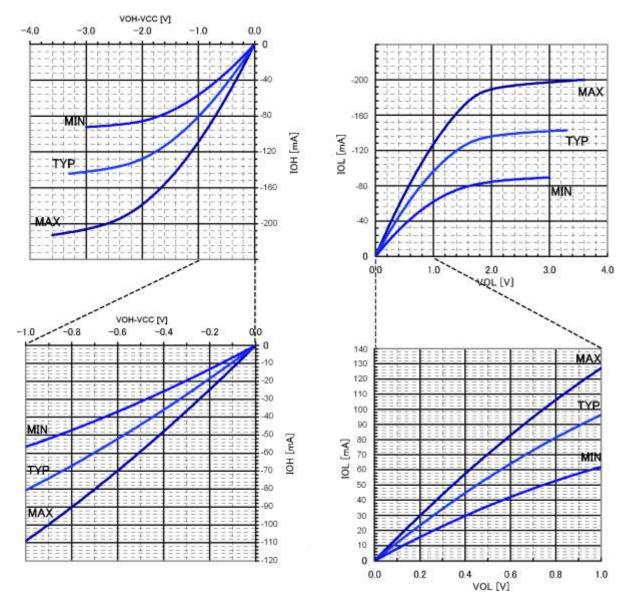


V-I Characteristics 3.3 Vstandard CMOS IO X type

Pins marked in Table 2-1 with Nominal Current I(mA)=±X

Condition:	MIN	Process=	Slow	Tj=	125°C	VCC=	3.0 V
	TYP	Process=Ty	pical	Tj=	25°C	VCC=	3.3 V
	MAX	Process=Fast		Tj=	-40°C	VCC=	3.6 V

Figure 15-3. CMOS IO X type, V-I curves





15.4 Power Consumption

Table 15-4.	ATPL210A Power	Consumption
-------------	----------------	-------------

Parameter	Condition	Symbol	Rating			Unit
Falameter		Condition Symbol	Min.	Тур.	Max.	Unit
Power Consumption	TA=25°C, VCC=3.3v	P ₂₅		260		mW
Power Consumption (worst case)	TA=85°C, VCC=3.6v	P ₈₅			355	

15.5 Thermal Data

Table 15-5. ATPL210A Thermal Data

Parameter	Symbol	LQFP144	Unit
		60 ⁽¹⁾	00.001
Thermal resistance junction-to-ambient steady state	R _{Theta-ja}	40 ⁽²⁾	°C/W

Notes: 1. Mounted on 2-layer PCB.

2. Mounted on 4-layer PCB.

Theta-ja is calculated based on a standard JEDEC defined environment and is not reliable indicator of a device's thermal performance in a non-JEDEC environment. The customer should always perform their own calculations/simulations to ensure that their system's thermal performance is sufficient.



15.6 Oscilator

Figure 15-4. External Crystal configuration

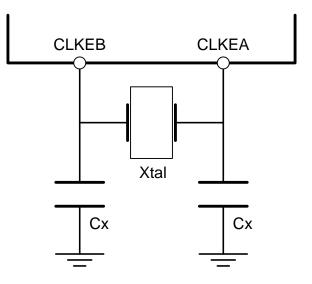


Table 15-6. External oscillator parameters

Parameter	Test	Symbol		Unit		
Falameter	Condition		Min.	Тур.	Max.	Onit
Crystal Oscillator frequency	fundamental	Xtal		20		MHz
External Oscillator Capacitance	See Figure 15- 4	Сх	5	18	30	pF
H-level Input Voltage		XVIH	2	-	VCC+0.3	V
L-level Input Voltage		XVIL	-0.3	-	0.8	v
External Oscillator Parallel Resistance		Rp	not needed			Ω
External Oscillator Series Resistance		Rs		not neede	d	\$2

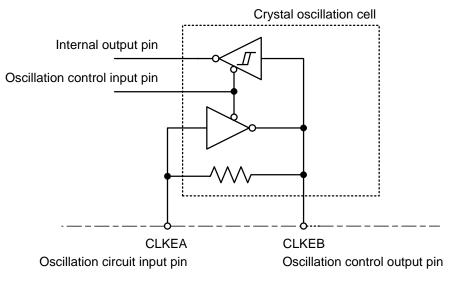
Notes: 1. The crystal should be located as close as possible to CLKEB and CLKEA pins.

2. Recommended value for Cx is 18pF. This value may depend on the specific crystal characteristics.

3. Crystal Stability/Tolerance/Ageing values must be selected according to standard PRIME requirements.





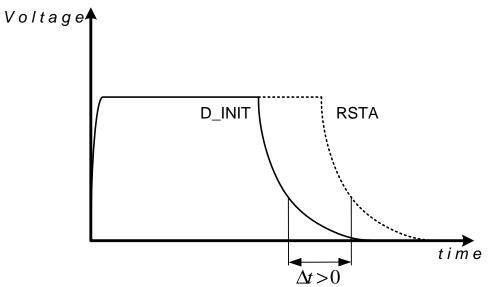




15.7 Power On

During power-on, D_INIT should be released before asynchronous reset signal RSTA in order to ensure proper system start up.





Not minimum time is required between both releases, $\Delta t > 0$, so a simple RC circuit is enough to satisfy this requirement.

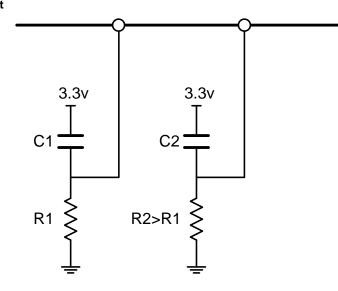
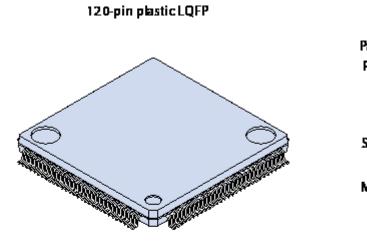


Figure 15-7. Example circuit



16. Mechanical Characteristics

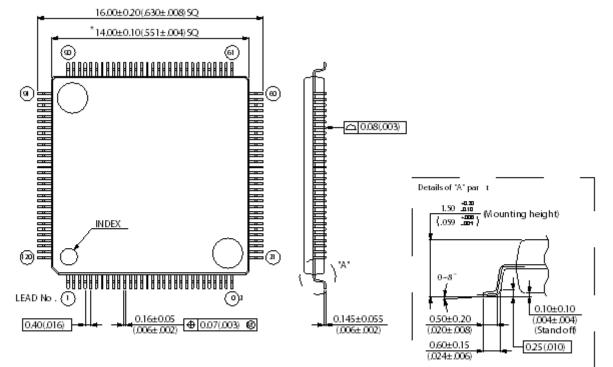
Figure 16-1. 120-lead LQFP Package Mechanical Drawing



Lead pitch	0.40 mm				
Pa ckage width · package length	14.0 mm · 14.0 mm				
Lead shape	Gullwing				
Sealing method	Plastic mold				
Mounting height	1.70 mm MAX				
Code (Reference)	P-LFQFP120-14 · 14-0.4 0				

120-pin plastic LQFP

Note 1) * :These dimensions do not include resin protrusion. Note 2) Pins width and pins thickness include plating thickness. Note 3) Pins width do not include tie bar cutting remainder.



Dimensions in mm (inches). Note: The values in parentheses are reference values.



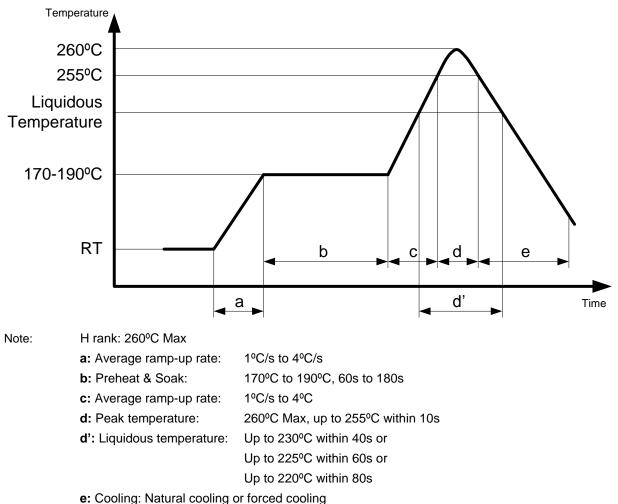
17. Recommended mounting conditions

17.1 Conditions of Standard Reflow

Table 17-1. Conditions of standard Reflow

Items	Contents				
Method	IR(Infrared I	Reflow)/Convection			
Times		2			
	Before unpacking	Please use within 2 years after production			
	From unpacking to second reflow	Within 8 days			
Floor Life	In case over period of floor life	Baking with 125°C +/- 3°C for 24hrs +2hrs/-0hrs is required. Then please use within 8 days. (please remember baking is up to 2 times)			
Floor Life Condition	Between 5°C and 30°C and also below 70%RH required. (It is preferred lower humidity in the required temp range.)				

Figure 17-1. Temperature Profile





17.2 Manual Soldering

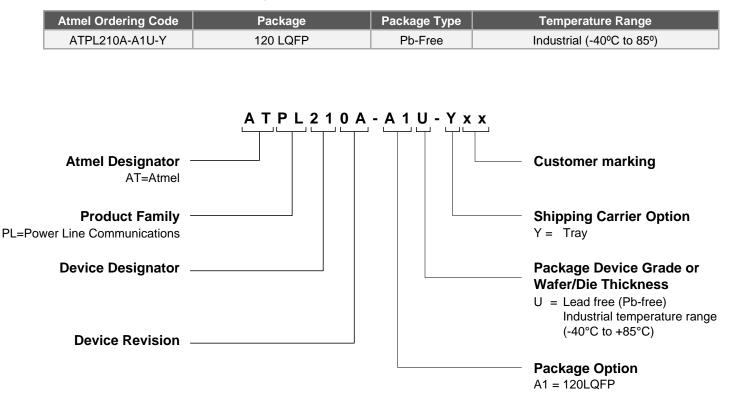
ltems	Contents	
Floor life	Before unpacking	Please use within 2 years after production
	From unpacking to Manual Soldering	Within 2 years after production (No control required for moisture adsorption because it is partial heating)
Floor life condition	Between 5°C and 30°C and also below 70%RH required. (It is preferred lower humidity in the required temp range.)	
Solder Condition	Temperature of soldering iron: Max 400°C, Time: Within 5 seconds/pin *Be careful for touching package body with iron	

Table 17-2. Conditions of Manual Soldering



18. Ordering Information

Table 18-1. Atmel ATPL210A Ordering Codes





19. Revision History

Doc. Rev.	Date	Comments
A	06/18/2012	Initial release







Enabling Unlimited Possibilities**

Atmel Corporation

1600 Technology Drive San Jose, CA 95110 USA **Tel:** (+1)(408) 441-0311 **Fax:** (+1)(408) 487-2600 www.atmel.com

Atmel Asia Limited

Unit 01-5 & 16, 19F BEA Tower, Millennium City 5 418 Kwun Tong Road Kwun Tong, Kowloon HONG KONG Tel: (+852) 2245-6100 Fax: (+852) 2722-1369

Atmel Munich GmbH

Business Campus Parkring 4 D-85748 Garching b. Munich GERMANY **Tel:** (+49) 89-31970-0 **Fax:** (+49) 89-3194621 Atmel Japan G.K.

16F Shin-Osaki Kangyo Building 1-6-4 Osaki Shinagawa-ku, Tokyo 141-0032 JAPAN **Tel:** (+81)(3) 6417-0300 **Fax:** (+81)(3) 6417-0370

© 2012 Atmel Corporation. All rights reserved. / Rev.: 43005A-ATPL-06/12

Atmel[®], Atmel logo and combinations thereof, and others are registered trademarks or trademarks of Atmel Corporation or its subsidiaries. Other terms and product names may be trademarks of others.

Disclaimer: The information in this document is provided in connection with Atmel products. No license, express or implied, by estoppel or otherwise, to any intellectual property right is granted by this document or in connection with the sale of Atmel products. EXCEPT AS SET FORTH IN THE ATMEL TERMS AND CONDITIONS OF SALES LOCATED ON THE ATMEL WEBSITE, ATMEL ASSUMES NO LIABILITY WHATSOEVER AND DISCLAIMS ANY EXPRESS, IMPLIED OR STATUTORY WARRANTY RELATING TO ITS PRODUCTS INCLUDING, BUT NOT LIMITED TO, THE IMPLIED WARRANTY OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE, OR NON-INFRINGEMENT. IN NO EVENT SHALL ATMEL BE LIABLE FOR ANY DIRECT, INDIRECT, CONSEQUENTIAL, PUNITIVE, SPECIAL OR INCIDENTAL DAMAGES (INCLUDING, WITHOUT LIMITATION, DAMAGES FOR LOSS AND PROFITS, BUSINESS INTERRUPTION, OR LOSS OF INFORMATION) ARISING OUT OF THE USE OR INABILITY TO USE THIS DOCUMENT, EVEN IF ATMEL HAS BEEN ADVISED OF THE POSSIBILITY OF SUCH DAMAGES. Atmel makes no representations or warranties with respect to the accuracy or completeness of the contents of this document and reserves the right to make changes to specifications and products descriptions at any time without notice. Atmel products are not intended, authorized, or warranted for use as components in applications intended to support or sustain life.

Mouser Electronics

Authorized Distributor

Click to View Pricing, Inventory, Delivery & Lifecycle Information:

Microchip:

ATPL210A-A1U-Y